Local Management of Trust Anchors (for the RPKI)

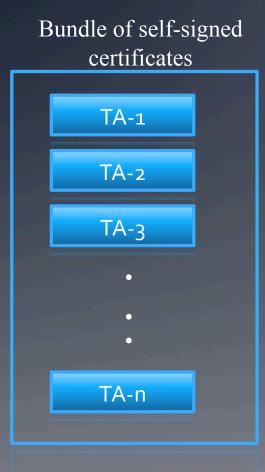
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### Local TA Management

- In principle, every RP should be able to locally control the set of TAs that it will employ
- In practice, most PKI applications do not provide good, local TA management capabilities
- The common form of a TA, a self-signed certificate, also limits a user's ability to impose constraints on it and on subordinate certificates
- The mechanisms described here are from a document that is now a WG item in SIDR, to provide a local TA management capability for the Resource PKI (RPKI)

# **Typical TA Configuration**



- In common practice, each selfsigned certificate contains no extensions that constrain it
- 5280 path validation algorithm would not impose such constraints, because they appear in a TA
- The "scope" of each TA may overlap, e.g., re name constraints
- A confusing (& dangerous) model for RPs

## Our Model: The RP is the TA!

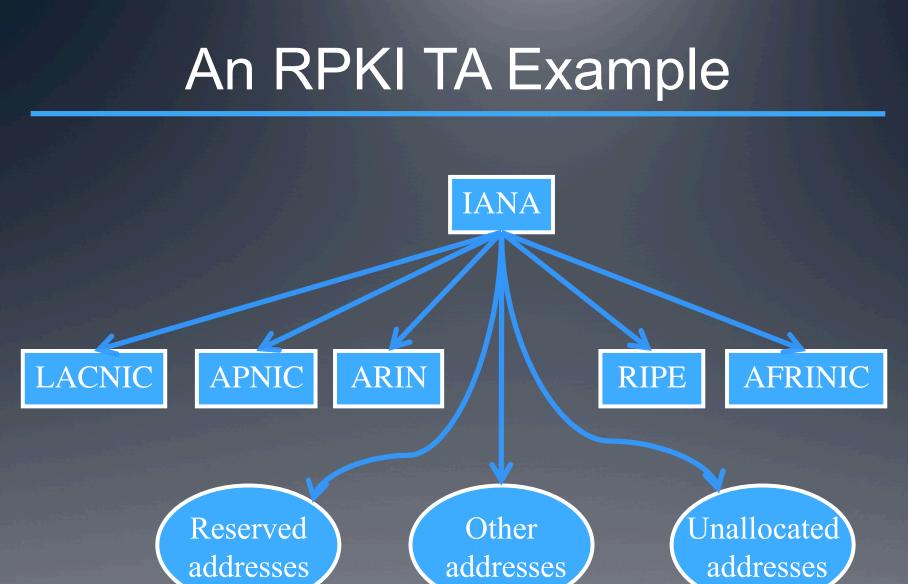
- The model we are pursuing calls for each RP to recognize exactly one TA, itself!
- The RP imports putative TAs (in the form of selfsigned certificates) and re-issues them under itself
- In so doing, the RP is empowered to insert constraints
- Because these certificates are now one step below the TA, normal certificate path processing (ala 5280) will impose those constraints (e.g., path length, policy, name, and RFC 3779)

## Why the RPKI Needs This

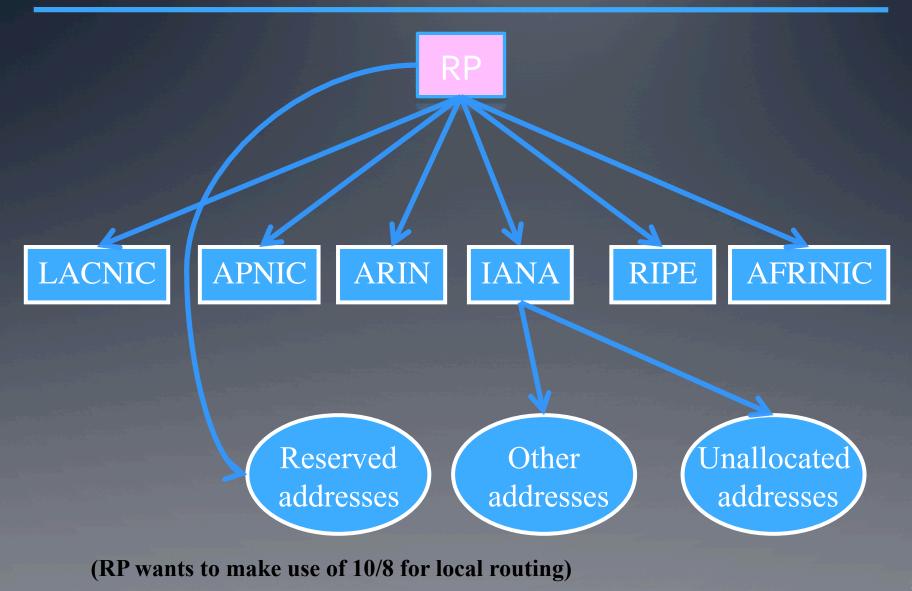
- The RPKI makes use of certificates (CA and EE) that contain Internet Number Resource (INR) extensions, as defined in RFC 3779
- The validation algorithm defined for these extensions requires that a path conform to subset rules, analogous to the name constraints extension, <u>all the</u> way to the TA for the path
- 3779 extensions imply that the RPKI is a hierarchy
- Some RPs may need to override the RPKI nominal hierarchy, e.g., to deal with RFC 1918 addresses, or for security reasons

