

IETF83-KARP



Key Management and Adjacency Management for KARP-based Routing Systems

J. William Atwood
Revathi Bangalore Somanatha
Concordia University, Montreal

Contents



- ❑ Problem Statement
- ❑ Architecture and Definitions
- ❑ Overview
- ❑ Key Scopes
- ❑ Place of KARP Proposals
- ❑ Context Identifier assignment
- ❑ Adjacency
- ❑ Our Design
- ❑ Questions

Problem Statement



- ❑ Ongoing work on key management
 - RKMP – for unicast pairings
 - MRKMP – for multicast associations on a shared LAN
 - GDOI and GDOI-IKEv2 – examples of a group management protocol
- ❑ Ongoing work on adjacency management
 - None that we are aware of
 - We will present some ideas and hope for feedback from the WG members

Definitions



❑ Administrative Domain (AD)

- Set of routers under a single administration
 - RFC 4375 provides a convenient definition (in the context of Emergency Management)
- An AD is not bigger than an autonomous system
 - Because we are dealing with Interior Gateway Protocols

❑ Domain Controller (DC)

- Specific to a particular routing protocol (RP), because “adjacency” may be defined differently for each RP
 - Rules may be the same for different protocols, but stored data will be different

Definitions..2



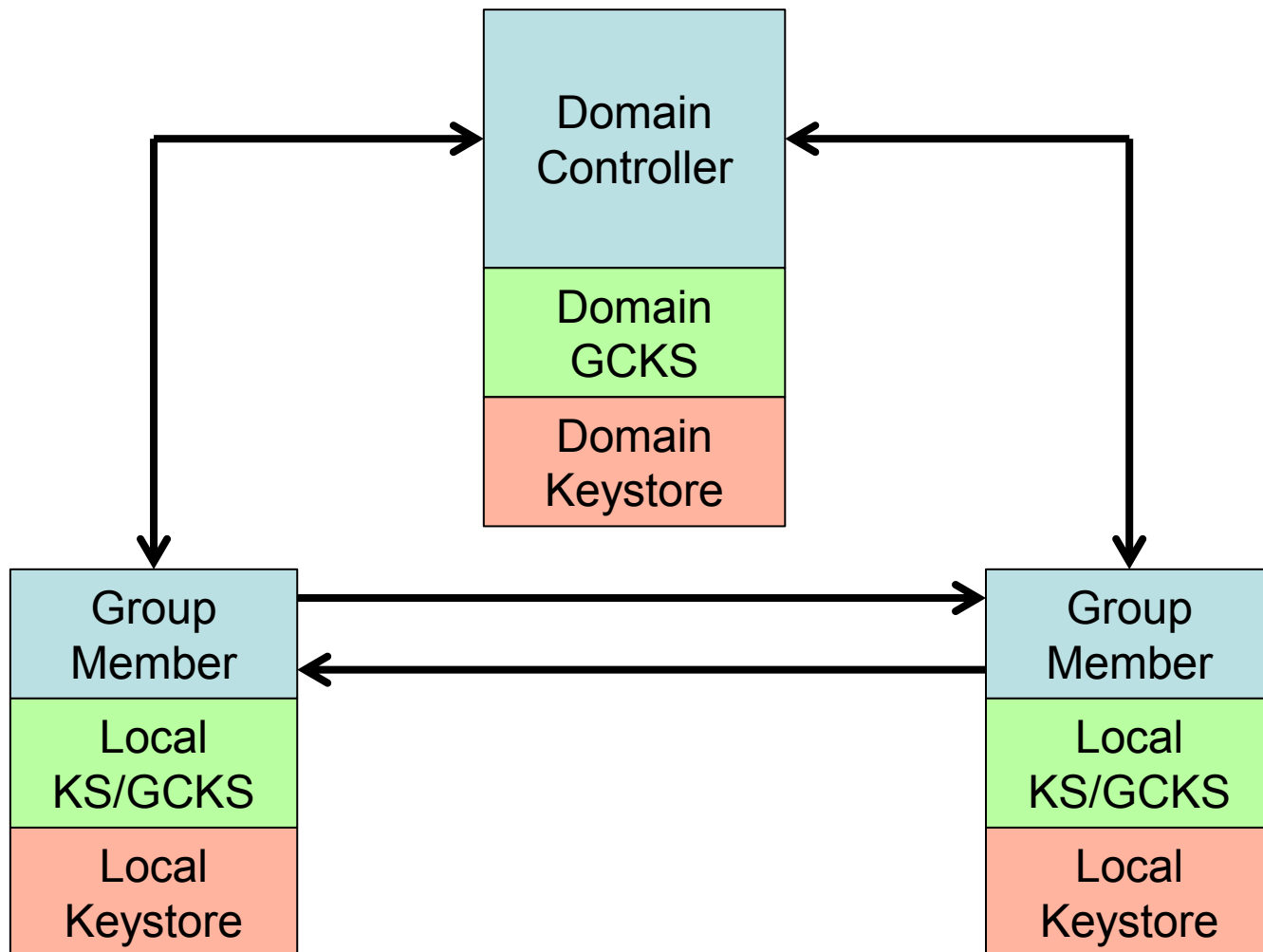
❑ Group Member (GM)

