#### draft-fujiwara-dnsop-bad-dns-auth-01

## IETF61 DNSOP 2004/11/07

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#### **Topics**

- Protecting cache server against misconfiguration of DNS authorititive servers
  - •We had a terrible cache server overload.
  - This is caused by misconfigured authoritative servers.
  - OProtection method is needed for ISP cache server.
- ☐ This topic has more issues about DNS protocol.
  - ODNS Response size consideration
  - ○EDNS0 implementation status
  - OEDNS0(with fragmentation) and IPv6
  - ODNS anycast vs TCP query (not mentioned in the draft -01)

## Cache server overload (1)

- ☐ There are some authoritative server misconfiguration
  - Olarge response size RRSet
    - ▶ Many(32) PTRs in one IP address
  - ono EDNS0
  - OTCP filtering at authoritative servers
- □ and this RRSet is major IP address, many clients query this address frequently.
- □What happenes as a result?
  - Oln every query, truncation occurs.
  - OAt that time, cache server queries again by TCP.
  - OBut TCP is blocked by filter.
  - Then cache server has many 'stocled' TCP SYN\_SENT status and makes high load.

## Cache server overload (2)

- □ Authoritative server misconfiguration can create significant overloads on cache servers.
  - OThis behavior was found through the observation of query traffic to/from ISP cache servers.
  - OAnd we reported it in NANOG32 meeting in October.
- □ Attacker can make a DoS attack to ISP cache servers using this problem.
  - OAttacker prepares an authoritative server and a RRSet with this problem.
  - OAttacker sends a lot of queries with spoofed source addresses, as if they are sent from various clients.
- □ From the ISP users view, failure of DNS cache server is almost equal to the failure of the Internet service itself.
- □We should protect DNS cache servers.

## How to decrease TCP sessions in Resolver server

- □One idea: do not query by TCP when truncation
  - Olt reduces TCP sessions.
  - OBut the answer which is supposed to be able to get it properly if it listens with TCP can't be resolved.
  - oit cannot cache any data (RFC 2308).
  - All resolving request, the cache server queries to all the authoritative servers by UDP.
  - ○More, it may violate RFC2181.
- □ Needs new cache/resolver server algorithm

## Cache/Resolver server algorithm improvement

#### □We propose

- OAs before, the cache server queries by UDP (w/wo EDNS0) and if TC bit is set, the cache server queries again the authoritative server by TCP.
- o(new)When queries for all authoritative servers are unsuccessful, the cache server caches that RRSet(name,class,type) as unresolvable.
- O(new)Next query for the same RRSet from stub resolvers, the cache server does not query to authoritative servers and answers "unresolvable".
- o(optional)Cache server marks misconfigured servers which does not respond TCP. (equivalent to BIND9's EDNS0 capability database.)

#### □ Protocol consideration

#### □RFC2308 section 7 - Other Negative Responses

• This section does not mention about TCP filtering.

▶UDP is OK

⊳no answer by TCP, no TCP reset

ORFC 2308 7.1: "... In either case a resolver MAY cache a server failure response. If it does so it MUST NOT cache it for longer than five (5) minutes, and it MUST be cached against the specific query tuple <query name, type, class, server IP address>."

#### □But

- ○5 minutes is too small to protect from DoS.
- In our case, authoritative server's misconfiguration lasted in about a half year.

#### □Our proposal

• For protecting cache servers, we recommend to cache unresolvable information for several hours.

## DNS Response size consideration

- □DNS response size lower than 512 octets
  - osafe with Original UDP DNS protocol
- □512 < DNS response size <= 1280 (IP/UDP)header size (1200? octets)
  - osafe with EDNS0 without IPv6 fragmentation
  - ○(on the present Internet, IPv4 is the same as IPv6.)
  - ○TCP is OK
- □1200? < DNS response size
  - OEDNS0 needs IP/IPv6 fragmentation
  - ○TCP is OK

## EDNS0 implementation status

- □ Question
  - ONow, EDNS0 requirement is "SHOULD". Is this OK?
  - ○When will EDNS0 requirement be "MUST" ?
- □This discussion is need for enum-wg.
- □RRSet may be large in ENUM.

## EDNS0(with fragmentation) and IPv6

#### □ According to RFC3226 section 3

O"All RFC 2535 and RFC 2874 compliant entities MUST be able to handle fragmented IPv4 and IPv6 UDP packets." (to support EDNS0 with large response size)

#### □But RFC2460 "IPv6 Specification" section 5

"the use of such fragmentation is discouraged in any application that is able to adjust its packets to fit the measured path MTU."

#### □ Question

EDNS0 with large response size requires IPv4 and IPv6 fragmentation. Is it OK? (I think OK.)

# DNS anycast vs TCP query (not mentioned in the draft -01)

- □TCP queries may work on DNS anycast with BGP.
  - Routing information may be stable for a short time.
  - Equal cost multi path doesn't occur in principle.
    - ▶But ECMP problem occurs in the Internet, more investigation is necessary.
  - OTCP communication may work for a short time.
  - ODNS queries using TCP is completed in a short time.
- □DNS anycast with IGP
  - "Equal cost multi path" problem can be solved with per flow routing.
- □ Need more consideration.

## A minimal requirement may be

- □DNS response size exceeds 512 octets
  - Othe authoritative name servers MUST permit TCP queries
  - or MUST support EDNS0
- □DNS response size exceeds 1200 octets
  - Othe authoritative name servers MUST permit TCP queries
  - ○AND MUST support EDNS0

## Summary

- □TCP support for DNS is now mandatory, but there are many authoritative servers which do not support TCP.
- □EDNS0 has IP/IPv6 fragment issues.
- □Still need for protection mechanism for DNS cache server.
- □This I-D should be separated as two I-Ds.
  - Negative cache issue
  - OToday's DNS requirements

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#### Questions

- □need Comments
- □ Please discuss in dnsop mailing list.