



Engineering, Operations & Technology  
Boeing Research & Technology

Research & Technology

# On-Demand Dynamic Route Optimization Between Tunnel Endpoints

IETF RTGWG – March 28, 2011

Fred L. Templin  
Boeing Research & Technology  
[fred.l.templin@boeing.com](mailto:fred.l.templin@boeing.com)

# Document Status

- **New Section 5.14 of ‘draft-templin-intarea-vet’**
- **Submitted March 14, 2011**
- **Originally specific to ISATAP tunnels**
- **Now being generalized to route optimization on any large links**

# Route Optimization for Large Links - Problem Statement

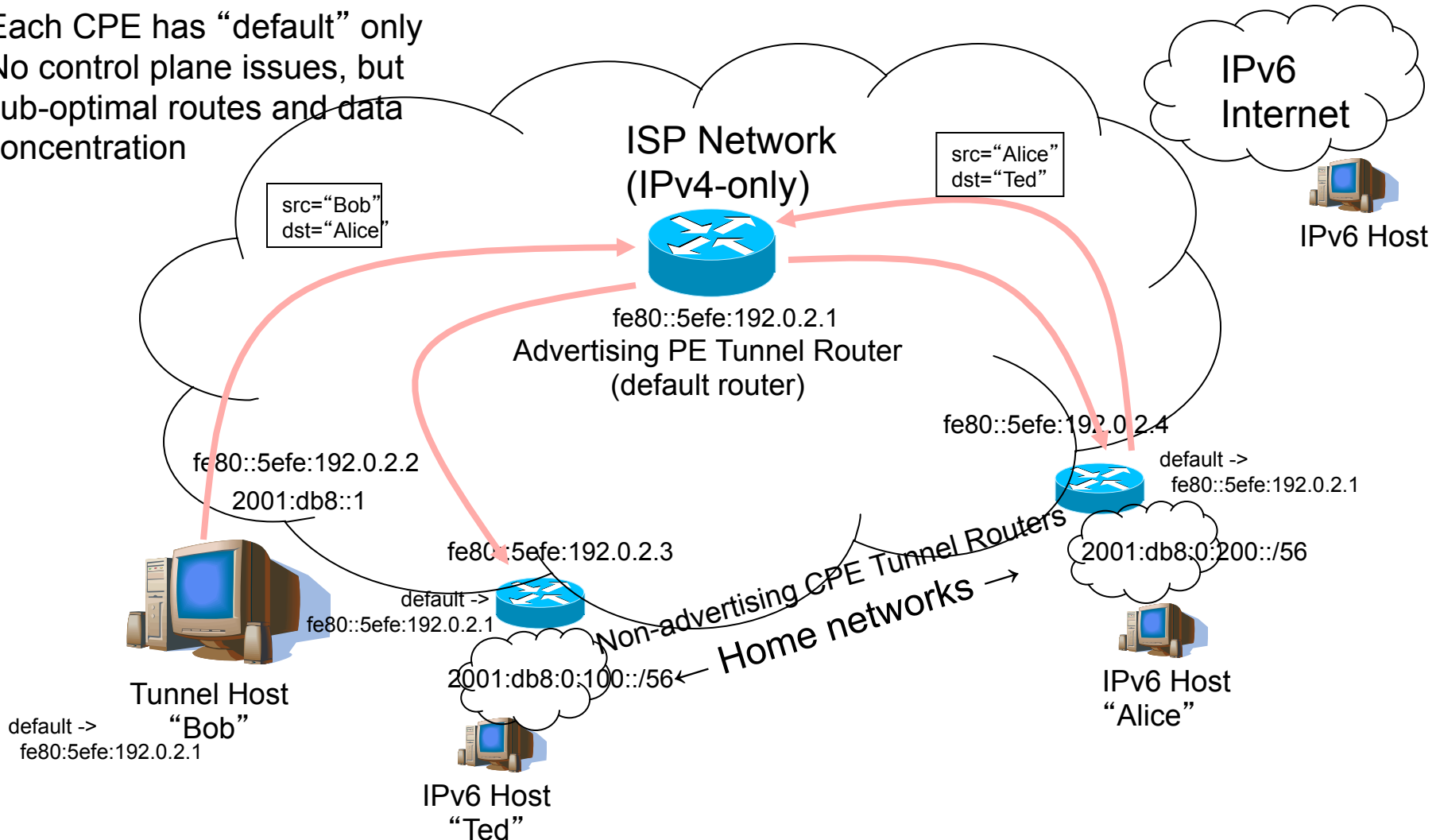
- **Some links connect many hosts and routers**
  - large campus LANs
  - bridged enterprise networks
  - cellular provider networks
  - aviation networks
  - disaster relief / defense networks
  - **Non-Broadcast, Multiple Access (NBMA) tunnels (e.g., ISATAP, 6over4, 6rd, vet, etc.)**
- **Traditional IGPs (e.g., RIPng, OSPFv3, etc.) don't scale well when the number of routes/routers is very large, or when many nodes are mobile**
- **Default routing via “hub” routers works, but:**
  - Triangular routing
  - Hubs forced to bear load in the data plane

