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# Efficient Secure BGP AS Path using **FS-BGP**

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# Outline

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- **Introduction**
  - **FS-BGP: Fast Secure BGP**
  - **Terminology**
  - **Quick review of S-BGP**
- FS-BGP
- Evaluation
- Discussion

# FS-BGP: **F**ast **S**ecure BGP

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- How to secure the path
  - CSA (**Critical path Segment Attestation**) to secure the AS path
  - SPP (**Suppressed Path Padding**) to protect the optimal path and prevent effective hijacking
- Security
  - All the authenticated paths are **feasible path**
  - Achieves **similar level of security as S-BGP**
- Computational cost (on backbone router)
  - Signing cost: **~0.6%** of S-BGP
  - Verification cost: **~3.9%** of S-BGP

# Terminology (1)

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- Feasible Path
  - Exist in the AS-level graph, and satisfies **import and export policies** of all ASes along the path
- Unfeasible Path
  - (1) Paths do **NOT** exist in the DAG
  - (2) Paths **violate import and export routing policies**

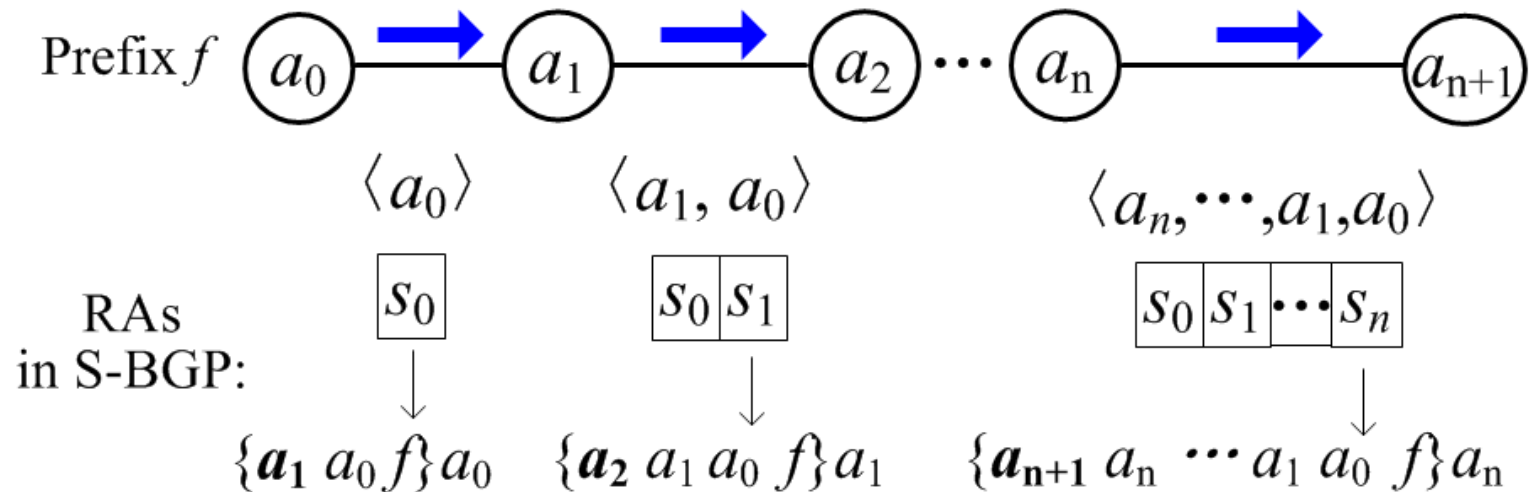
# Terminology (2)

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- Three categories of Feasible Path
  - **Outdated Path**: path announced but temporarily not available
  - **Current Path**: currently using and announcing path
  - **Not announced Path**: feasible but not announced, because BGP only announce the current optimal path each time

