Algorithms for computing Maximally Redundant Trees for IP/LDP Fast-Reroute

draft-enyedi-rtgwg-mrt-frr-algorithm-00

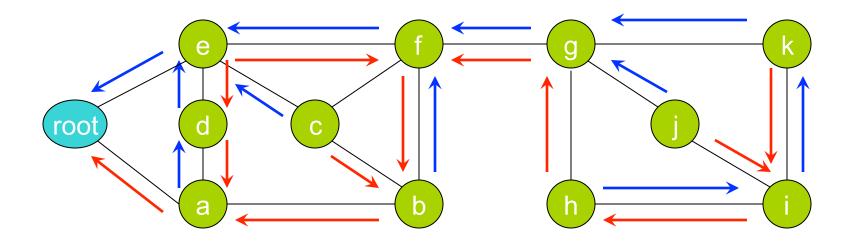
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MRT

- Maximally Redundant Trees
 - A pair of directed spanning trees
 - The common root is reachable along both of them
 - The two paths along the two trees are maximally disjoint



Why do we need this draft?

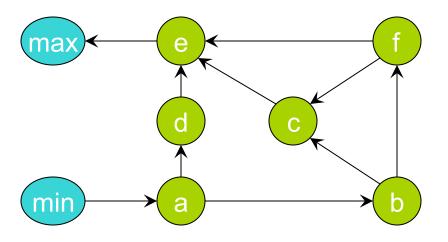
- We need a pair of MRTs rooted at each node
 - All the nodes should compute the same!
 - We will need standardization for MRT computation (algorithm) or results of that computation.

Principles

Partial order
ADAG
Blocks and GADAG

Partial order

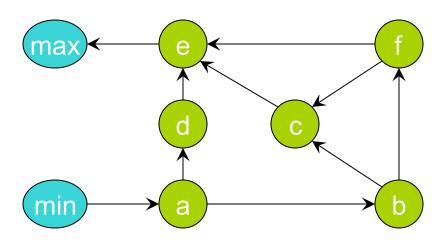
- Partial order of a set (e.g. set of nodes)
 - A relation like a normal set
 - Except: not all the elements can be compared
 - For some a and b neither a<b nor a>b
- Graph representation:
 - Directed Acyclic Graph (DAG)



- min<a<b<f<c<e<max
- a<d<e

Finding node-disjoint paths

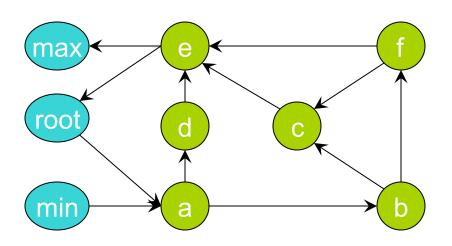
- Suppose that
 - We have a partial order of nodes
 - Exactly one min and max
 - Each node (except min and max) has a lower and greater neighbor
- Walk down and up
 - Min and max are reached
 - The two paths are node-disjoint!



- min<a<b<f<c<e<max
- a<d<e

Two paths to the same node

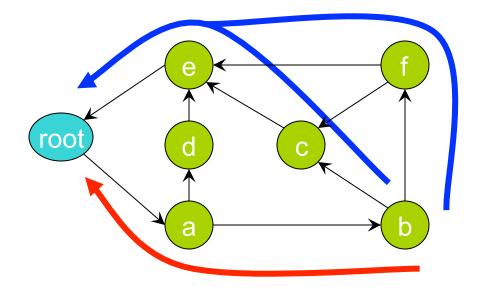
- DAG is not enough
 - Let min and max be the same node!
- Resulting graph is an Almost DAG (ADAG)
 - There is a single node, the root, such that without the root it is a DAG



- mipt<a
b<ff<c<e<mopst</pre>
- a<d<e

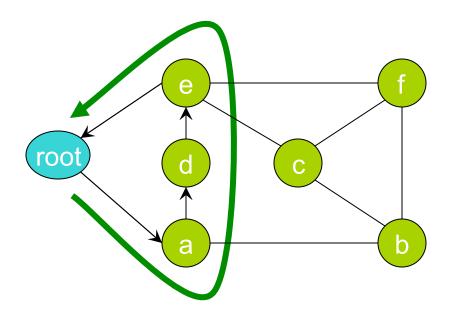
Redundant paths to the root

- Blue path:
 - Nodes must increase
- Red path:
 - Nodes must decrease
- Load sharing is possible