## Making this Work in the RPKI

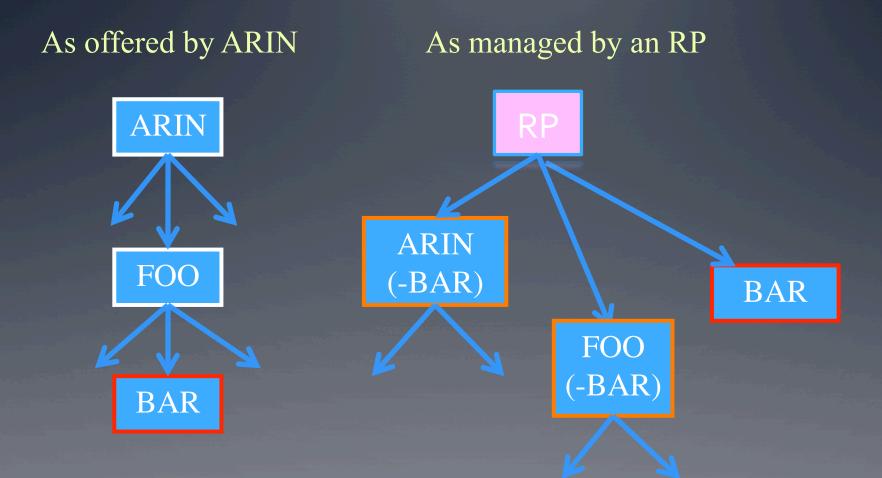
- We will need to be able to create new certificates, often with modified RFC 3779 extensions
- To make this work
  - The RP's TA certificate must contain RFC 3779 extensions encompassing <u>all</u> addresses and <u>all</u> AS#s
  - The RP TA re-issues certificates with new 3779 extensions
    - Delete overlapping 3779 data as needed
    - Re-issuing targeted certificates directly under the RP TA
    - Re-parenting ancestors of re-issued certificates under the RP TA
  - The RP can also override certain fields of the re-issued certificate using a "constraints file"



## **RPKI** with Local Control



## A More Detailed Example



(RP trusts its own knowledge of BAR's address allocation and does not want any action by ARIN or FOO to override that knowledge)

### Elements of the Solution

#### • Constraints file

- Resource re-writing algorithm
  - Target processing
  - Ancestor processing
  - Tree processing
  - TA re-homing
- Path discovery
- Revocation
- Expiration

### Constraints File

- The RP creates (or acquires from an authoritative source) a constraints file specifying IP address space and AS# resources for target certificates
  - Certificates are specified by SKI, thus the constraints file must be updated when the targets rekey
- The constraints file also allows the RP to control rewriting certain fields in the re-issued certificates
  - Validity dates
  - CRLDP
  - AIA
  - Policy Qualifier OID

## **Resource Rewriting Algorithm**

- There are four stages to the algorithm
  - Target processing
    - Certificates that match a given SKI have their resources rewritten to those specified in the constraints file
  - Ancestor processing
    - Ancestors of targets are processed to ensure RFC 3779 rule compliance (remove target certificate resources)
  - Tree processing
    - The entire tree of certificates is searched, and certificates with resources that conflict with any target resources are modified to remove the conflict
  - TA re-issuing
    - All TAs in the original hierarchy are re-issued under the RP's TA

### Implications of this Model

- This algorithm creates two parallel hierarchies: the original certificate hierarchy and the para-certificate hierarchy
- There are implications for path discovery, since a certificate can now have an original parent and a para-parent
- There are implications for revocation
- There are also implications for expiration, since the constraints file allows rewriting the validity interval of para-certificates

### Path Discovery

- Path discovery prefers the para-certificate hierarchy
- If a certificate has a para-parent, that para-parent will be used to form the certificate path
- If the certificate has only an original parent, but that parent was a target specified in the constraints file, or an ancestor of such a target, then path discovery fails
  - This can occur if the RP has revoked the paracertificate, the original certificate is still present, and the Local TA tool has not yet been run to regenerate the para-certificate
- If the certificate has only an original parent, and the parent is not a target, or the ancestor of a target, path discovery can proceed up the original chain

### Revocation

- The original hierarchy and the para-hierarchy are disjoint; revocation of a certificate in one does not affect the other
- Para-certificates are all issued by the RP, so only the RP can revoke them
- Original certificates can still be revoked by their issuers
- Because of the modified path discovery rule, revocation of any para-certificate will cause path discovery to fail until the para-certificate has been replaced or regenerated

### Expiration

- The constraints file allows the RP to specify notBefore and notAfter for all para-certificates
  - This is a global rewrite rule, not a per-certificate rewrite rule
- As a result, expiration of the original certificate does not necessarily imply that the corresponding paracertificate expires at the same time
- Expiration of a para-certificate affects path discovery in the same way as revocation of a para-certificate

# Questions?

