

MLS Message Protection

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What's in the draft today...

Initial draft for Message Protection.

- A. Defines an Application Key Schedule (AKS)
to go from the Epoch Secret to the Application messages encryption keys.

- B. Defines which algorithms to use and on which objects
to protect Application messages against active network attackers,
with improved resistance to traffic analysis using optional padding.

Handshake Key Schedule

```
init_secret_[n-1] (or 0)
  |
  v
update_secret -> HKDF-Extract = epoch_secret
  |
  +--> Derive-Secret(., "app", GroupState_[n])
  |     = application_secret
  |
  v
Derive-Secret(., "init", GroupState_[n])
  |
  v
init_secret_[n]
```

Application Key Schedule

Two main ways of chaining secrets: interleaving or in parallel independently from what we chain (Application secret or message encryption key)

- Group chaining of the secret (interleaving)
- Participant chaining of the secret (parallel)

Both would need anyway need to provide security properties such as:

- Forward Secrecy for secrets / messages
- PCS for secrets / messages

Application Key Schedule

Group chaining of the secret (interleaving)

- + Reduced complexity (storage)
- + Improved Forward Secrecy (no unused key stored for a long time)
- Reduced ability to handle out-of-order messaging for high frequency transmissions

Participant/Sender chaining of the secret (parallel)

- + Well-known design
- + Able to handle out-of-order messaging
- Higher complexity (storage)

Handshake Key Schedule

```
init_secret_[n-1] (or 0)
  |
  v
update_secret -> HKDF-Extract = epoch_secret
  |
  +--> Derive-Secret(., "app", GroupState_[n])
  |     = application_secret
  |
  v
Derive-Secret(., "init", GroupState_[n])
  |
  v
init_secret_[n]
```

Application Key Schedule

Derive a Participant Application Secret for each potential sender.

```
application_secret
  |
  v
Derive-Secret(., "app sender", [sender])
  |
  v
application_secret_[sender]_[0]
```

Application Key Schedule

Move forward the Participant Application Secret for each message.

```
application_secret_[sender]_[N-1]
  |
  +--> HKDF-Expand-Label(., "nonce", "", nonce_length)
  |   = write_nonce_[sender]_[N-1]
  |
  +--> HKDF-Expand-Label(., "key", "", key_length)
  |   = write_key_[sender]_[N-1]
  V
Derive-Secret(., "app upd", "")
  |
  V
application_secret_[sender]_[N]
```


Application Message Protection

Expected Secrecy Properties

- + Forward Secrecy for each message inside a Participant chain
- + Forward Secrecy for each message across Group Updates
- + Post-Compromise Security for messages across Group Updates

Expected Authentication Properties

- + Weak Authentication from Group membership
- + Strong Authentication from Signatures (currently)

Application Message Protection

Signatures inside the ciphertext for Privacy reasons

```
struct {
    uint8  group[32];
    uint32 epoch;
    uint32 generation;
    uint32 sender;
    opaque content<0..2^32-1>;
} MLSSignatureContent;
```

Application Message Protection

Encryption via standard AEAD constructions and optional padding

```
struct {
    opaque content<0..2^32-1>;
    opaque signature<0..2^16-1>;
    uint8 zeros[length_of_padding];
} ApplicationPlaintext;

struct {
    uint8 group[32];
    uint32 epoch;
    uint32 generation;
    uint32 sender;
    opaque encrypted_content<0..2^32-1>;
} Application;
```

Summary

We have a reasonable initial draft for message protection

- Can we have PCS for all messages ? Do we want it ?
- We need some data to evaluate the computation and storage performance for the per-participant chaining and eventually the group chaining
- Should we keep signatures or use different constructions ?

We will need to have a look at...

- AES 128 vs AES 256
- AES GCM vs AES SIV
- PQ Primitives sizes