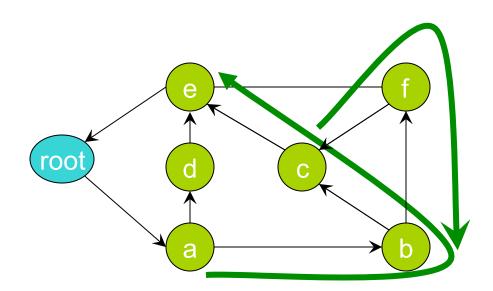
Finding an ADAG (2-connected networks)

- Phase 1 basic partial ADAG
 - Find a partial ADAG for a cycle containing the root
 - Use either direction
 - Extend partial ADAG into all nodes



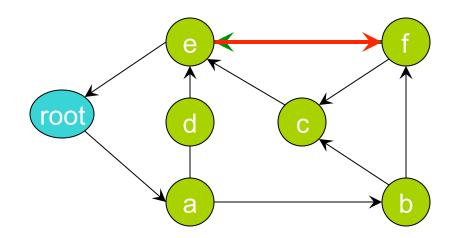
Finding an ADAG (2-connected networks)

- Phase 2 extending
 - Find a path from one "ready" node to the another
 - Nodes along the path must not be ready (except the endpoints)
 - Add the path to the ADAG in a "proper" direction



Adding not used links

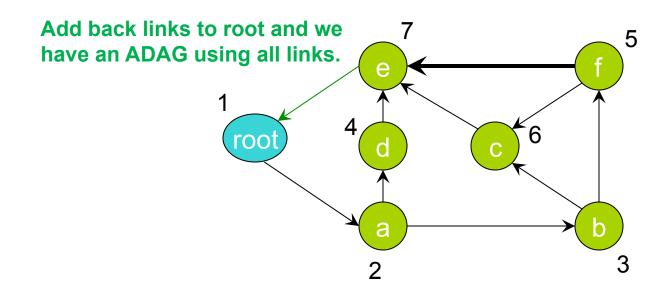
Some links may be out of the ADAG





How can ordering be kept up?

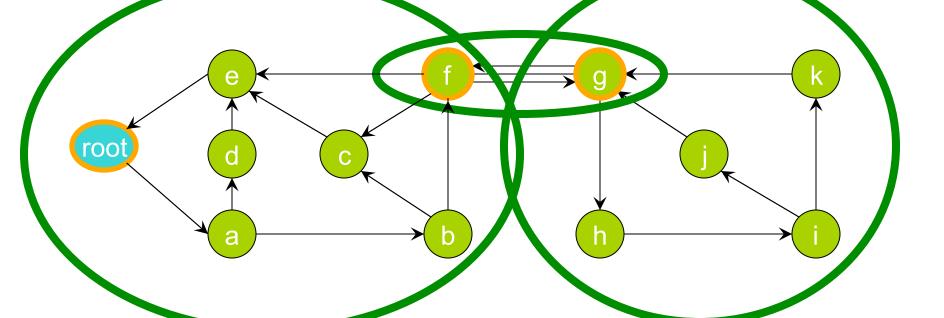
- ADAG is almost a DAG
 - Let root be now only the smallest one
 - Now, it's a DAG, create a topological sort
 - This is a total order
 - Add extra links with respect to this



What if the network is not 2-connected?

- We need to split the graph into blocks
 - Block:
 - Maximally 2-connected subgraph
 - Two connected nodes
 - (Isolated node)
 - Each block has its local-root
 - That is the cut-vertex towards the root
 - Compute an ADAG in all the blocks
 - This is a Generalized ADAG

Generalized ADAG



- Block1: root, a, b, c, d, e, f
- Block2: T. o.
- Block3: g, n, i, j, k

The algorithm

MRTs in a block
MRTs in the whole network

How to Find MRTs

- If it is complex, then we break the problem down
 - Transform network into its blocks
 - Find ADAGs in each block
 - Connect up the ADAGs to make a GADAG
 - Add all the other links in with the proper directionality
- From a GADAG, compute your next-hops to each destination
 - First for those in the same block
 - Destinations outside the block inherit their nexthops from a proxy in the block