- Any router within the Administrative Domain
 - Note that depending on the keying model in use, we may form smaller “groups”

❑ Neighbor

- The set of routers that are adjacent to a particular router

Architecture



Overview



- ❑ Three issues for discussion
 - Key scope
 - Context Identifier assignment
 - Adjacency management

Overview..2



□ Key scope

- The subset of the GMs where a key is valid
- Two extreme examples
 - One key for whole region
 - Different keys for each interface for each sender

□ Context Identifier assignment

- MUST be centralized for multicast inter-router communication
 - SPI assignment for unicast IPsec contexts is receiver-based
 - SPI assignment in IPsec cannot be receiver-based when there are multiple receivers

Overview..2



- ❑ Adjacency control
 - If active, MUST be centrally managed
 - Otherwise, the router MAY use (insecure) neighbor discovery

- ❑ This implies that there must be a central (domain) controller
 - Our design tries to minimize the need to communicate with this central controller, especially when re-booting

- ❑ We are trying to prepare for adjacency control

Key scope: 4+1 cases



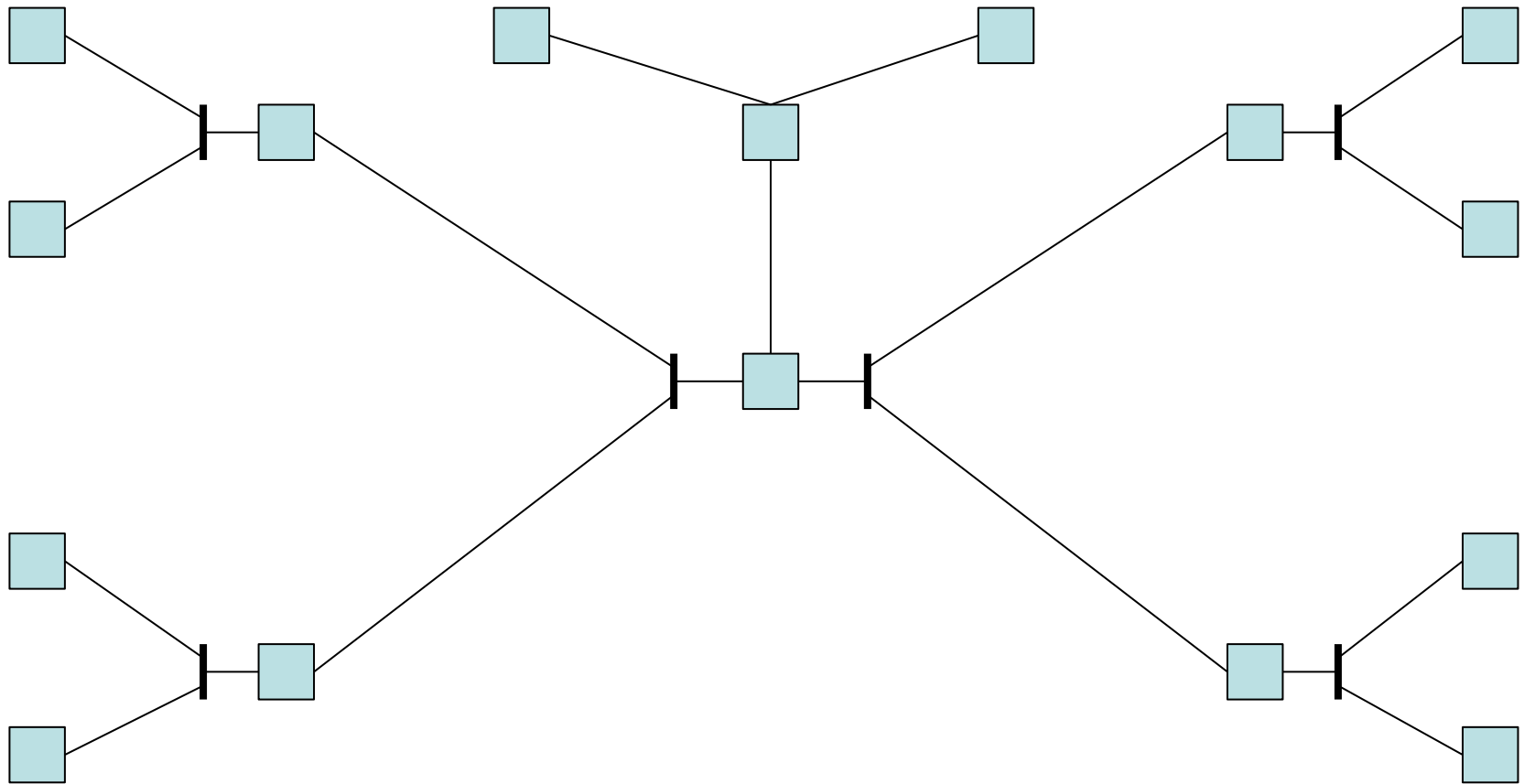
- ❑ One key for the AD
 - Very large attack surface
 - Key must be determined by the Domain Controller
- ❑ One key per shared LAN
 - Smaller attack surface
 - Key can be determined locally
 - By mutual agreement
 - By electing a local GCKS for that LAN