# Signatures in S-BGP

- Route Attestations (RA) to secure the path



# Pros and Cons of S-BGP

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- Actually signed the whole path, including the recipient AS
- **Pros:** the most secure schema
- **Cons**
  - Unbearable computational cost, so many paths.
  - Long Exp-date: unable to defend replay attack
  - Short Exp-date: destroy the whole system

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- **FS-BGP: Fast Secure BGP**
  - **CSA: Critical Segment Attestation**
  - SPP: Suppressed Path Padding
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# Announcement Restrictions in BGP

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- Best route announcing
  - **Temporary** restriction
  - Local preference and other metrics
- Selective import & export policy
  - **Persistent** restriction
  - ***Neighbor based import and export:***  
contracts (\$\$) are between neighbor Ases
  - Feasible path: exist in AS-level graph & obey the policy

# Critical Path Segment

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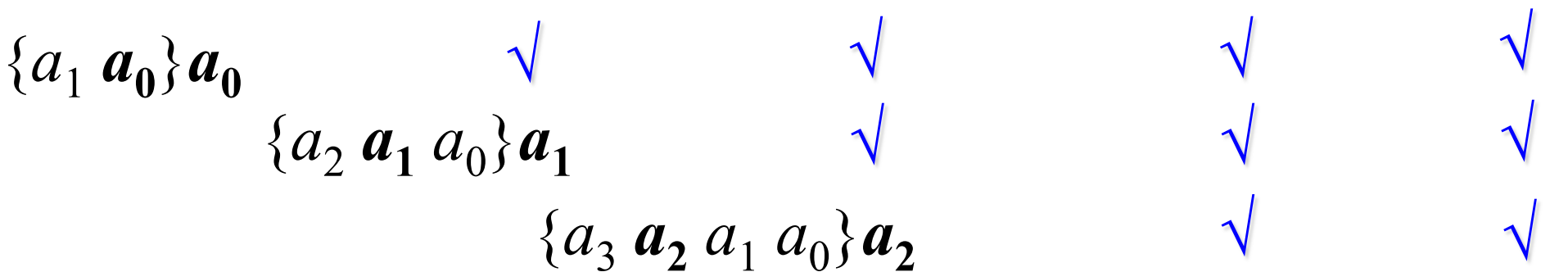
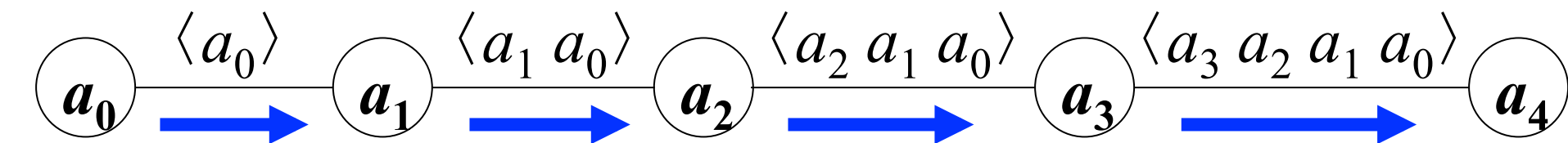
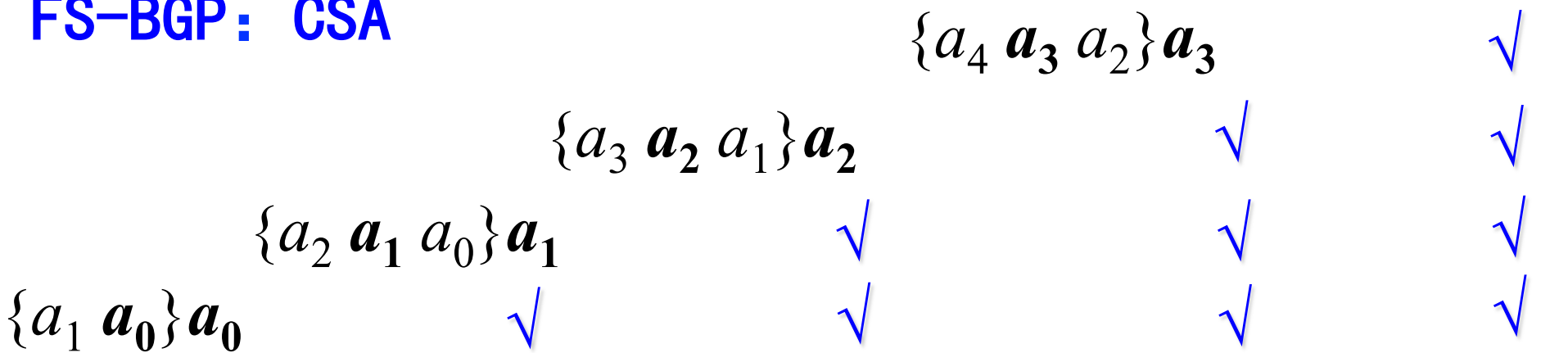
- In path:  $p_n = \langle a_{n+1}, a_n, \dots, a_0 \rangle$ , the Critical Path Segment  $c_i$  owned by  $a_i$  is

