#### ALTO Use Case: Resource Orchestration for Multi-Domain, Geo-Distributed Data Analytics

draft-xiang-alto-multidomain-analytics-02

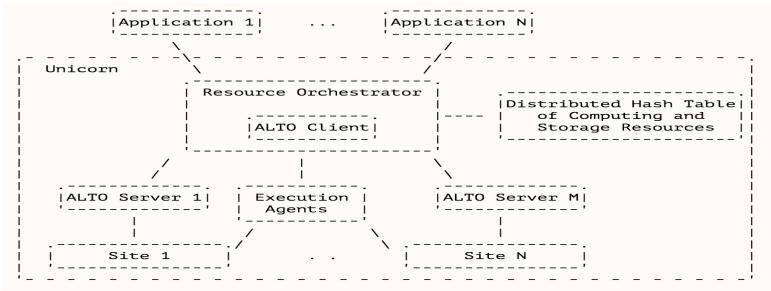
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## Takeaway from IETF 101

• Substantial updates for document review:



Three-phase resource discovery

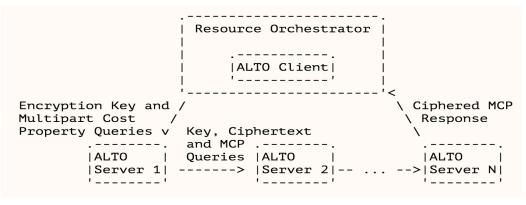
- 1. ALTO EPS to discover locations and properties of computing and storage resources;
- 2. ALTO ECS to discover the connectivity between computing and storage resources
- 3. ALTO PV extension to discover the networking resource sharing between flows.
  - Propose an ALTO extension to support accurate, privacy-preserving resource discovery across multiple domains.

## Update for IETF 102

- Two technical updates for the resource abstraction discovery phase (Phase 3).
  - Update the design of the privacy-preserving multi-domain resource abstraction aggregation protocol.
    - The new design does not require a chaining aggregation process between different ASes.
  - Introduce a super-set projection technique to improve the scalability.

#### Phase 3: Resource Abstraction Discovery

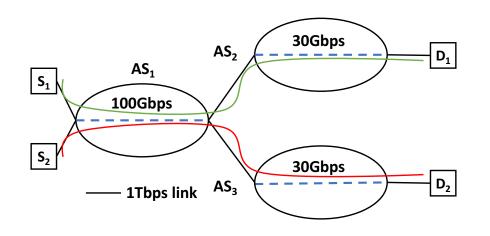
 Previous design: an ALTO-extension for privacy-preserving multidomain resource information aggregation, which returns the intersected capacity region of all networks in ALTO PV extension.



# Representation of capacity region after the aggregation:

 $\begin{array}{l} 69x_1+61x_2+11x_{11}^s+58x_{21}^s+50x_{31}^s=4340,\\ 71x_1+118x_2+49x_{11}^s+22x_{21}^s+69x_{31}^s=7630,\\ 170x_1+184x_2+95x_{11}^s+75x_{21}^s+89x_{31}^s=14420,\\ 59x_1+129x_2+34x_{11}^s+25x_{21}^s+95x_{31}^s=7000, \end{array}$ 

• Example:



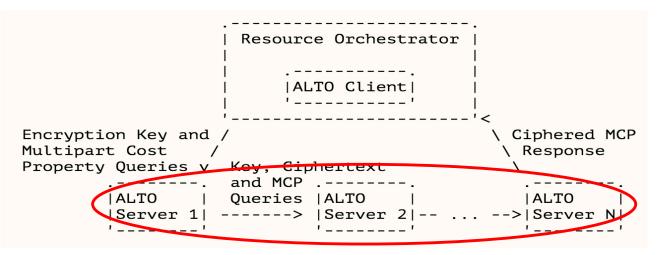


 $\begin{aligned} &\Pi_1(F_1): \{x_1+x_2 \leq 100\} \\ &\Pi_2(F_2): \{x_1 \leq 30\} \\ &\Pi_3(F_3): \{x_2 \leq 30\}. \end{aligned}$ 

