

Integrating DTN and MANET Routing

2006-11-09

DTNRG – IETF 67 (slides from ACM CHANTS workshop)

Jörg Ott Dirk Kutscher Christoph Dwertmann jo@netlab.tkk.fi dku@tzi.org nermal@tzi.org

© 2006 Jörg Ott · Dirk Kutscher



TZi Center for Computing Technologies

Motivation: Ad-hoc Networking for Human Users

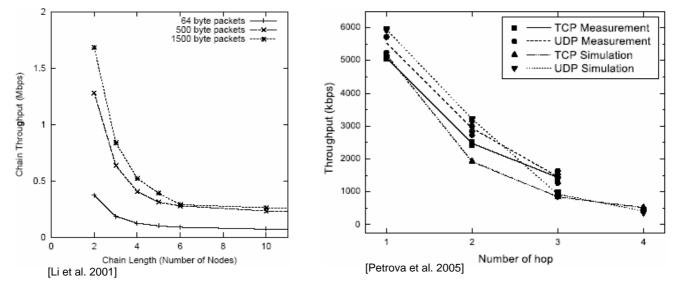
- Target field: interpersonal communication and Internet access
- Node density in environments without controlled deployment will often be sparse
 - Heterogeneity of consumer devices: different brands, versions, capabilities
 - Also: users may not want to cooperate
 - An end-to-end path may exist—but often will not
 - Calls for asynchronous communications: DTN
- User experience benefits from immediate end-to-end interactions
 - Users have grown expectations over time
 - Application protocols designed in this way anyway
 - Anything but end-to-end may only be second best
 - Yet better than not communicating at all





Motivation: Limitations in End-to-End Performance

 Communication performance degrades with increasing number of wireless hops in ad-hoc networks



Path stability decreases with distance depending on mobility

© 2006 Jörg Ott · Dirk Kutscher



TZi Center for Computing Technologies

3

Approach: Integrating DTN and AODV

- Let the application (and ultimately the user) decide
- Provide an interface to learn about the ad-hoc environment
 - Rather than completely abstracting from (i.e., hiding) it
- Combine AODV route search with DTN node location
 - Return available end-to-end route + path length
 - Return available DTN nodes in the vicinity + path lengths
- Application then chooses whether to use synchronous end-to-end or asynchronous hop-by-hop (DTN) communication
 - Example: Simply use DTN as fallback if end-to-end fails
 - Example: Path length, expected performance and stability vs. data volume



b-end

5

- Let the application (and ultimately the user) decide
- Provide an interface to learn about the ad-hoc environment.
 - R Applicability Statement: No, this does not work with existing applications
- Com
 - R<mark>But:</mark>