# Specific Example – Tunnels over ISP Network

- **Many ISP networks still IPv4-only**
- **Growing customer requirement for IPv6**
- **IPv6-in-IPv4 tunnels seen as a near-term solution:**
  - **Tunnel as virtual NBMA link connecting many thousands of CPE routers**
- **PE routers in hub-and-spokes, mesh, or partial mesh**
- **Default routing via PE router works, but:**
  - **Traffic between CPEs takes longer paths than necessary**
  - **PE routers have to bear considerable load**
- **If each CPE is delegated IPv6 address/prefix, need a way to discover inter-CPE routes**

# NBMA Tunneling over ISP Network (ISATAP example)

- Each CPE has “default” only
- No control plane issues, but sub-optimal routes and data concentration

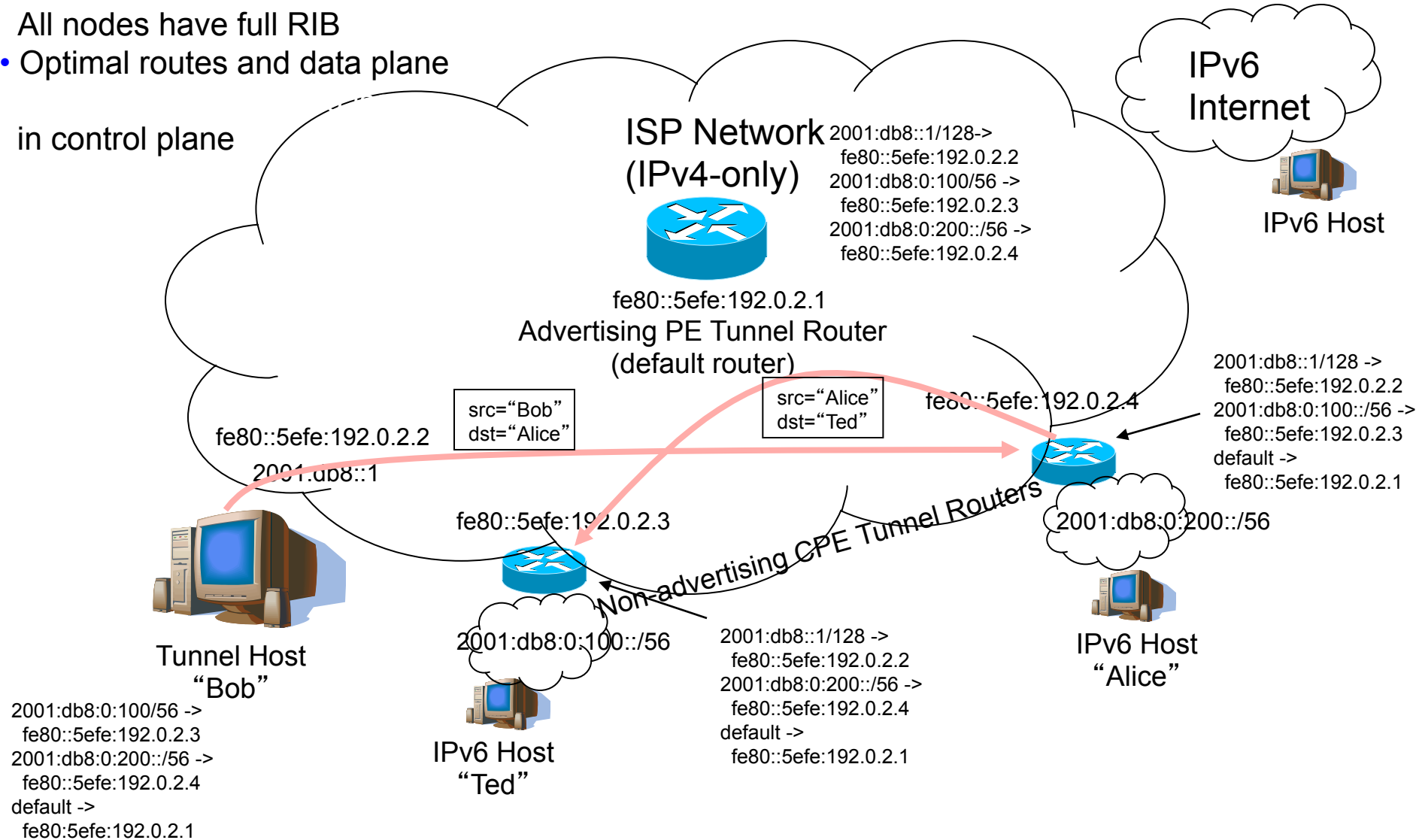


# Dynamic IPv6 IGP Between Tunnel Routers

All nodes have full RIB

- Optimal routes and data plane

in control plane



# Requirements for Dynamic Routing on Large Links

- **R1: Zero configuration on link nodes**