$$c_i = \begin{cases} \langle a_1, a_0 \rangle & \text{for } i = 0 \\ \langle a_{i+1}, a_i, a_{i-1} \rangle & \text{for } 0 < i \leq n \end{cases}$$

- Those adjacent AS triples actually describes part of routing policy of the corresponding owner
  - $c_i = \langle a_{i+1}, a_i, a_{i-1} \rangle$  means  $a_i$  can (**and already**) announce routes to  $a_{i+1}$  which are import from  $a_{i-1}$
  - If every owner sings the critical segment in a current announcing path, the consequence ASes will be able to verify the whole path

$\{msg\}a_i$ : signature of  $msg$  signed by  $a_i$

### FS-BGP: CSA



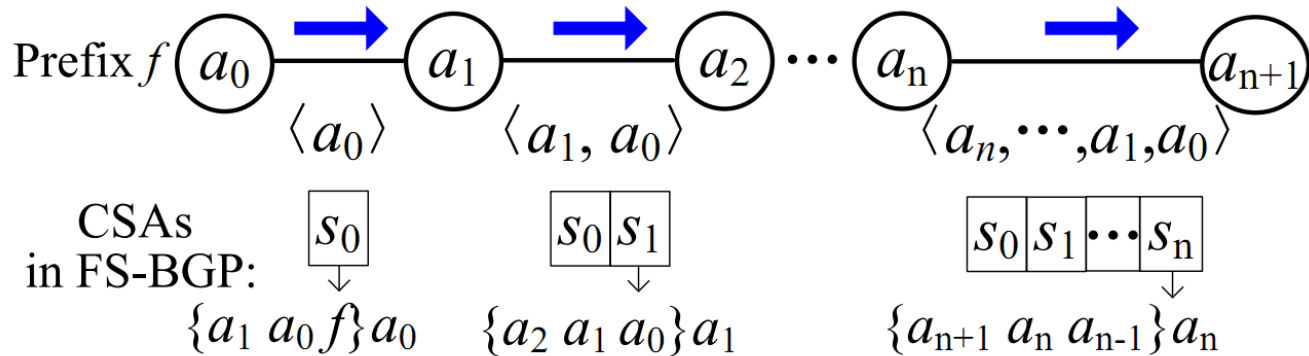
### S-BGP: RA

$\{a_4 a_3 a_2 a_1 a_0\}a_3$  ✓

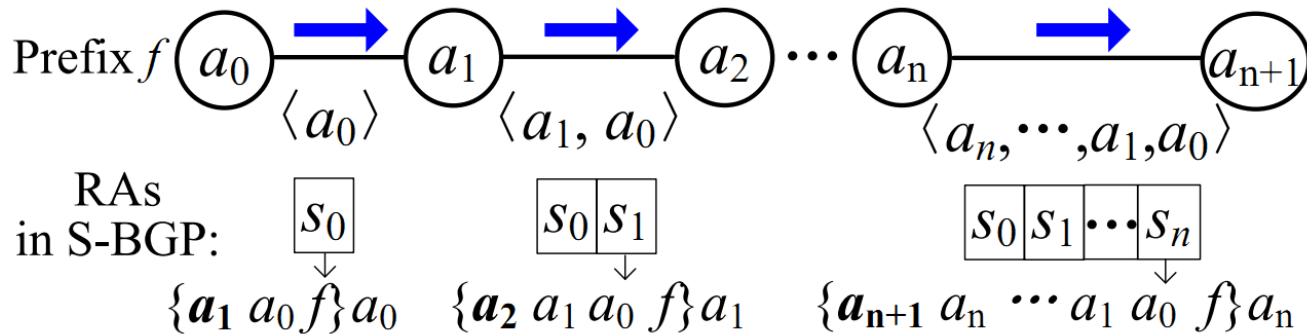
# Signatures in FS-BGP and S-BGP

Signatures for the path:  $p_n = \langle a_{n+1}, a_n, a_{n-1}, \dots, a_0 \rangle$

**FS-BGP**

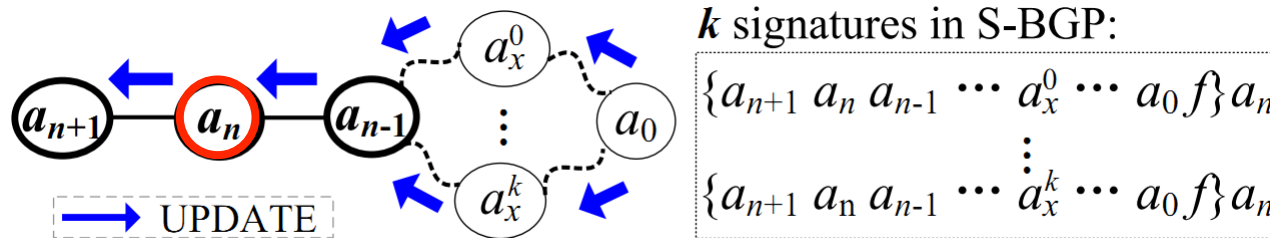


**S-BGP**



# Cost Reduction

- (# total critical segment)  $\ll$  (# total AS path)
- If we use a small cache, the cost will be sharply decreased



- S-BGP:  $a_n$  receives  $k$  paths, signs  $k$  signatures
- FS-BGP:  $a_n$  receives  $k$  paths, signs **1** signature

# Outline

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- Introduction
- **FS-BGP: Fast Secure BGP**
  - CSA: Critical Segment Attestation
  - **SPP: Suppressed Path Padding (Optional)**
- Evaluation
- Discussion

# CSA achieves Feasible Path Authentication

- Paths can be verified in FS-BGP are all feasible paths [Theorem 1]