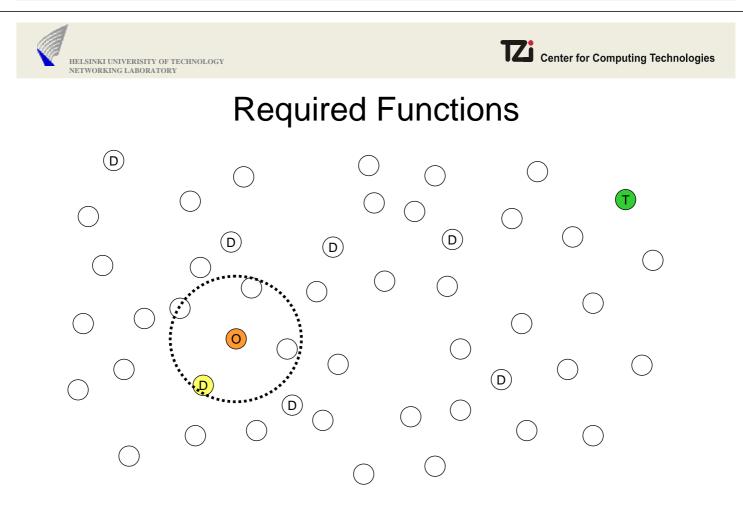
HELSINKI UNIVERISITY OF TECHNOLOGY NETWORKING LABORATORY

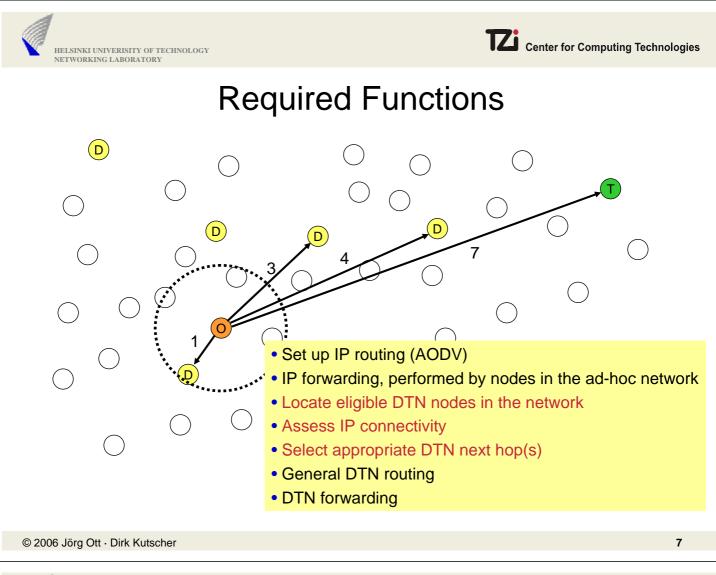
- Existing applications don't work well with mobility anyway.
- Appl

Or as Short-term: use proxies as interim means for migration.

- Long-term: design application protocols in a more suitable way.
- Example: Deth length, expected performance and stability value data valume Special Application: DTN routers

© 2006 Jörg Ott - Dirk Kutscher







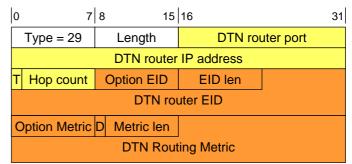
HELSINKI UNIVERISITY OF TECHNOLOGY NETWORKING LABORATORY **TZi** Center for Computing Technologies

AODV Extensions

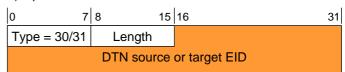
Additional attributes in RREQ and RREP packets

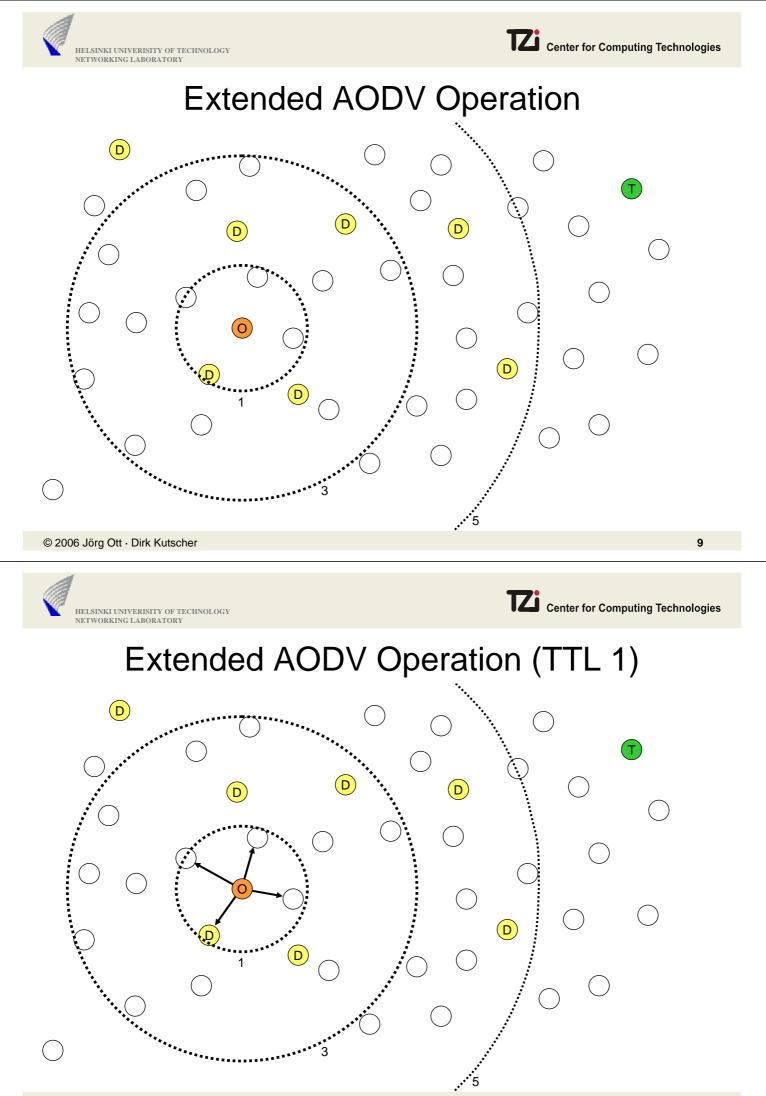
- Indicate extension support
- Report DTN node contact info
- Optional
 - Convey source, target EID
 - Include routing metrics
- Modified processing and timer handling rules
- Workable with plain AODV nodes