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MRTs in a single block

- As the computing router S: From the GADAG, can use SPF and reverse SPF to find next-hops to all destinations in the same block
 - SPF gives nodes definitely greater
 - rSPF gives nodes definitely lesser
 - Remaining nodes are not ordered
- Then use some simple rules

MRTs in a single block: source perspective

- Find greater and lesser nodes
- Rules
 - 1. If **S < D** –
 - 2. If **S > D** increase to root
 - 3. No order decrease to root
 - 4. If D=root increase to root
 - 5. If S=root increase to D

- increase to D

decrease to root

decrease to D

increase to root

decrease to root

decrease to D

Routing table of **node c** (S = c):

+ root 0 d
a b

Dest (D)	Rule Used	Blue Next-hop	Red Next-hop
а	2	е	b
b	2	е	b
d	3	b	е
е	1	е	b
f	2	е	f
root	4	е	b

MRTs in a single block: destination perspective

- Find greater and lesser nodes
- Rules
 - 1. If **S < D** –
 - 2. If **S > D** -
 - 3. No order decrease to root
 - 4. If d=root increase to root
 - 5. If s=root increase to D

increase to D

increase to root

decrease to root

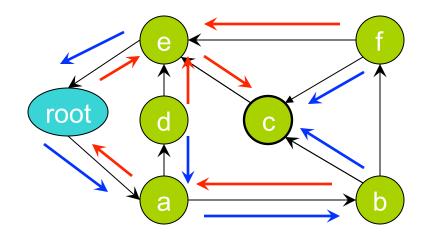
decrease to D

increase to root

decrease to root

decrease to D

Destination: node c (D = c)



Src (S)	Rule Used	Blue Next-hop	Red Next-hop
а	1	b	root
b	1	С	а
d	2	a	е
е	3	root	С
f	2	С	е
root	1	а	е

How to Find MRTs

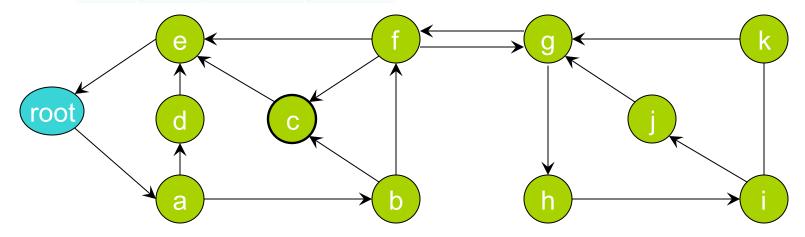
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Inter-block MRTs

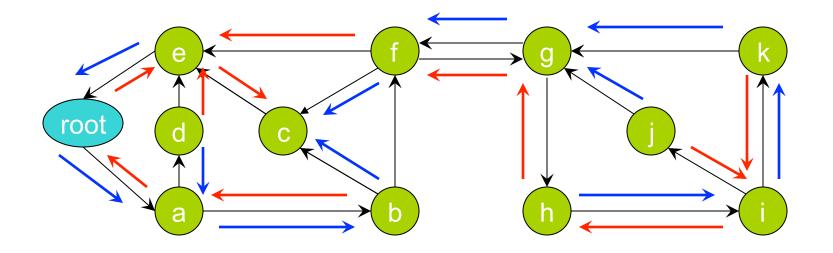
Proxy node: the last vertex in the block to the destination (along *any* path)

Dest (D)	Rule Used	Blue Next-hop	Red Next- hop
а	2	е	b
b	2	е	b
d	3	b	е
е	1	е	b
f	2	е	f
root	4	е	b

Dest (D)	Proxy	Blue Next- hop	Red Next- hop
g	f	е	f
h	f	е	f
i	f	е	f
j	f	е	f
k	f	е	f



Example – destination is node C



- Block1: root, a, b, c, d, e
- Block2: e, f
- Block3: f, g, h, i, j

Summary

- Algorithm
 - Find GADAG
 - ADAG in each block
 - Add not used links
 - Find next-hops along the MRTs
 - Do an SPF and an rSPF to find ordered nodes
 - Use rules to find NHs your block
 - Find proxy nodes

Thanks for the attention