Representation of capacity region before the aggregation: 4

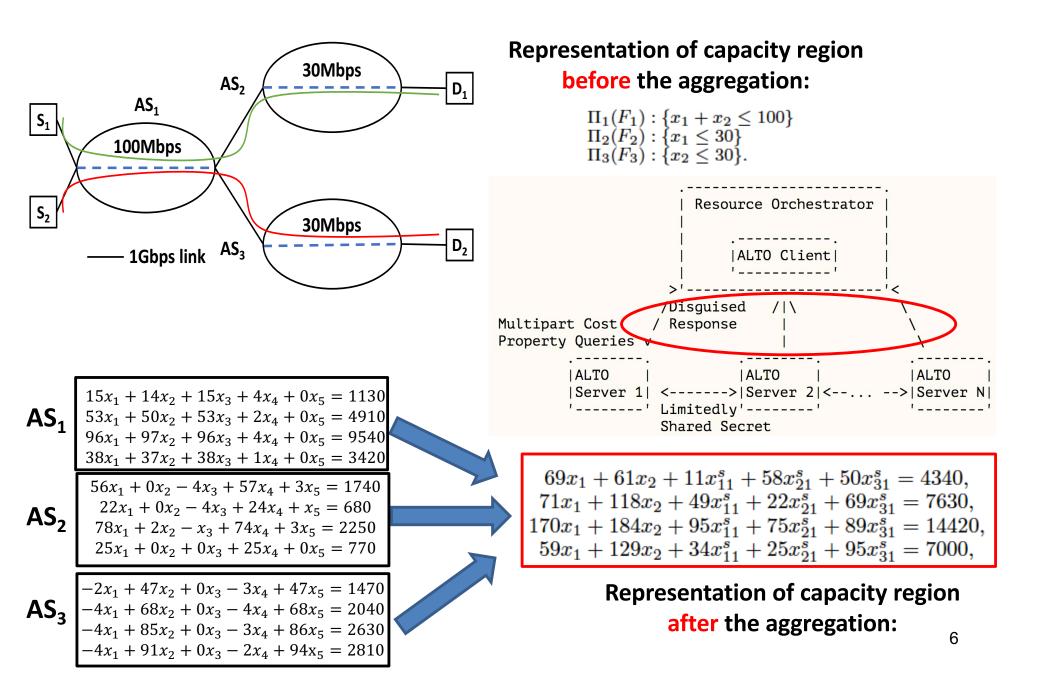
#### Phase 3: Resource Abstraction Discovery

 Issue: The aggregation process is a chaining process across all ALTO servers, which will take a long time for the ALTO client to get the final result.

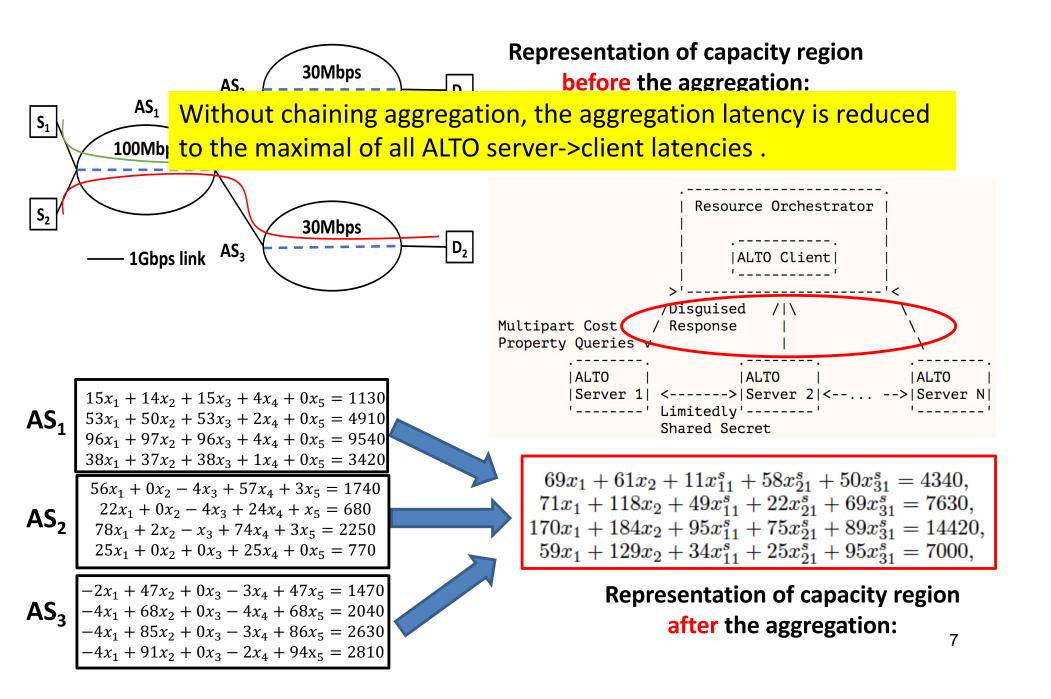


- In the -02 draft, we design a new privacy-preserving aggregation service, which does not require the chaining aggregation process.
- Basic idea: Each ALTO server disguises its own set of linear inequalities in the ALTO-PV response with an obfuscating algorithm we developed.