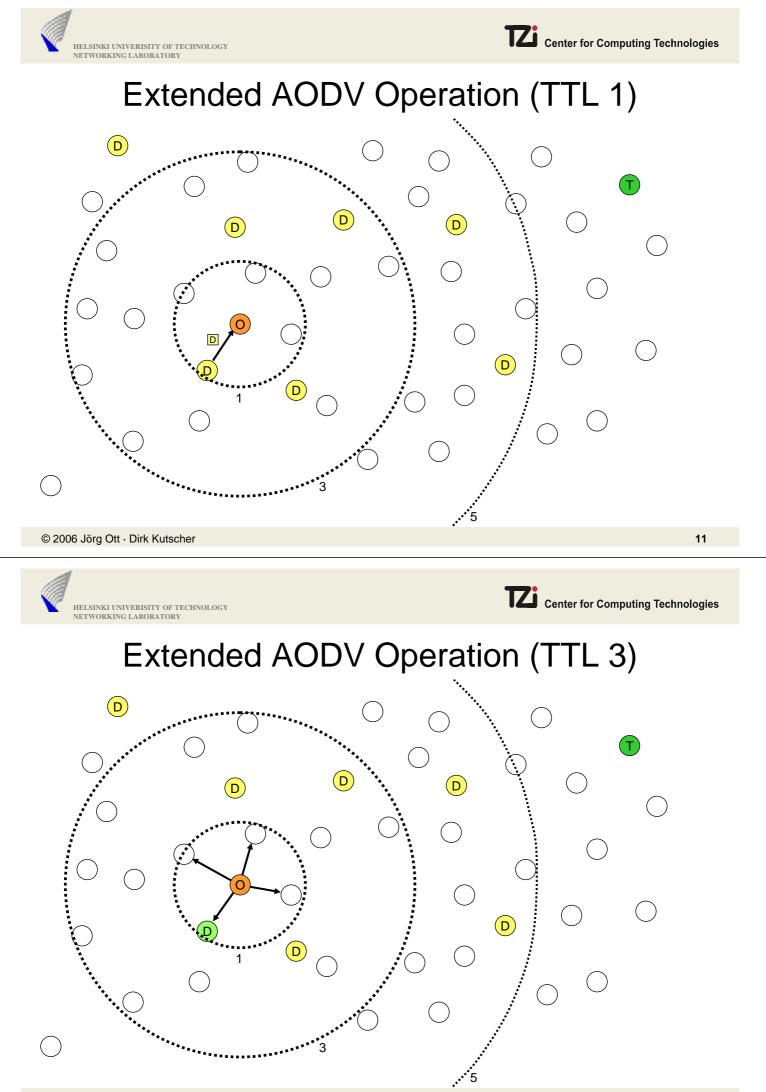
a) DTN router info + optional metric

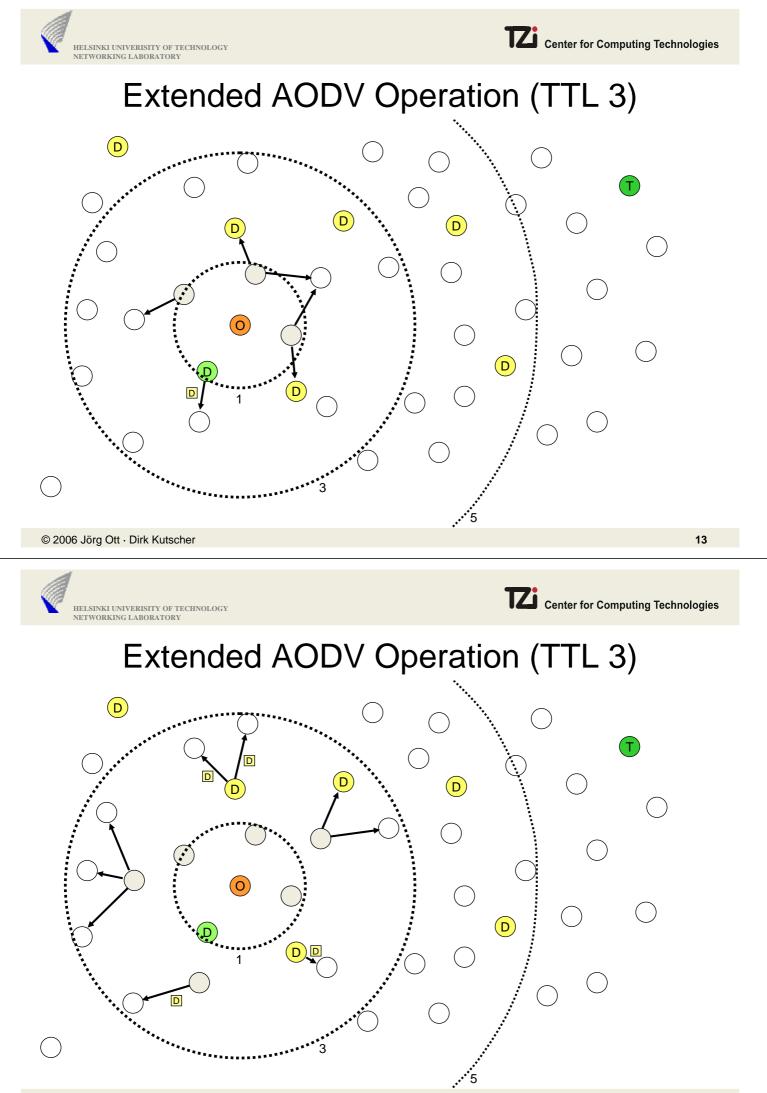


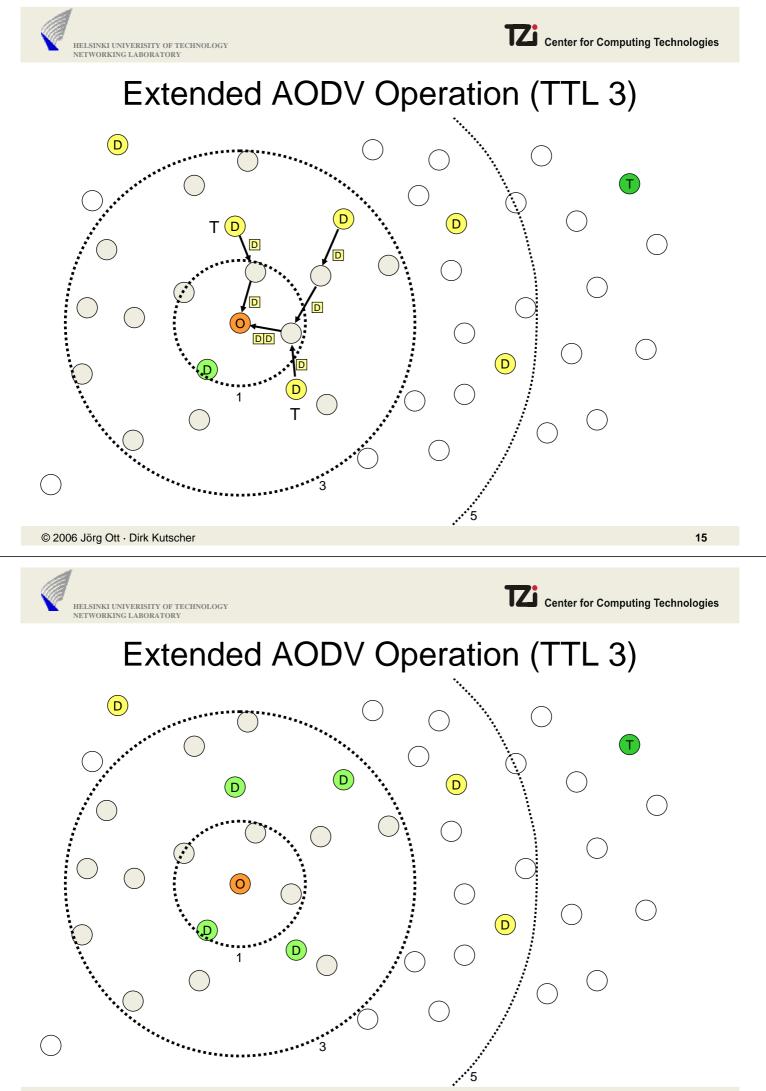
b) Optional DTN EID

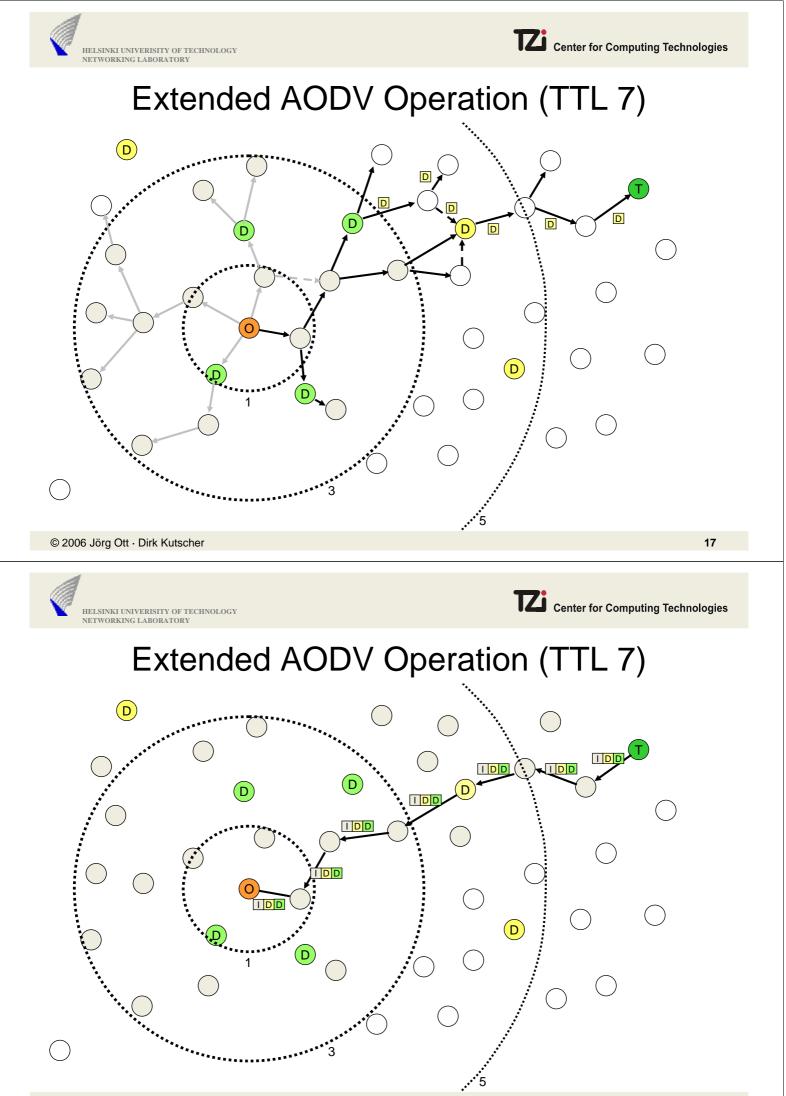