Key scope..2

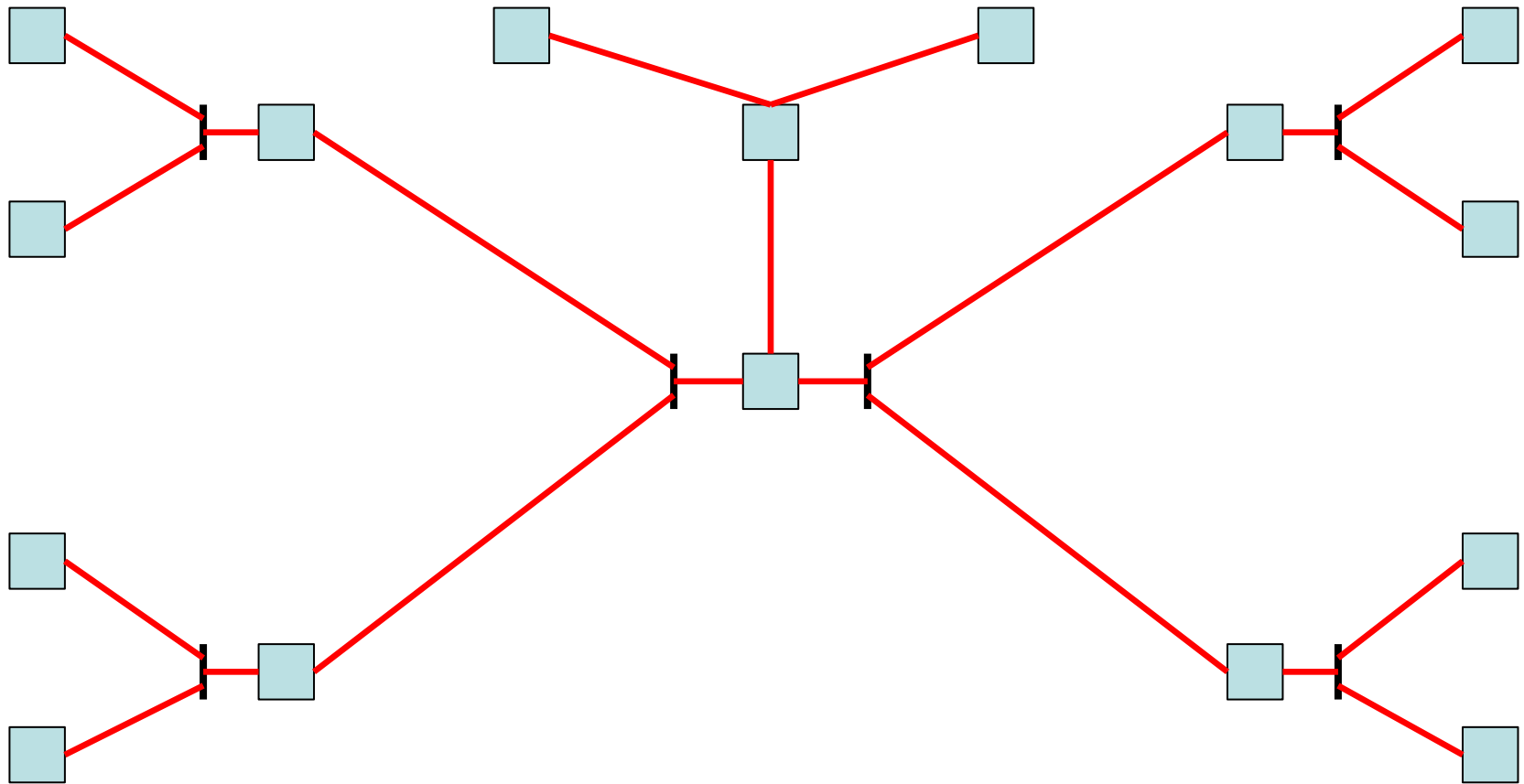


- ❑ One key per sending router
 - Even smaller attack surface
 - Key is determined by sending router, and distributed to its legitimate neighbors
- ❑ One key per interface per sending router
 - Smallest attack surface
 - Keys are determined by sending router
- ❑ Two keys per pair of routers
 - Unicast IPsec (IKE, IKEv2)
 - Application layer security (TLS)

