

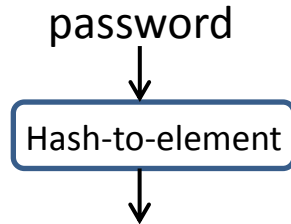
draft-harkins-tls-pwd

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- What
 - Certificate-less ciphersuites
 - Authentication using a password
 - Resistance to off-line dictionary attack
 - No, it's not patented
- Why (...not SRP)
 - EC support
 - Finite cyclic group is not fixed for each user, allowing a ciphersuite's hash and cipher to be of commensurate strength
 - Parlay a simple passcode into a certificate using a RFC 5967-style request and a RFC 5751-style (degenerate, certs only) response
 - Same key exchange used in other protocol for data plane protection, nice to do the same thing for control plane protection—straight forward way to provide consistent, system-wide security
 - Commodity-purchased smart energy device with limited GUI
 - Misuse-resistant TLS

How it Works (very broadly)

Alice generates Password Element



PE = password element

Alice generates 2 random numbers

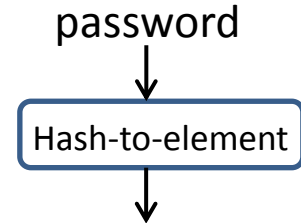
$\text{rnd-a, mask-a} \xleftarrow{\$} Z_q$

Alice sends scalar and element to Bob

$\text{scalar-a} = (\text{rnd-a} + \text{mask-a}) \bmod q \quad \text{-->}$

$\text{element-a} = \text{PE}^{-\text{mask-a}} \bmod p \quad \text{-->}$

Bob generates Password Element



PE = password element

Bob generates 2 random numbers

$\text{rnd-b, mask-b} \xleftarrow{\$} Z_q$

Bob sends scalar and element to Alice

$\text{scalar-b} = (\text{rnd-b} + \text{mask-b}) \bmod q \quad \text{<--}$

$\text{element-b} = \text{PE}^{-\text{mask-b}} \bmod p \quad \text{<--}$

Alice and Bob generate pre-master secret

$(\text{PE}^{\text{scalar-b}} * \text{element-b})^{\text{rnd-a}} \bmod p = \text{pre-master-secret} = (\text{PE}^{\text{scalar-a}} * \text{element-a})^{\text{rnd-b}} \bmod p$

How it works (changes to TLS)

```
enum { ff_pwd, ec_pwd } KeyExchangeAlgorithms;
```

```
struct {  
    opaque salt<1..2^8-1>;  
    opaque pwd_p<1..2^16-1>;  
    opaque pwd_g<1..2^16-1>;  
    opaque pwd_q<1..2^16-1>;  
    opaque ff_sscalar<1..2^16-1>;  
    opaque ff_selement<1..2^16-1>;  
} ServerFFPWPParams;
```

```
struct {  
    opaque salt<1..2^8-1>;  
    ECPParameters curve_params;  
    opaque ec_sscalar<1..2^8-1>;  
    ECPoint ec_selement;  
} ServerECPWPParams;
```

```
struct {  
    select (KeyExchangeAlgorithm) {  
        case ec_pwd:  
            ServerECPWPParams params;  
        case ff_pwd:  
            ServerFFPWPParams params;  
    };  
} ServerKeyExchange;
```

```
struct {  
    opaque ff_cscalar<1..2^16-1>;  
    opaque ff_celement<1..2^16-1>;  
} ClientFFPWPParams;
```

```
struct {  
    opaque ec_cscalar<1..2^8-1>;  
    ECPoint ec_celement;  
} ClientECPWPParams;
```

```
struct {  
    select (KeyExchangeAlgorithm) {  
        case ff_pwd:  
            ClientFFPWPParams;  
        case ec_pwd:  
            ClientECPWPParams;  
    } exchange_keys;  
} ClientKeyExchange;
```

- diff v00 v01
 - Salting password on server side
 - Mitigation of side channel attack on the process of hashing into an elliptic curve
 - Editorial changes: security considerations, justification/purpose

- OK, what do I want

- Ask WG to accept document and move it forward as a Proposed Standard

or, at the very least

- Stable, published specification
- Codepoints for pwd ciphersuites

```
CipherSuite TLS_FFCPWD_WITH_3DES_EDE_CBC_SHA = ( TBD, TBD );
```

```
CipherSuite TLS_FFCPWD_WITH_AES_128_CBC_SHA = (TBD, TBD );
```

```
CipherSuite TLS_ECCPWD_WITH_AES_128_CBC_SHA = (TBD, TBD );
```

```
CipherSuite TLS_ECCPWD_WITH_AES_128_GCM_SHA256 = (TBD, TBD );
```

```
CipherSuite TLS_ECCPWD_WITH_AES_256_GCM_SHA384 = (TBD, TBD );
```

```
CipherSuite TLS_FFCPWD_WITH_AES_128_CCM_SHA = (TBD, TBD );
```

```
CipherSuite TLS_ECCPWD_WITH_AES_128_CCM_SHA = (TBD, TBD );
```

```
CipherSuite TLS_ECCPWD_WITH_AES_128_CCM_SHA256 = (TBD, TBD );
```

```
CipherSuite TLS_ECCPWD_WITH_AES_256_CCM_SHA384 = (TBD, TBD );
```