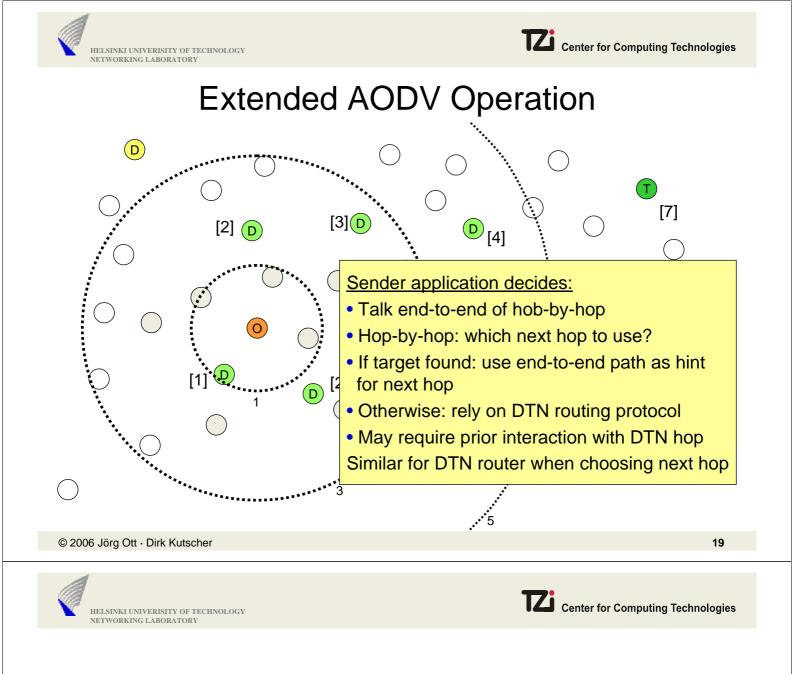








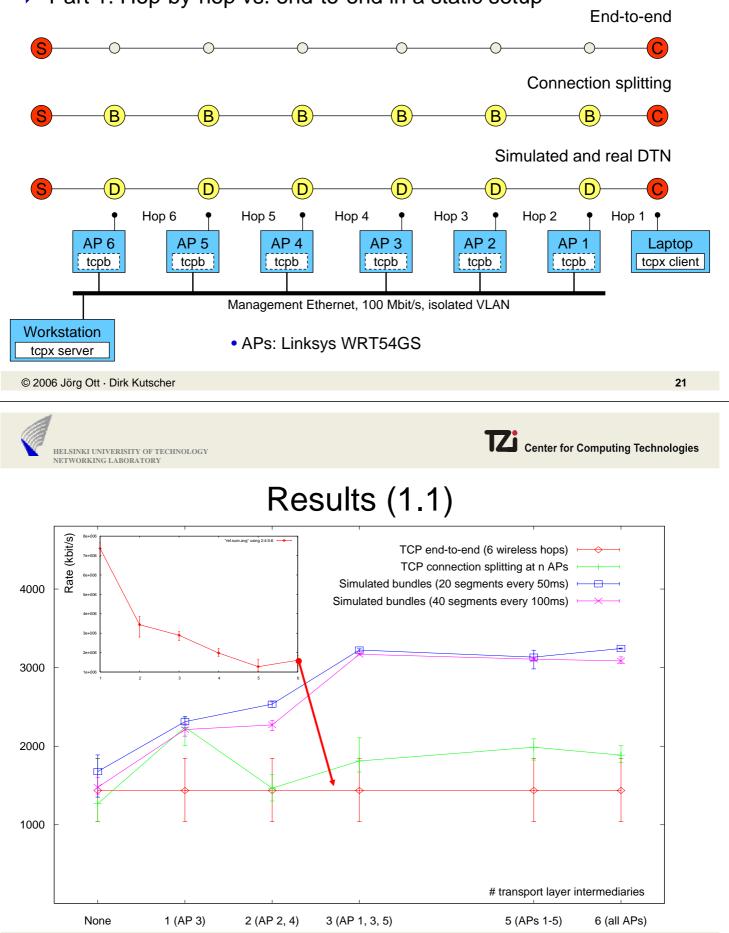




Performance Observations

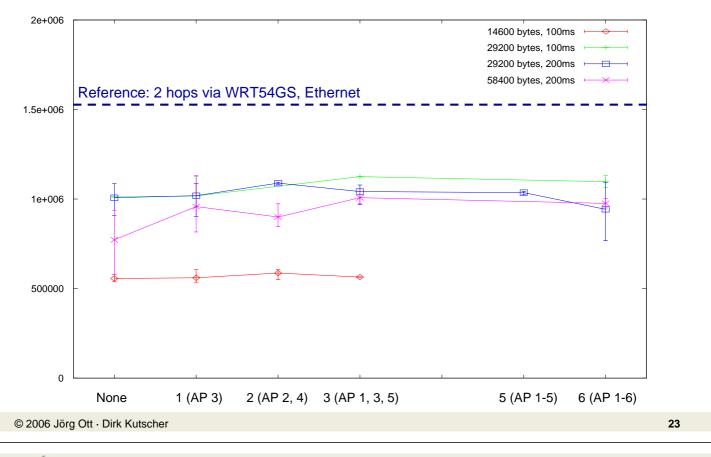
Measurements (1)

Part 1: Hop-by-hop vs. end-to-end in a static setup





Results (1.2): dtnd