#### New Privacy-Preserving Aggregation Service: An Example



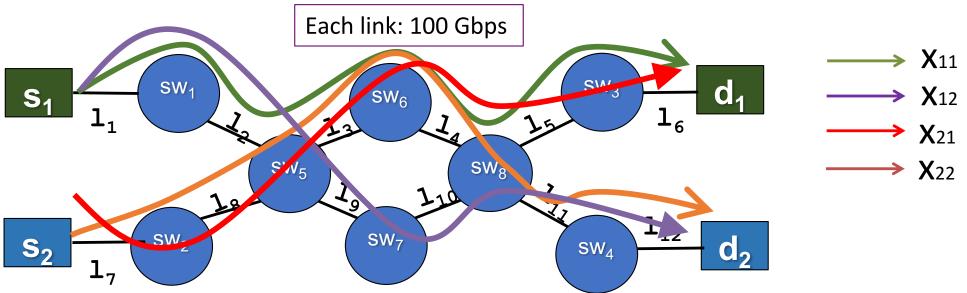
#### New Privacy-Preserving Aggregation Service: An Example



#### Improve Scalability: Super-Set Projection

- Issue: In collaborative science experiments, the number of data analytics jobs is huge. Repeatedly querying ALTO servers and making ALTO servers compute the responses for every new job would raise scalability issue.
- Current design proposal: Super-set projection.
- **Basic idea**: For each ALTO server in each network, let it precompute the routing information, and the resource sharing information for a set of flows, whose source and destination is the combination of all ingresses and egresses of the networks.
  - E.g., Assume a network with M ingress and N egress, precompute the route and bandwidth sharing for a set of M\*N flows.
  - When a new PV query comes in, the ALTO server can projects the precomputed set of linear inequalities for the M\*N flows based on the ingresses and egresses of the flows in the PV query, to get the resource sharing information for this query.

#### Super-Set Projection: Example



- Only two ingresses  $(I_1, I_7)$  and two egresses  $(I_6, I_{12})$
- Pre-computed set of linear inequalities:  $\begin{array}{l} x_{11} + x_{12} \leq 100, \, \text{for link} \, \{l_1, l_2\} \\ x_{11} + x_{21} \leq 100, \, \text{for link} \, \{l_5, l_6\} \\ x_{11} + x_{21} + x_{22} \leq 100, \, \text{for link} \, \{l_3, l_4\} \\ x_{21} + x_{22} \leq 100, \, \text{for link} \, \{l_7, l_8\} \\ x_{12} + x_{22} \leq 100, \, \text{for link} \, \{l_{11}, l_{12}\} \\ x_{12} \leq 100, \, \text{for link} \, \{l_9, l_{10}\} \end{array}$
- When a PV query comes with two flows (s1, d1) and (s2, d2), the projected result is:  $x_{11} \leq 100$ , for link  $\{l_1, l_2, l_5, l_6\}$

 $x_{11} \leq 100, \text{ for link } \{l_1, l_2, l_5, l_6\}$   $x_{11} + x_{22} \leq 100, \text{ for link } \{l_3, l_4\}$  $x_{22} \leq 100, \text{ for link } \{l_7, l_8, l_{11}, l_{12}\}$ 

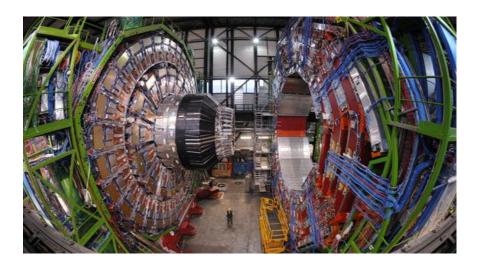
#### Next Steps

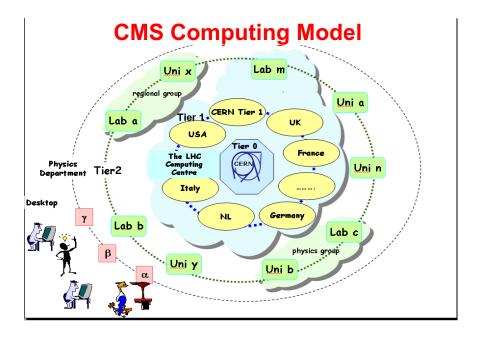
- Privacy-preserving information aggregation:
  - This is a first proposal to address the security/privacy concerns of using ALTO.
  - Interests in moving it to a formal extension (standard track)?
- Super-set projection
  - The current design focuses on resource abstraction discovery (phase 3).
  - Sending the pre-computed set of linear inequalities to the ALTO client, who can do projection by itself, could further reduce the discovery latency.
  - How to extend this design to the other two phases (endpoint and path discoveries) without raising additional privacy concerns?
- The overall system is under the final review phase of SuperComputing'18.

# Backup slides

#### Recap: Multi-Domain, Geo-Distributed Data Analytics

- Settings: Different organizations contribute various resources (e.g., sensing, computation, storage and networking resources) to collaboratively collect, share and analyze extremely large amounts of data.
  - Example: the CMS experiment in Large Hardon Collider.