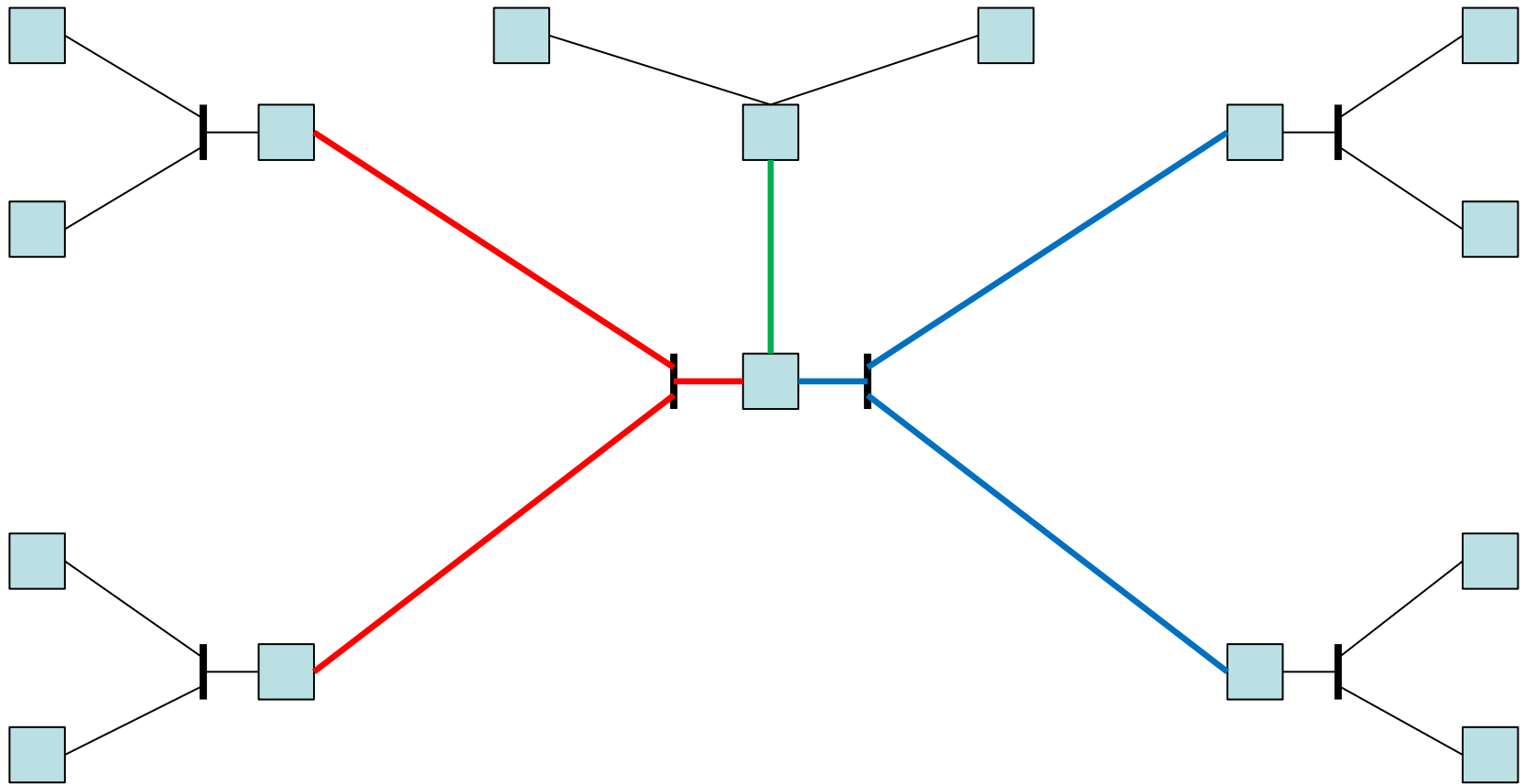
Example network



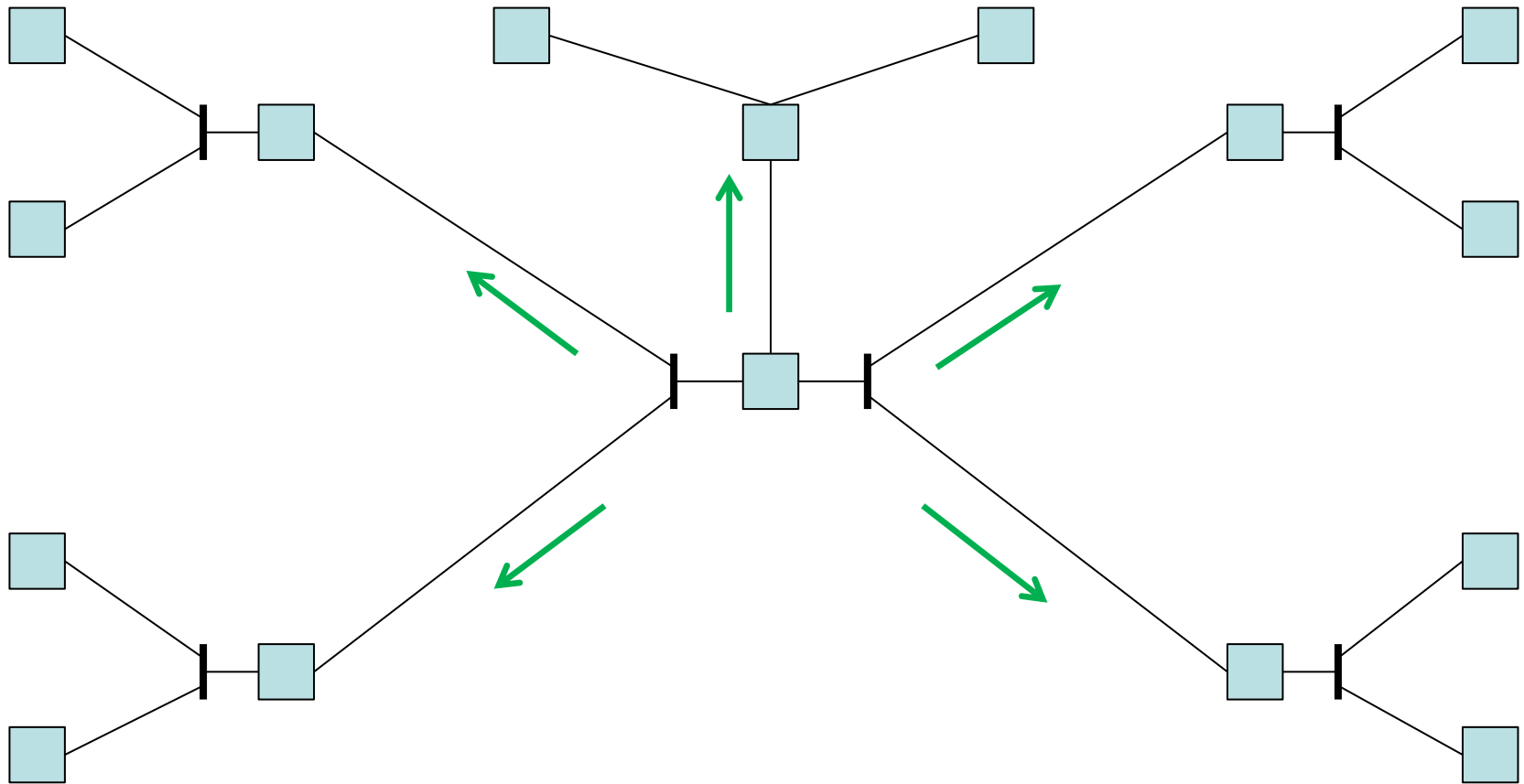
Single Key



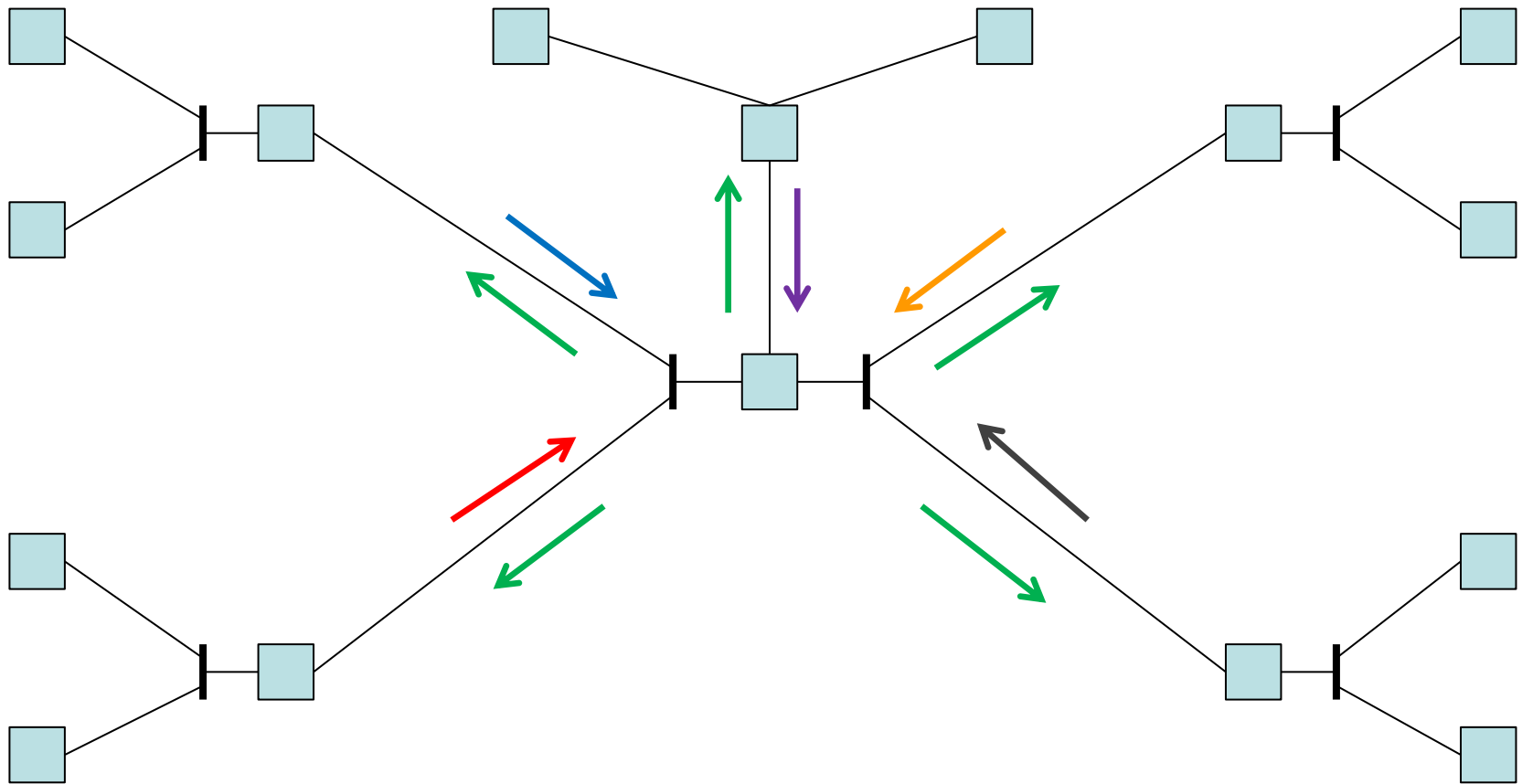
One Key per LAN



One Key per Sender: One Outgoing Key



One Key per Sender: Five Incoming Keys



Place of KARP proposals



❑ RKMP

- Used to establish peer-to-peer relationships
- Assumes a router identification method exists

❑ KMPRP

- Additional details of exchanges
- Deals with key rollover

Proposals..2



❑ MRKMP

- Focuses on the election of a local GCKS for the “One Key per LAN” model
- Assumes a router identification method exists
- Deals with router reboots
- Cannot deal with adjacency management

❑ GDOI/GDOI-IKEv2

- Does not take into consideration keying groups (key scopes)
- Does not deal with adjacency management

Context Identifier (CI) assignment



- ❑ One key
 - Context Identifier (e.g., SPI) to be used can be defined in the RFC, or by the administrator for the domain
- ❑ All other cases
 - Since there will potentially be multiple recipients of the group information, the CIs for each “mini-group” MUST be centrally assigned (i.e., by the Domain Controller)
 - There is probably a very nice graph-coloring problem inside this...