HELSINKI UNIVERISITY OF TECHNOLOGY NETWORKING LABORATORY **TZi** Center for Computing Technologies

Center for Computing Technologies

Measurements (2)

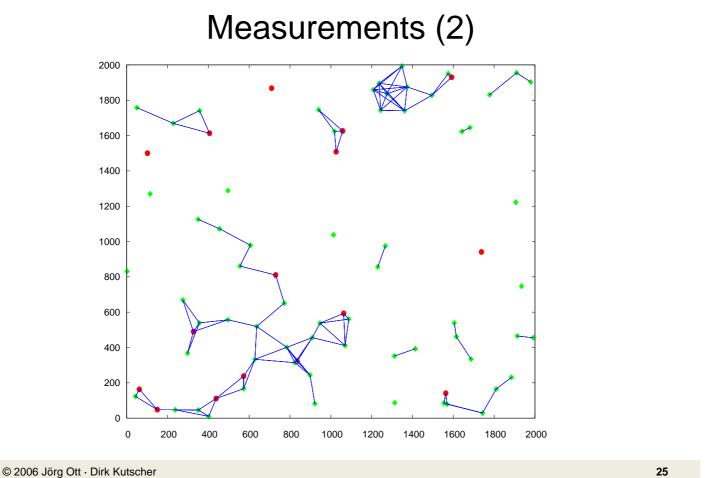
Emulation environment

- Dedicated emulation machine: 2 x Intel dual core @ 3.7 GHz (16 GB RAM)
- Xen 3.0 (unstable) for up to 100 virtual machines (32 MB memory each)
- Debian 2.6.16.13
- AODV UU from Uppsala University + AODV extensions implemented
- ns-2 + nsemulation patch to simulate motion + wireless connectivity
- tcpdump

AODV Extension Overhead

- · Repeated ping from one node to another
- 2000m x 2000m area, 200m wireless communication range, 0...5 m/s
- 20, 40, 80 nodes
- 0, 5, 10, 20, 50, 100% of which support DTN
- Compare against plain AODV overhead