Signed paths  
in S-BGP

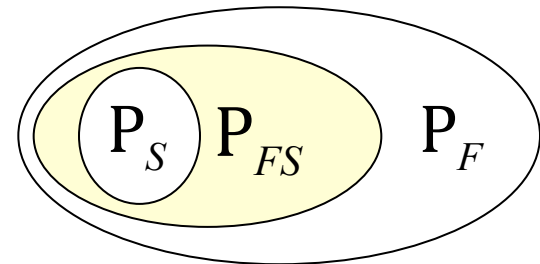
$P_S$

Signed paths  
in **FS-BGP**

$\subset$   $P_{FS}$

All feasible  
paths

$\subset$   $P_F$



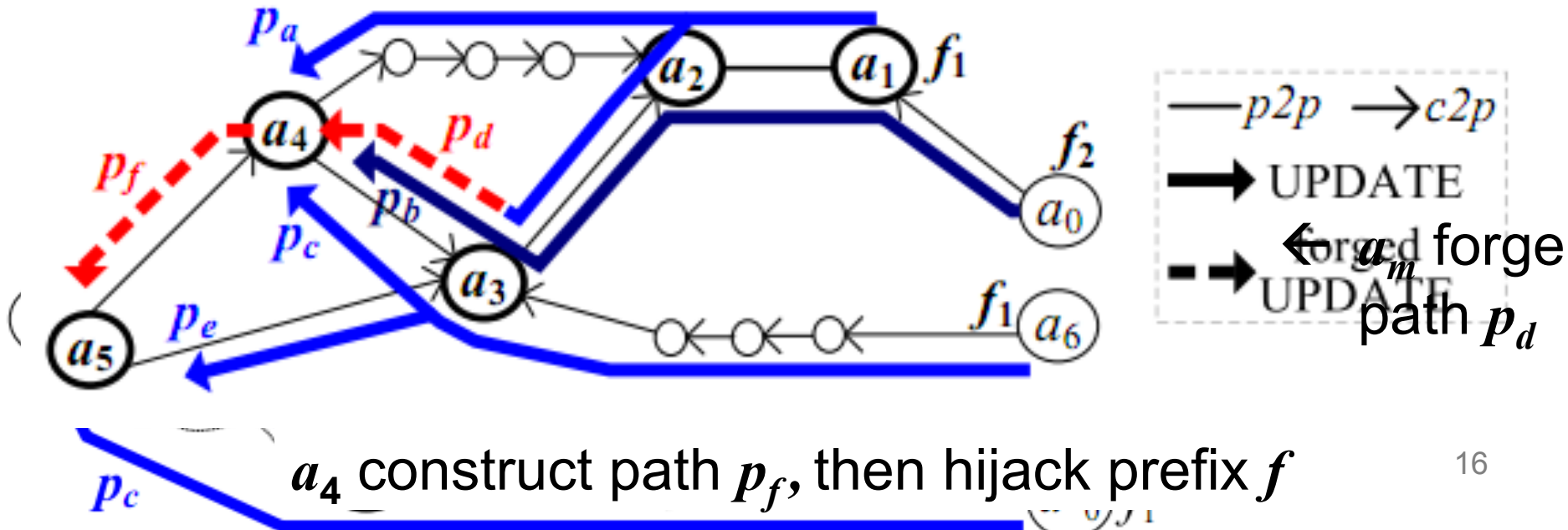
1. Outdated path
2. Current path

1. Outdated path
2. Current path
3. **Revealed path**

1. Outdated path
2. Current path
3. **All not announced path**

# Forge a path in FS-BGP is possible

- **Forged path (Revealed path)** in FS-BGP
  - Using authenticated path segments, manipulator can construct forged path, which is feasible but currently not announced.





