

# SHA-1 Hash Function Replacement

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# What is the problem?

- All crypto hashes deployed today – **broken**
  - To variable extent, but still...
- Strongest hash currently deployed – SHA-1
  - Design based on intuition, not science
  - **Practical attacks are expected within one year**
    - **Either attack complexity will improve to  $2^{50}$  (predicted lower bound)**
    - **Or distributed (and/or supercomputer) search will succeed at  $2^{63}$  attack**
    - ***Or both?***

# Why not SHA256?

- Ultimately SHA256 planned as solution **but**
  - Again intuition-based design that failed twice (SHA0 & SHA1), no known lower bounds
  - Requires radically new implementation
  - Different parameter size, etc.
  - Deployment expected by 2010 – not soon enough!
  - Performance sucks (and will for a few years)
  - Shares design weakness with SHA-1, plus
    - Non-linear code – security by confusion, not science
    - Non-linearity without analysis can lead to disastrous attacks
  - If truncated to 160 bits – problems compound

# Replace SHA-1 with what?

- **Leading attack against SHA-1 – differential**
  - Offers practical method of finding collisions
  - All other practical attacks rely on this one
- **Math shows why differential attack possible**
  - Weak key schedule (message expansion)
  - Low minimum Hamming distance
- **Math also shows how to foil this attack**
  - And therefore invalidate other attacks as well
- **SHA1-IME is implementation of this defense**

# Structure of SHA1-IME

- Same as SHA-1 in FIPS 180-1 and FIPS 180-2
- Minor change to message expansion
  - Old code (part of message expansion):
    1. *for*(*t* = 16; *t* < 80; *t*++)
    2.  $W[t] = \text{ROL1}(W[t-3] \wedge W[t-8] \wedge W[t-14] \wedge W[t-16]);$
  - New code:
    1. *for*(*t* = 16; *t* < 36; *t*++)
    2.  $W[t] = (W[t-3] \wedge W[t-8] \wedge W[t-14] \wedge W[t-16]) \wedge$
    3.  $\text{ROL13}(W[t-1] \wedge W[t-2] \wedge W[t-15]);$
    4. *for*(*t* = 36; *t* < 80; *t*++)
    5.  $W[t] = (W[t-3] \wedge W[t-8] \wedge W[t-14] \wedge W[t-16]) \wedge$
    6.  $\text{ROL13}, (W[t-1] \wedge W[t-2] \wedge W[t-15] \wedge W[t-20]);$
- Provably secure against differential attacks



# Deploy SHA1-IME because

- SHA1-IME leaves API and PKCS unchanged (same input and output size)
- Performance hit minor – about 5% in software (possibly 10% in hardware)
- SHA1-IME is provably secure – proven lower bound on collision probability
- Differential attack estimated  $2^{150}$  probability
- SHA1-IME – easiest to get FIPS certification if you already certified SHA1
  - As it is a small change to already-certified FIPS 180-1, process much faster
- Code change miniscule – easier to do
  - Both software and firmware (ASIC may be in trouble ☺)

**NO PATENTS!**

# References

- Uri Blumenthal [uri.blumenthal@intel.com](mailto:uri.blumenthal@intel.com)
- Charanjit Jutla [csjutla@watson.ibm.com](mailto:csjutla@watson.ibm.com)
  - ☞ Please direct math questions to Charanjit ☺
- Anindya Patthak [patthak@gmail.com](mailto:patthak@gmail.com)
- Specification in draft-irtf-cfrg-sha1-ime-00.txt
  - URL to follow, also being submitted to CFRG

# THANK YOU!