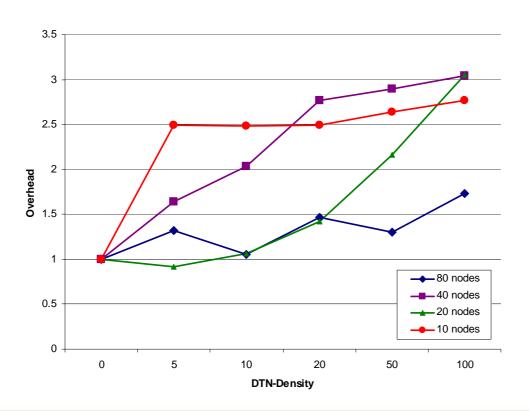




TZi Center for Computing Technologies

HELSINKI UNIVERISITY OF TECHNOLOGY NETWORKING LABORATORY

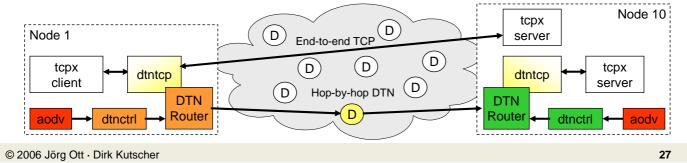






Measurement (3)

- dtnd from the DTN reference implementation (Intel, UC Berkeley)
- tcpx + dtntcp
 - Send via plain TCP end-to-end whenever a path exists upon transfer initiation
 - Otherwise fall back to DTN
- Initial tests with chain of 10 nodes and on/off links
- Proof of concept emulations
 - 2000m x 600m area, 10 nodes: random walk, 0...20m/s, 29200 bytes every 150s
 - Hack: n > m routing
 - Transmission from node 1 to node 10 in regular intervals
 - DTN augmented TCP-based delivery well in most cases with only 5min lifetimes





TZi Center for Computing Technologies

Discussion: API and Applications

- Current approach: implicit signaling (e.g., TCP connect())
 - Little chance to communicate EIDs
 - Approach limited to locating DTN routers
 - Cumbersome retrieval of "search results"
 - Correlation may be tricky
 - How long to wait? (TCP timeout? ...?)
 - Protocol-dependent feedback does not work with all protocols
- Better: dedicated API for explicit control
 - Trigger route search independently
 - Allow providing and retrieving parameters as needed
 - Current thoughts: interfacing through dtnd vs. bypassing dtnd

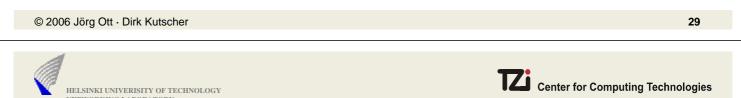
Applications need to be adapted for mobility anyway

Discussion: Routing Integration

- Simple case: just find DTN routers
 - Implicit service location

HELSINKI UNIVERISITY OF TECHNOLOGY TWORKING LABORATORY

- Advanced case: learn about DTN route metrics
 - Data volume + processing overhead vs. improved routing decisions
- How much subsequent handshaking is needed?
 - Mutual authentication (which may decide about accepting a bundle)
 - Possible other reasons for not accepting a bundle
 - How much information can and should we convey before? Taking into account that AODV is expensive
 - How will DTN routing protocols work in ad-hoc environments?
- How to embed routing hints in DTN bundles?
 - (and processing/queuing hints for the local DTN router)



Conclusion and Next Steps

- AODV extensions enables route search to locate DTN routers
 - Re-using a reactive routing protocol as opposed to separate service location
 - Insights into requirements for underlying routing protocols
- Allows applications to make e2e vs. hop-by-hop tradeoffs
- Allows DTN routers to discover peers beyond next hop
- Allows DTN routers to forward packets instead of bundles
- DTN-based communication may outperform end-to-end TCP

Next steps:

ETWORKING LABORATORY

- Interaction with emerging DTN routing protocols
- Performance tradeoffs in different MANET environments
- API, applications and policies
- Internet Draft describing all the details