Adjacency



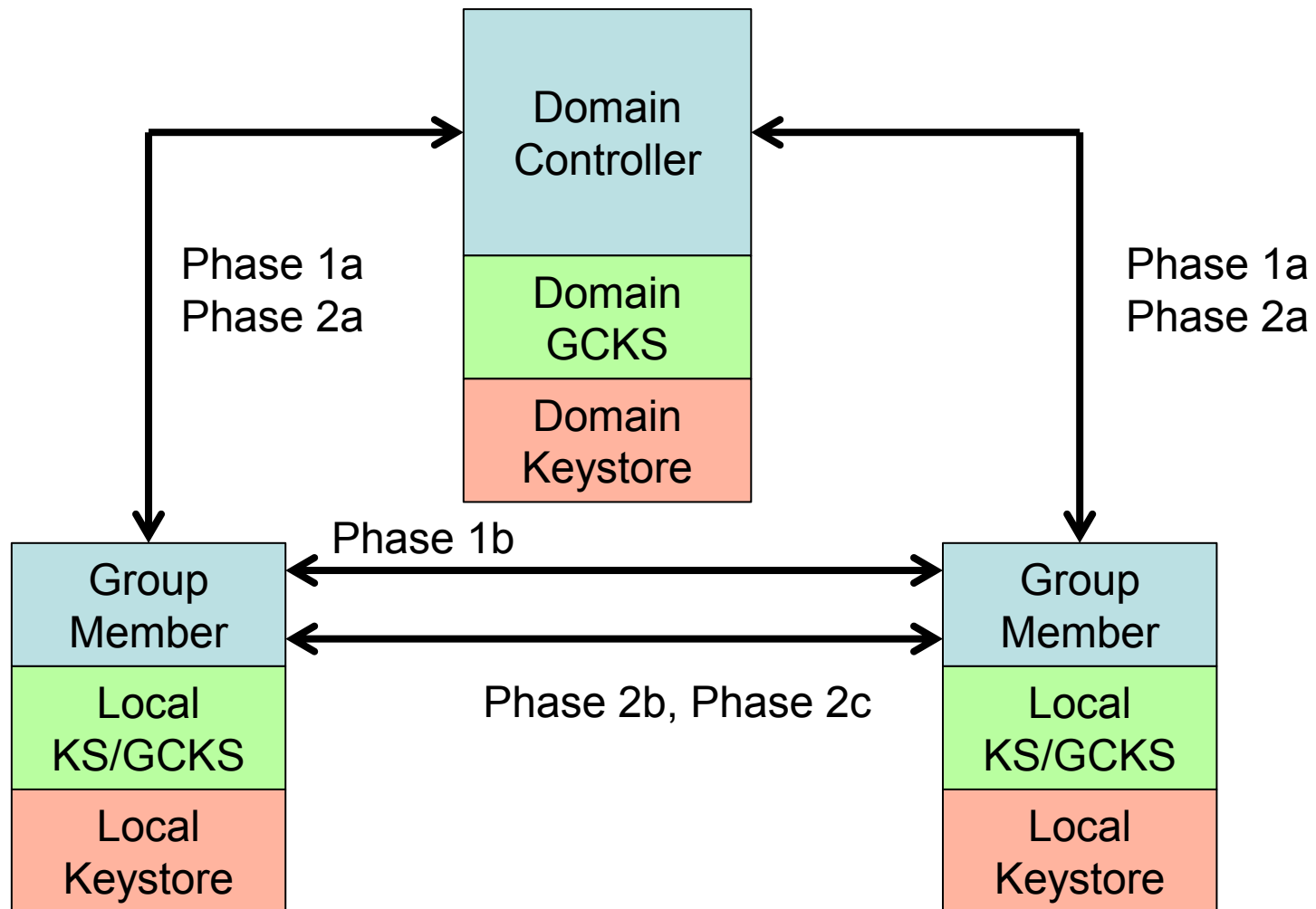
- ❑ Each router is assigned an “identity”
 - An FQDN, an arbitrary string, a PKI certificate, etc.
- ❑ Adjacency control can take a variety of forms
 - A neighbor is discovered, accept it
 - A neighbor has a valid certificate
 - (it is a valid router, but not necessarily adjacent to me)
 - A neighbor is permitted to be adjacent to me
- ❑ The last case **MUST** be centrally controlled
- ❑ The design must not prevent use of the other models (i.e., the disabling of adjacency control)

Our design



- ❑ We are exploring a design that
 - allows all of the above key scope models
 - allows us to control adjacency of routers
- ❑ Our intention is to specify the actors and the exchanges, and then formally validate the security of these exchanges using AVISPA

Key Management Phases: Between Components



Keying Phases: 1



□ Phase 1a

- Establish secure path and mutual authenticity between Domain Controller and individual Group Members
 - To be used to distribute information for use by the GM to identify and authenticate its neighbors

□ Phase 1b

- Establish secure path and mutual authenticity between adjacent Group Members
 - To be used to distribute parameters that will be used by the GM to send information to its neighbors (i.e., routing protocol control packets)

Phase 1 comments



- ❑ A single phase 1 MAY be used for all routing protocols on a particular router (for example, both OSPF and PIM), especially if their concept of “neighbor” is the same
- ❑ Phase 1a is the Phase 1 for IKE for the Domain Controller<->GM exchange
- ❑ Phase 1b is the Phase 1 for IKE for the GM<->GM exchange
 - It will happen only after the Phase 2a exchange occurs

Phase 1 comments..2



- We may need to find a good way of labeling the “keying group” that is being referenced:
 - How do I differentiate between the group on interface “x” and the group on interface “y”?
 - Is there a way to describe the interfaces that will be stable, and can be understandable to both the GM and the Domain Controller?