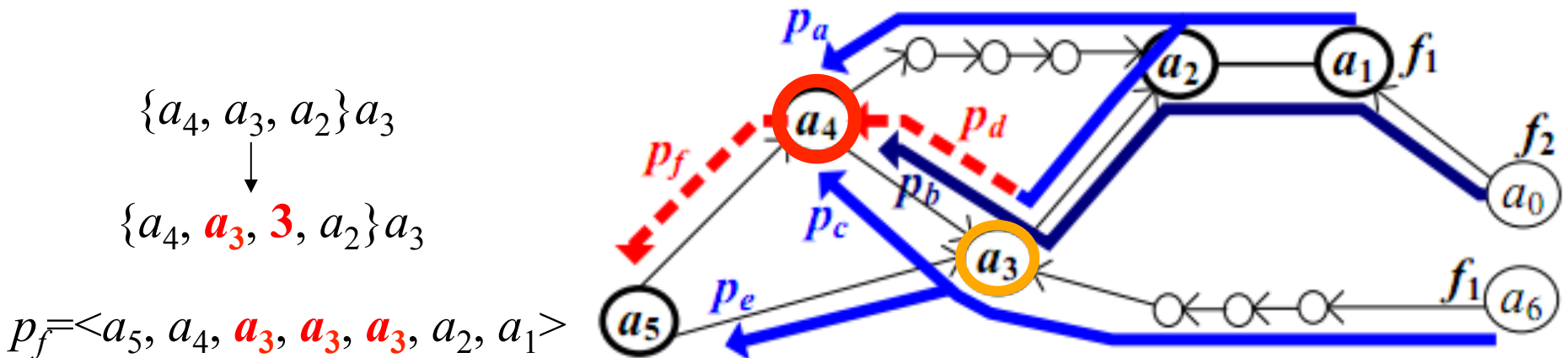
# Conditions of **Effective Hijacking**

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- **Effective hijacking**: the traffic is not forwarded by the attacker under normal status.
- (1) Forged path is **still feasible**, and only temporarily not received by the attacker!
- (2) Forge a path in FS-BGP is **very difficult**
  - Must be constructed using received authenticated path segments
  - Must not be announced by the intermediate AS
  - Can NOT be shorter than 5 hops [Theorem 2]
- (3) Only **short enough forge-path** can be used for a effective hijacking [Theorem 3]

# Prevent Effective Hijacking

- Using ASPP, can guarantee that attacker can not concatenate short enough forge path
- Short enough**: shorter than the optimal path (longest live-time)



# SPP: Suppressed Path Padding

- Suppressed Path: paths with lower local preference in the decision process