- **R2: Security based on chain-of-trust**
- **R3: Scale to support lots of nodes**
- **R4: Off-load performance-critical hub routers**
- **R5: Direct node-to-node route optimization**
- **R6: Route optimization for both routers and hosts**
- **R7: Support multiple levels of hierarchy**
- **R8: Do not circumvent ingress filtering**
- **R9: Do not expose packets to loss due to black holes**
- **R10: Support mobility**

# Idea – Use ICMP Redirects

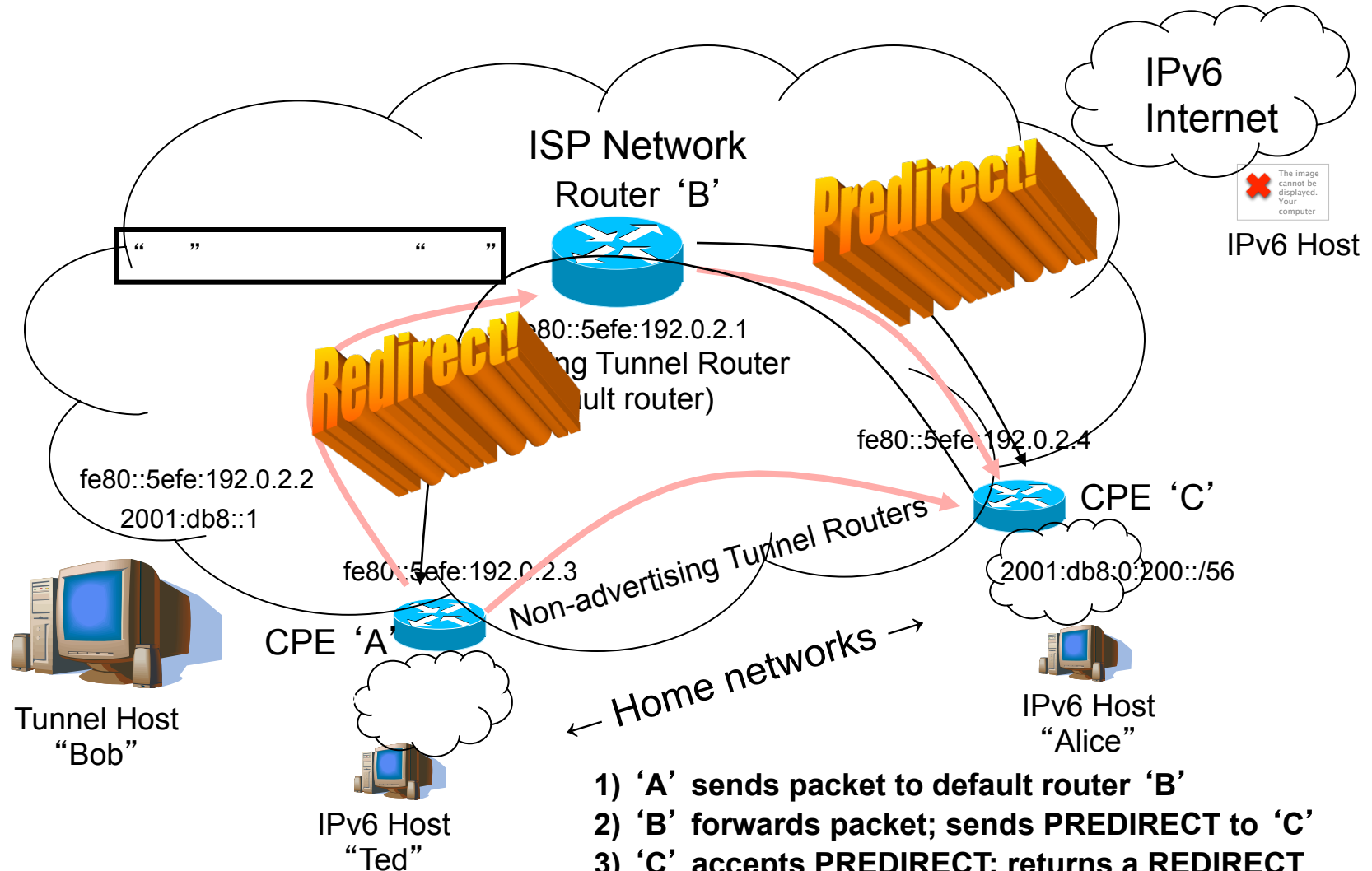
- **ICMP Redirect mechanism works between neighbors on multiple access links (i.e., physical or virtual)**
- **However, classical 1-way redirection does not support coordination between neighbors:**
  - **target has no way of knowing that the source is authorized to produce packets using a given source address (source could be spoofing)**
  - **source has no way of knowing that the target is prepared to accept its packets directly**
- **Also, ICMP Redirection doesn't support router-to-router redirects**