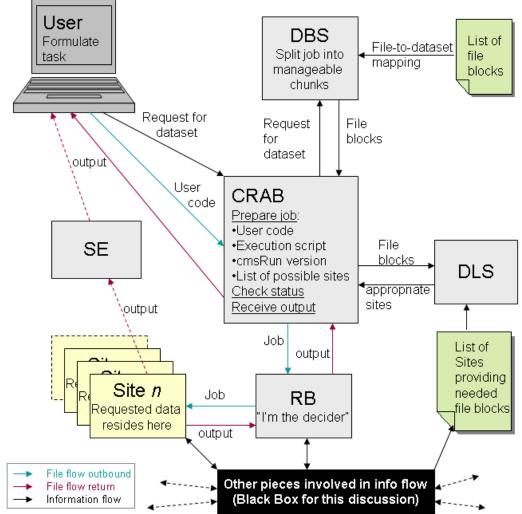




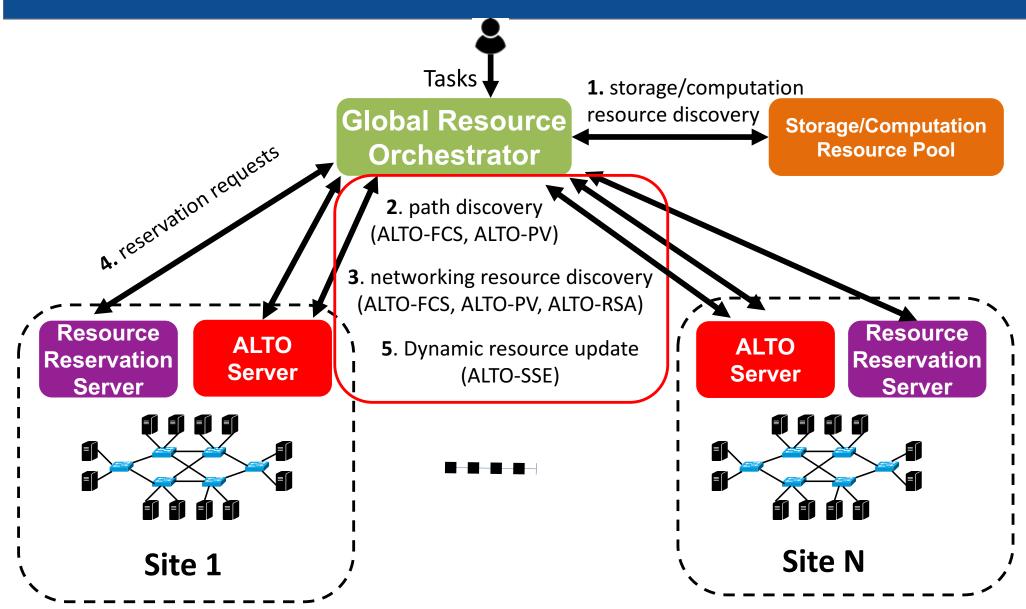
#### Current CMS Data Analytics Work Flow

# Factors determining data analytics task delay.

- Task decomposition (parallelization).
- Data transmission from input dataset location to computation nodes.
- Data transmission from computation nodes to output dataset sites.
- Current CMS workflow.
  - Simple, manual parallelization.
  - Opportunistic, network-unaware computation node assignment.
  - Opportunistic, network-unaware output stage out.

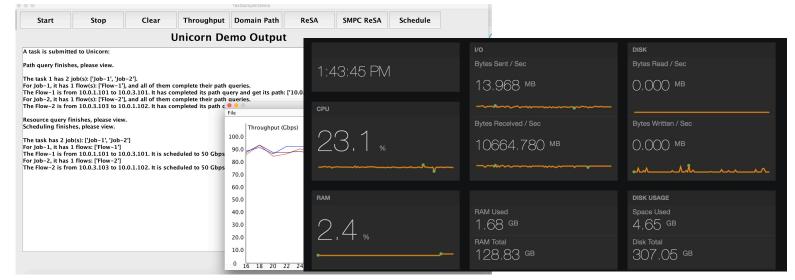


#### Architecture



#### **Unicorn Implementation and Demonstration**

- Orchestrator: ~2700 LoC Python code
- ALTO server: ~3000 LoC Java code
- Resource reservation server:
  - fast data transfer (FDT), FireQoS, OpenvSwitch, etc.
- Network controllers: OpenDaylight, Kytos
  - ONOS and Ryu are under development
- Demonstrated on different topologies at SuperComputing 2017 [2].



[2] Xiang, *et.al.*, "Unicorn: Unified Resource Orchestration for Multi-Domain, Geo-Distributed Data Analytics", 15 in INDIS Workshop 2017.