- Suppressed path may shorter than optimal path

Compute  $k_i$  :

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Algorithm 1 Suppressed Path Padding

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input: local ACS  $a_i$ , neighbor ACS  $a_j$

output:  $k_i$ : Path categories: e

1: ~~if  $a_i >_1$  then~~ **1. Suppressed Path**

2: return

3:  $k_i \leftarrow 1$

4: for all  $p \in \text{neighbors}(i)$  **2. Sub-optimal Path**

5:  $opt(p) \leftarrow \text{local\_preference}(p)$

6: ~~if  $PL(p) < PL(i)$  then~~ **3. Optimal Path**

7:  $k_i \leftarrow k_i + 1$

8: return  $k_i$

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Basic decision process:

1. Highest Local Preference (LP)
2. Shortest Path Length (PL)
3. Tie Breaks (TB)

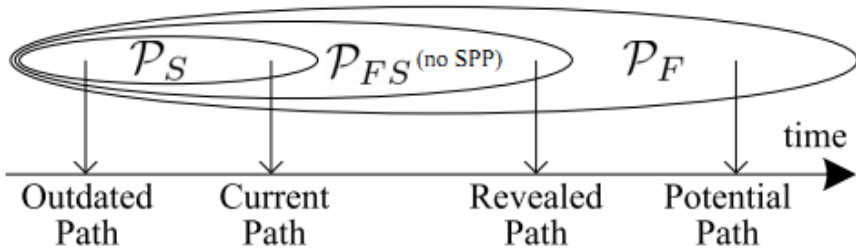
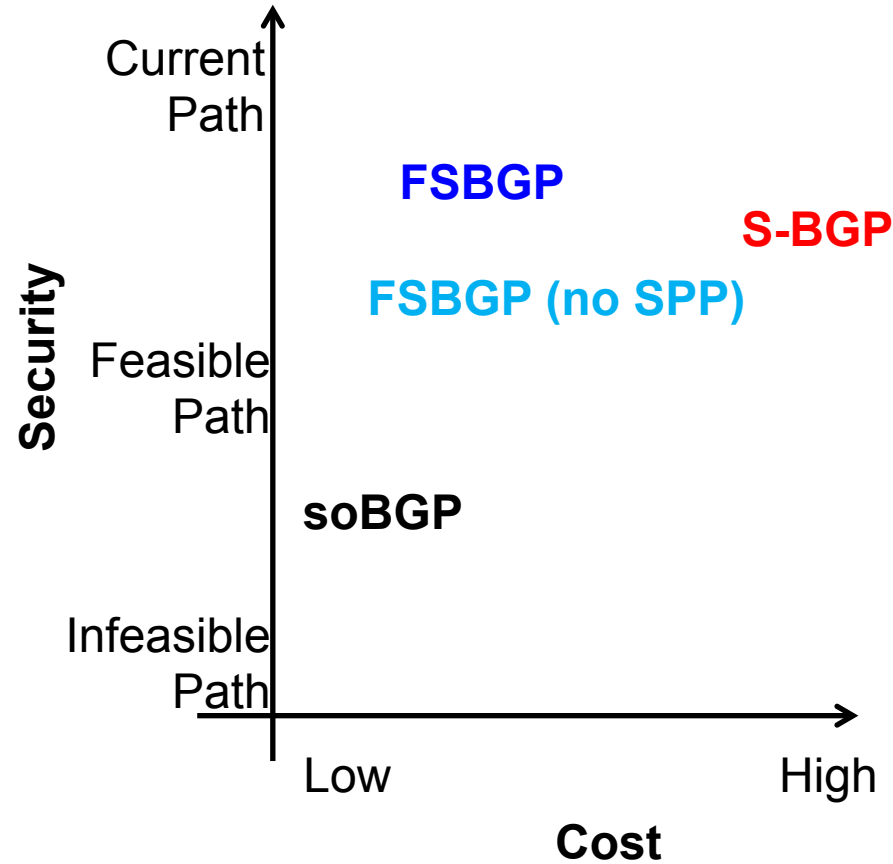
# Outline

