



DoS vulnerability of TCP by acknowledging not received segments

draft-azcorra-tcpm-tcp-blind-ack-dos-01

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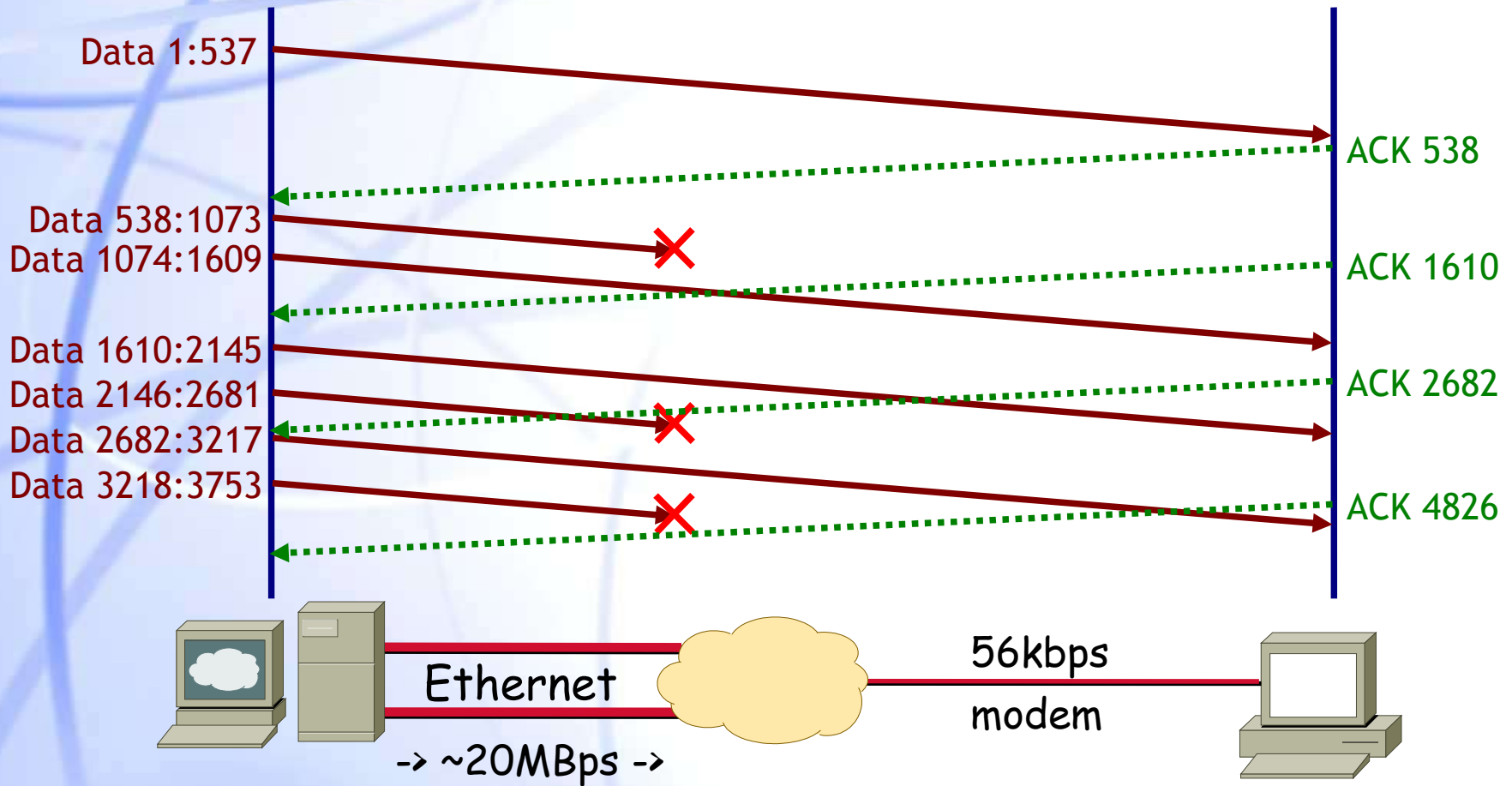
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Expeditious Blind ACK DoS

TCP “normal” sender

TCP “malicious” receiver



DoS attack to information servers!!



Proposed Solutions (I)

- **These modifications could be used independently or together**
- **1. Starting a slow start procedure**
 - ◆ **A server MUST start the slow start procedure if it receives an ACK for unsent data, but within the transmitter send window:**
 - **$SND.NXT < SEG.ACK \leq SND.UNA + SND.WND$**
 - ◆ **Intended to penalize an attacker**
 - ◆ **This modification couldn't be used by a malicious THIRD party to attack a legitimate TCP connection**



Proposed Solutions (II)

- **2. Matching SEG.ACK and (SEG.SEQ + SEG.LEN)**
 - ◆ **A server SHOULD randomize segment boundaries in the range [MS, a*MS]**
 - MS = current maximum segment size
 - a = number in the range [0,1]
 - ◆ **A server should accept an ACK, but NOT increase its CWND variable, if the SEG.ACK fulfils these 2 conditions simultaneously:**
 - $SND.UNA \leq SEG.ACK \leq SND.NXT$
 - $SEG.ACK \neq (SEG.SEQ + SEG.LEN)$ of one of the unacknowledged segments that have been sent