# Solution – “Augmented” Redirection

- **Initial packets from source go through “anchor” default router that is trusted by both nodes**
- **Anchor sends PREDIRECT message forward to target:**
  - target accepts the PREDIRECT since it trusts the anchor
  - target sends back a REDIRECT via the anchor
- **Anchor proxies the REDIRECT back to the source:**
  - source accepts the REDIRECT since it trusts the anchor
  - source sends future packets directly to target
- **Redirects / Predirects include Route Information Options (RIOs) that include prefix/length instead of just destination**
- **Redirects / Predirects can be router-to-router (i.e. and not just router-to-host)**

# Augmented Redirection in ISP Network



- 1) 'A' sends packet to default router 'B'
- 2) 'B' forwards packet; sends PREDIRECT to 'C'
- 3) 'C' accepts PREDIRECT; returns a REDIRECT
- 4) 'B' PROXIES REDIRECT back to 'A'
- 5) 'A' forwards future packets directly to 'C'

# Protocol Details

- **On receipt of PREDIRECT, Target creates FIB entry in “FILTERING” state and sets a 40sec timer**
  - **Target accepts packets only while in FILTERING state**
- **On receipt of REDIRECT, Source creates FIB entry in “FORWARDING” state and sets a 30sec timer**
  - **Source forwards packets only while in FORWARDING state**
- **State cleared when timers expire**
  - **Future packets from source re-start redirection process**
- **Source can send periodic Predirects to elicit Redirects from Target**
  - **Keeps route optimized while packets are flowing**

# Additional Considerations

- **Route Optimization is asymmetric in the forward direction from source to target**
  - If target sends packets in the reverse direction, a separate route optimization process is used (both directions managed independently)
- **Mobility supported:**
  - if final destination moves to a new network point of attachment, Target delivers the packet and returns a NACK
- **Backward compatibility supported:**
  - Based on standard ICMP Redirect messages (two new bits taken Reserved field)
  - Legacy nodes harmlessly ignore the messages and continue to use the default router

# Next Steps

- **Publish standalone document generalized to route optimization over any multiple access link that supports redirection**