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- Introduction
- FS-BGP: Fast Secure BGP
- **Evaluation**
  - **Security Level**
  - **Computational Cost**
- Discussion

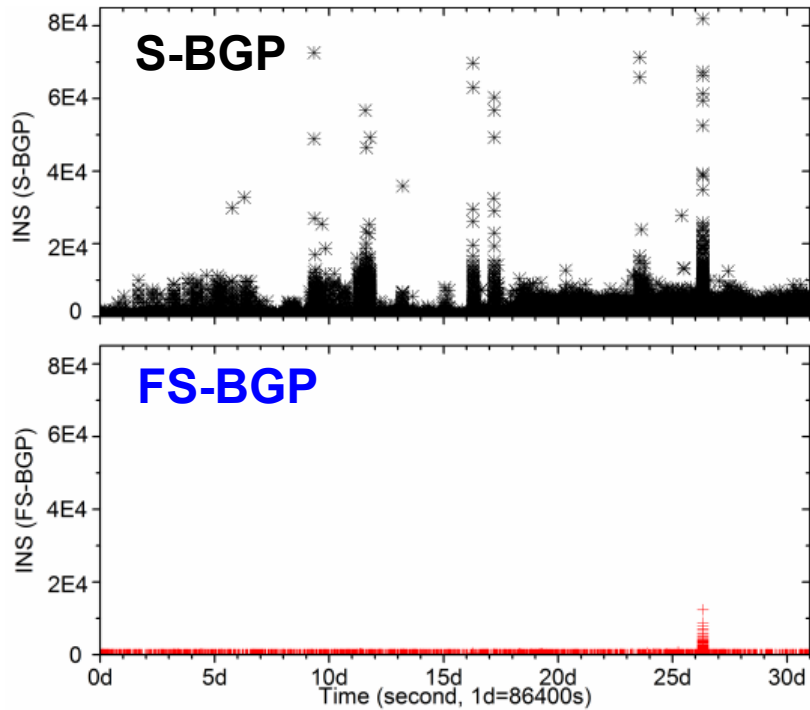
# Security Level

Type of Attack	FS-BGP	S-BGP	FS-BGP (no SPP)	so BGP
Inefficient hijack	✓	✓	✓	✓
False origin AS	✓	✓	✓	✓
Infeasible path	✓	✓	✓	✗
Feasible path				
Potential path	✓	✓	✓	✗
Revealed path	✓*	✓	✗	✗
Outdated path	✓*	✗	✗	✗
Policy violating [10]	✗	✗	✗	✗
Link-cut [4]	✗	✗	✗	✗