Keying Phases: 2



❑ Phase 2a

- Allows a GM to establish the identity of its neighbors (or be given the rules for establishing these identities)

❑ Phase 2b

- The GM contacts these identified neighbors
- Establishes their authenticity and legitimacy

❑ Phase 2c

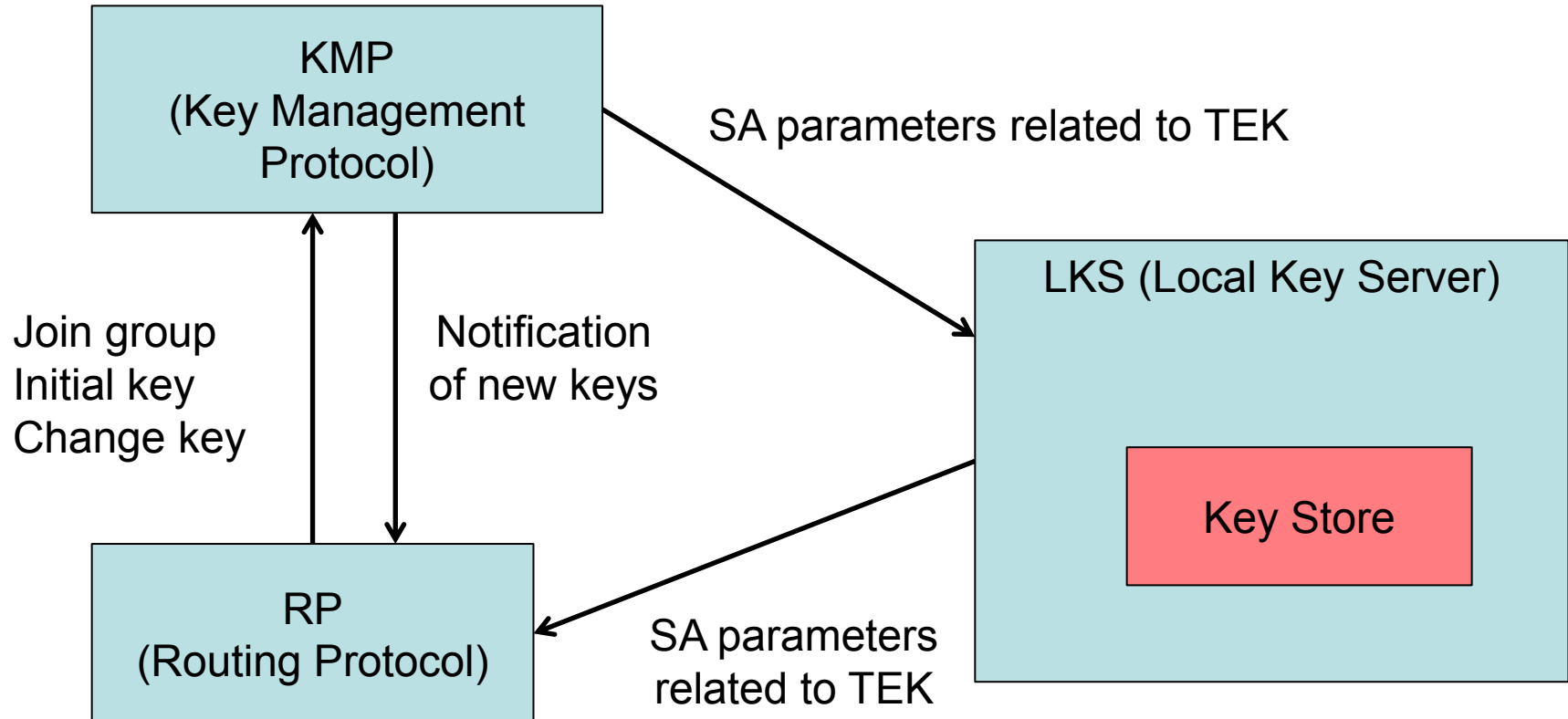
- The GM exchanges the information with its neighbors that will be used to send the routing protocol control packets (e.g., PIM-SM Hello)

Phase 2 comments



- ❑ If policy is transferred in Phase 2a, this should be done using standard policy-specification mechanisms
 - We are currently exploring the availability of such mechanisms within the IETF and elsewhere
- ❑ Depending on the rules provided in Phase 2a, parts or all of Phase 2b or Phase 2c may be suppressed

Key Management Exchanges: Within GMs



Our questions



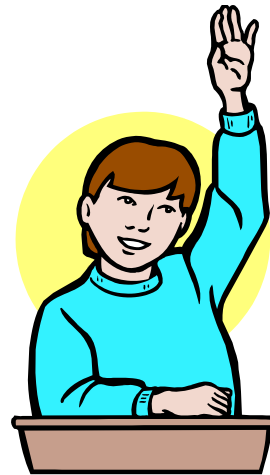
- ❑ Is this a reasonable model for the interactions that will occur?
- ❑ Are there things that we have left out that should be included?
- ❑ Any other comments?

Our plan



- ❑ There are some details of the interactions still to be worked out
- ❑ The modeling is in progress
- ❑ We expect to report on progress at IETF 84 in Vancouver

Thank You!



Questions?