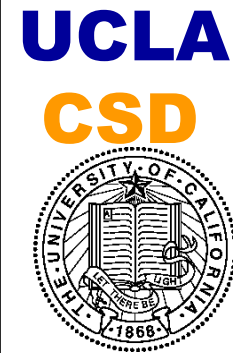


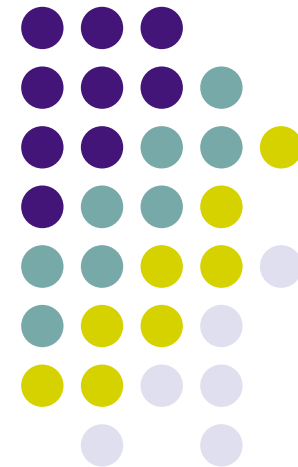
TCP over Network Coding



Presented by
Nadia Boukhatem
Télécom Paris

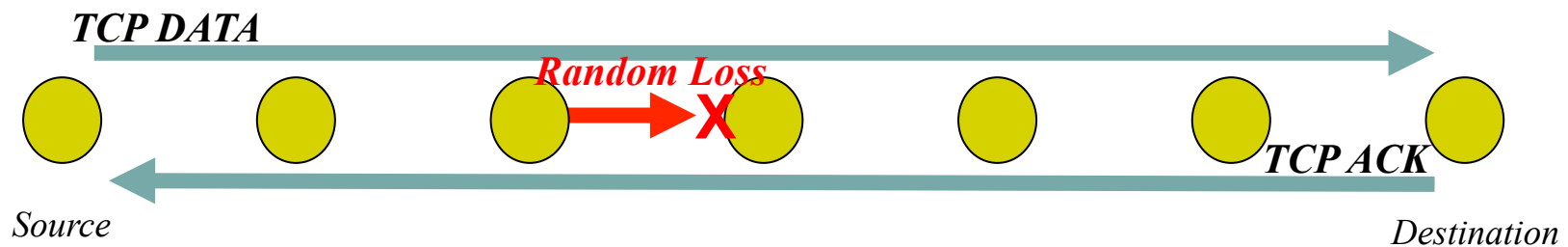
Joint Work with
Chien-Chia Chen, Mario Gerla, M.Y. Sanadidi

Network Research Lab
CSD, UCLA



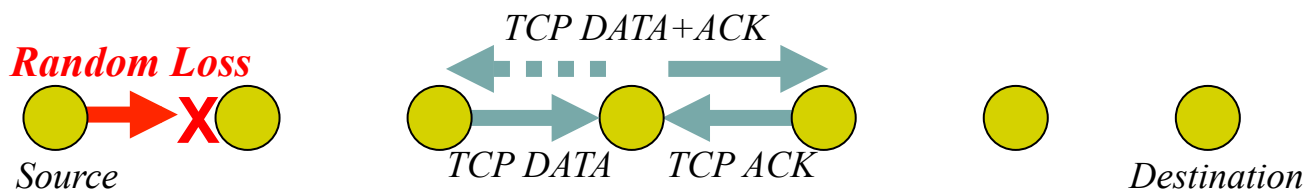
TCP in Wireless Networks

- Problems
 1. **random** losses are misinterpreted as **congestion**
 2. TCP DATA and ACK flows **contend** for the same shared medium



