Client-aided Congestion Management for TCP

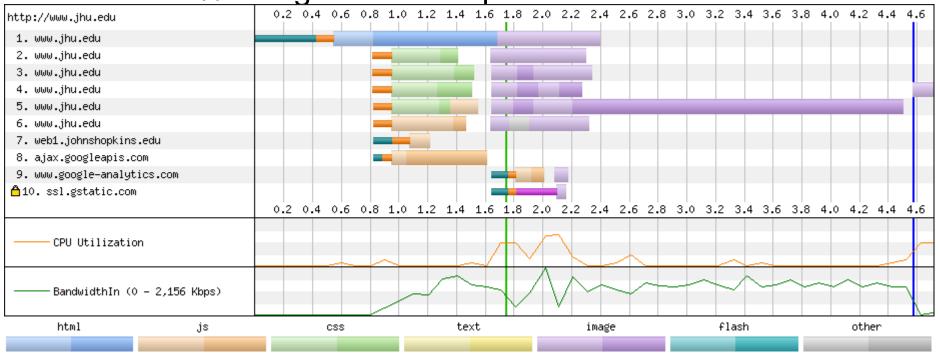
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Motivation: TCP throttles app performance

- Apps today use lots connections
 - Even with intelligent ADF multiplexing, e.g., SPDY
 - Persistent connection is common practice
- Every new connection has to (re)discover the network

~90% Google HTTP responses delivered in initial



A smart (or not-so-dumb) transport should ...

- Share network states among connections
 - Past and current active ones
 - Save or amortize reprobing time, e.g., slowstart
- A congestion manager (CM) on top of connections, *not* inside connections
 - New connection starts fast (as if it's never been disconnected)
 - Should recover fast and avoid timeout at all cost
 - N connections are as good as 1
 - Disincentivize parallel connections

Sender-side CM approach

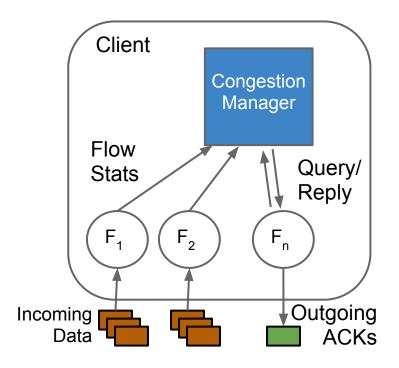
- Theory and practice
 - RFC 2140 TCP Ctrl Block Sharing by J. Touch '97
 - Congestion Manager by H. Balakrishnan, SIGCOMM '98
 - Ensemble-TCP by L. Eggert, CCR '00
 - SCTP, '00
 - Structured Stream Transport by B. Ford, SIGCOMM '07
 - Multi-path TCP
- Pros
 - Easy: sender traditionally holds all CC states
 - Fast deployment: maybe one side change only
- Cons
 - Scale: connections to same dst must hit same (physical) host
 - Difficult with large server farm load-balancing
 - Need big cache for the ever-growing Internet
 - Fragile: many devices/paths behind one client IP due to NAT

Can the client help?

- The client is the hub of *its* connections
 - Naturally the place for caching and sharing
 - Scales well
 - NAT is not a problem
- The client often knows better about the bottleneck: last-hop
 - Link properties: wired, wifi, or cellular
 - Link rate: edge vs 3G
 - Link failure and recovery
 - Dormant or active
- E.g., why RTO backoff then slow-start when a client can hint the sender the broken cellular link has recovered

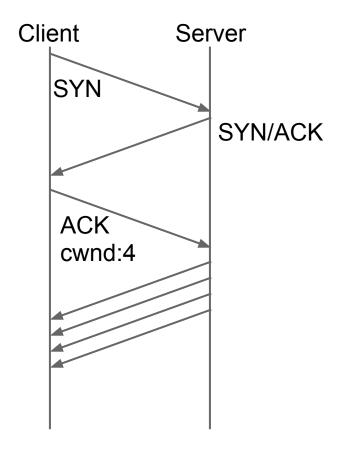
Great Snipe: a client-based congestion management

- A new TCP CC framework
 - Not a new congestion control algorithm
- Move congestion control to the client
 - Connections on the same path share one cwnd
 - also RTT, loss rate, reordering, etc
 - Network properties cached at the client
 - Use options for signalling
- Server handles e2e reliability
 - \circ $\,$ Detect and recover losses



Client-based congestion control

- Client as data receiver
 - Client maintains size of congestion window (cwnd)
 - Client passes cwnd to sender in ACKs
 - Sender limits # of outstanding packets to cwnd
- Benefit
 - Allows cwnd caching and reuse
- Client as data source
 - Same as TCP today



Implementing standard AIMD

- Connections on the same path share one cwnd (acwnd)
- On startup
 - Slow start with IW10 if no prior history
 - cwnd = acwnd / N otherwise
- On losses
 - Server performs traditional loss recovery and informs client
 - o acwnd = ssthresh
 - Reduce once across multiple losses or connections
 - acwnd = 1
 - If nothing received from the same dst for last RTO
- After cwnd reduction
 - acwnd += 1 per RTT
- Upon completion
 - acwnd remains same

Research issues / opportunities

- A pure client-based maybe overkill
 - What if client just guides the server somehow
- Sender announces the backlog to allow better acwnd allocation?
- Track one way delay (OWD)
 - New delay-based congestion control?
- Co-exist with traditional TCP and other protocols
 E.g., interactive or real-time protocols
- Detect shared bottlenecks among different paths
- Reusing / sharing other states, e.g., loss rates, reorderig, etc
- Energy efficiency

Conclusion

- Congestion control should be on top of individual logic connections
- Server-based congestion manager has practical scale issue
 - Client may offer interesting opportunities to improve CC today!
 - Often knows the network better
 - Naturally the sharing point
 - Scale well
- Great Snipe: move CC to client and on top of indiv. connections
 - Still in early development stage
 - Will release to the public for testing like Laminar
- Feedback & ideas welcome! ycheng@google.com