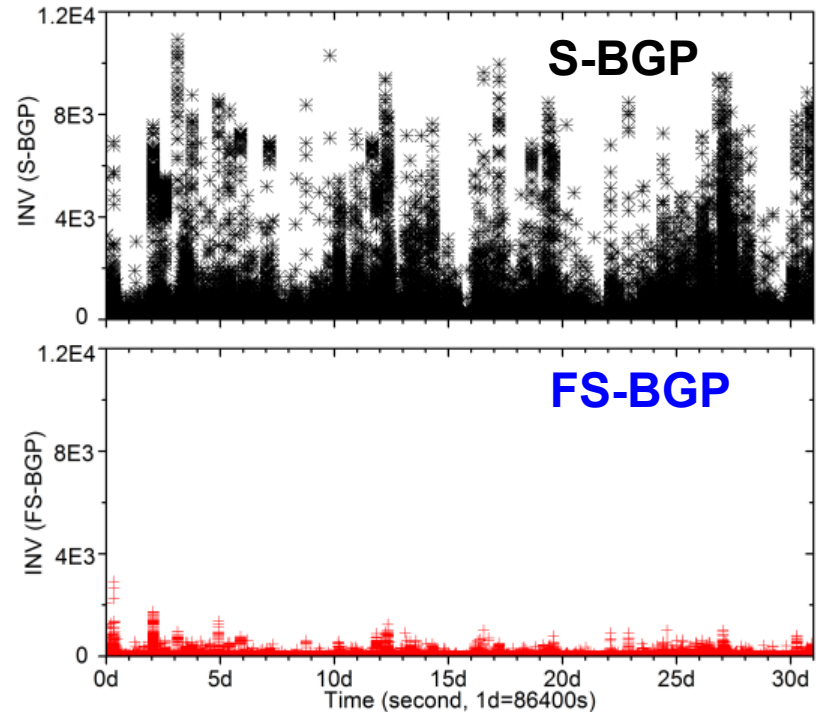


# Computational Cost

- 30 days' real BGP updates from backbone routers



# signings in every second



# verifications in every second

# Outline

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- Introduction
- FS-BGP: Fast Secure BGP
- Evaluation
- **Discussion**
  - **Support complicated routing policies**
  - **Protect privacy**

# Complicated Routing Policies

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- AS may use complicate route filters to describe their routing policies

– Prefix filter:

*export: to AS1 announce RS-ABC*  
*export: to AS2 announce 213.153.0.0/19*  
*export: to AS3 announce AS3^16-24*

← Included feasible prefixes into CSA

– Path filter:

*export: to AS4 announce <^AS4\$>*

← Sign whole path

– Origin filter:

*export: to AS5 announce AS-EFG*  
*export: to AS6 announce AS6*

← Included feasible origins into CSA

- FS-BGP can flexibly support route filters



# Revisit the route filters

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- Quantity of route filter
  - According our statistical result in IRR database, **only a very small portion** of policies use route filters
- Purpose of route filter
  - Some (i.e., origin/path filter) are set for **security considerations**, rather than policy requirements.
  - Others (i.e., prefix filter) are set for traffic engineering, to identifying the **preference of a route**, rather than the feasibility of a path

# Privacy Protection

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- Privacy: customer list ...
- FS-BGP can protect privacy data
  - Message spreading manner is same to BGP
  - Path segments not reveal additional info.
  - Path segments can only be passively received by valid BGP UPDATE receivers
  - Do NOT offer any kinds of public accessible policy database

# Next step: call for **WG** adoption

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- Acknowledgement
  - Greatly appreciate comments of *Russ White*
- Review
  - FS-BGP: Fast Secure BGP
    - CSA: Critical Segment Attestation
    - SPP: Suppressed Path Padding (Optional)
  - Evaluation
    - Security level: **similar security level as S-BGP**
    - Computational cost: **reduced the cost by orders of magnitude**
    - Support complicated routing policies
    - Protect privacy

**Thanks!**