ComboCoding— Combined Intra- and Inter-Flow Coding

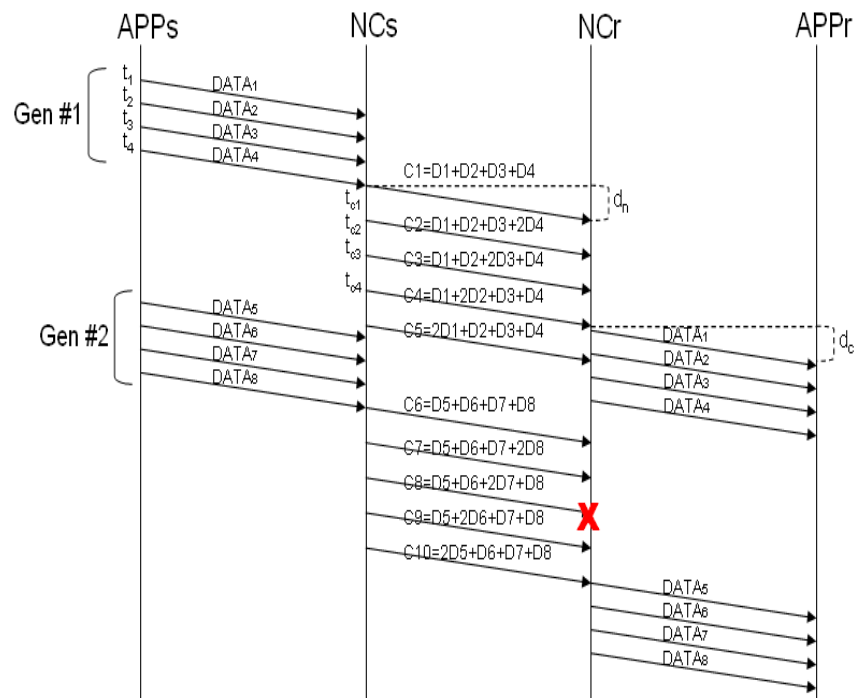
- To mitigate high loss
 → **Intra-Flow Coding** (Pipeline Coding)
 - Uses random linear coding to recover losses
- To mitigate DATA-ACK interference
 → **Inter-Flow Coding** (PiggyCode)
 - Opportunistically XOR DATA and ACK at relays
 - Mixing only DATA and ACK within the same TCP flow
- **Transparent to Upper/Lower Layers**



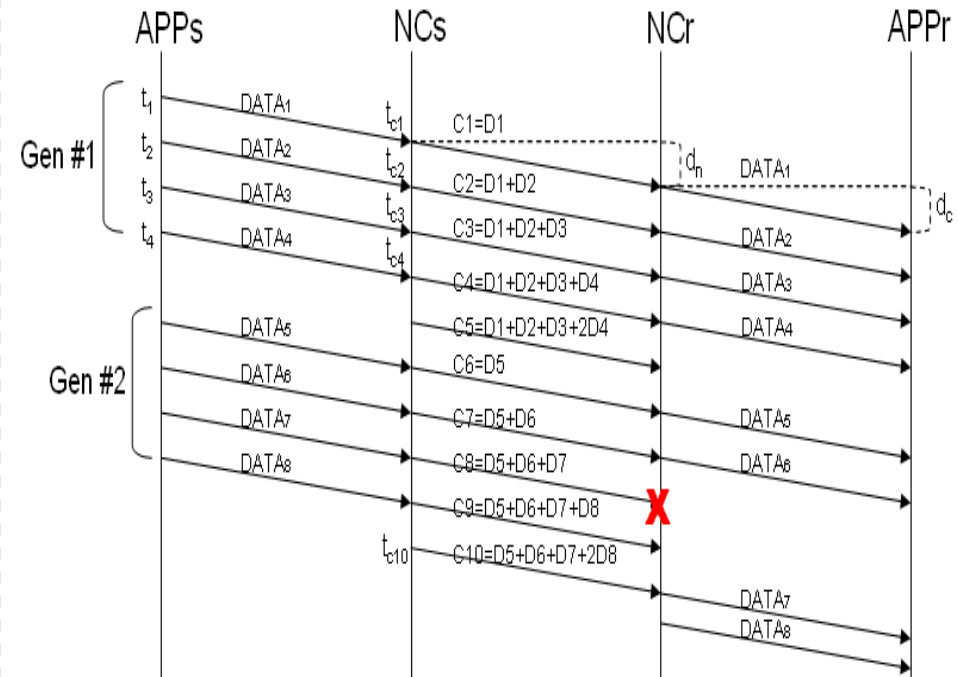
Intra-Flow Coding (Pipeline Coding)

- Benefits
 - Reduced delay
 - Improved throughput
 - Transparency to higher layers

Batch Coding



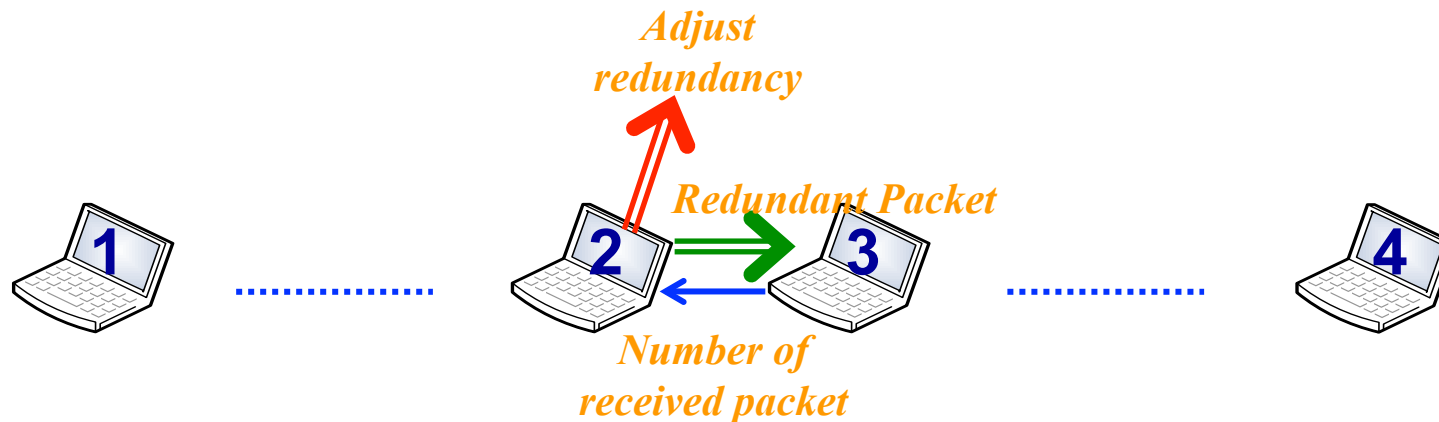
Pipeline Coding



2012/03/27

Adapt to Varying Losses

- Link error rates are changing at all times



- Each node stamps “**number of received packet**” in packets header
 - Upstream node receives it
 - It adjusts **link** coding redundancy based on **successful delivery** (to the next hop)

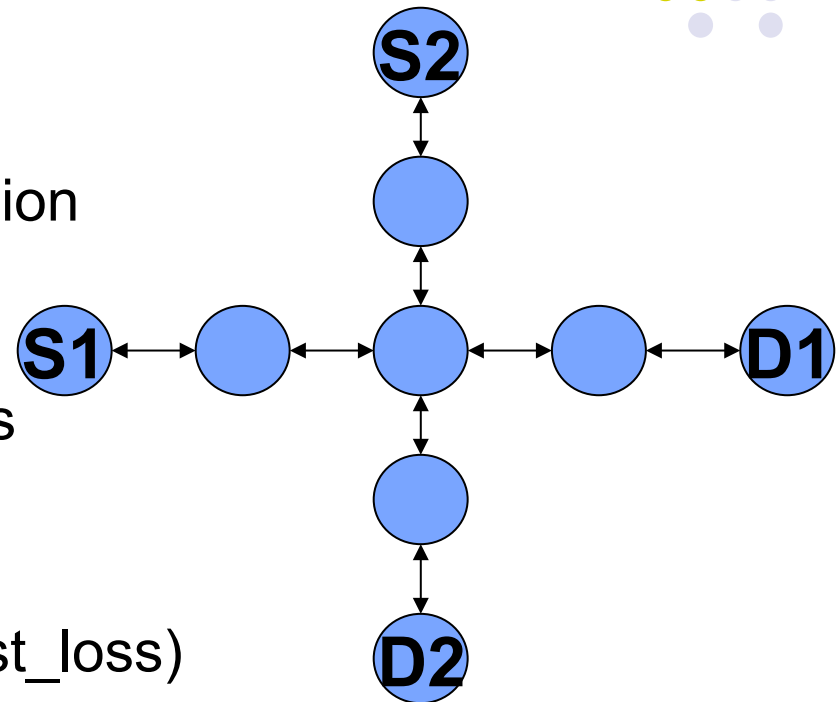
Simulation Evaluation

- Simulator: Qualnet 4.5
- Coding schemes are implemented at Network Layer as a special type of routing protocols
- TCP-NewReno is chosen for the transport layer protocol
- Our previous work has shown the effectiveness of our coding scheme
- We focus on **fairness** and **friendliness** comparisons in this presentation

Simulation Configuration



- 802.11g Unicast at 54Mbps
 - CSMA/CA
 - RTS/CTS is DISABLED
 - MAC ACK and MAC retransmission (up to 7 times)
 - Promiscuous Mode ENABLED
- Traffic: 2 FTP/TCP-NewReno Flows
- Gen size: 16
- Base Redundancy $K = 0.65$
- Adaptive Redundancy = $K + 1/(1 - \text{est_loss})$



Simulation Setup

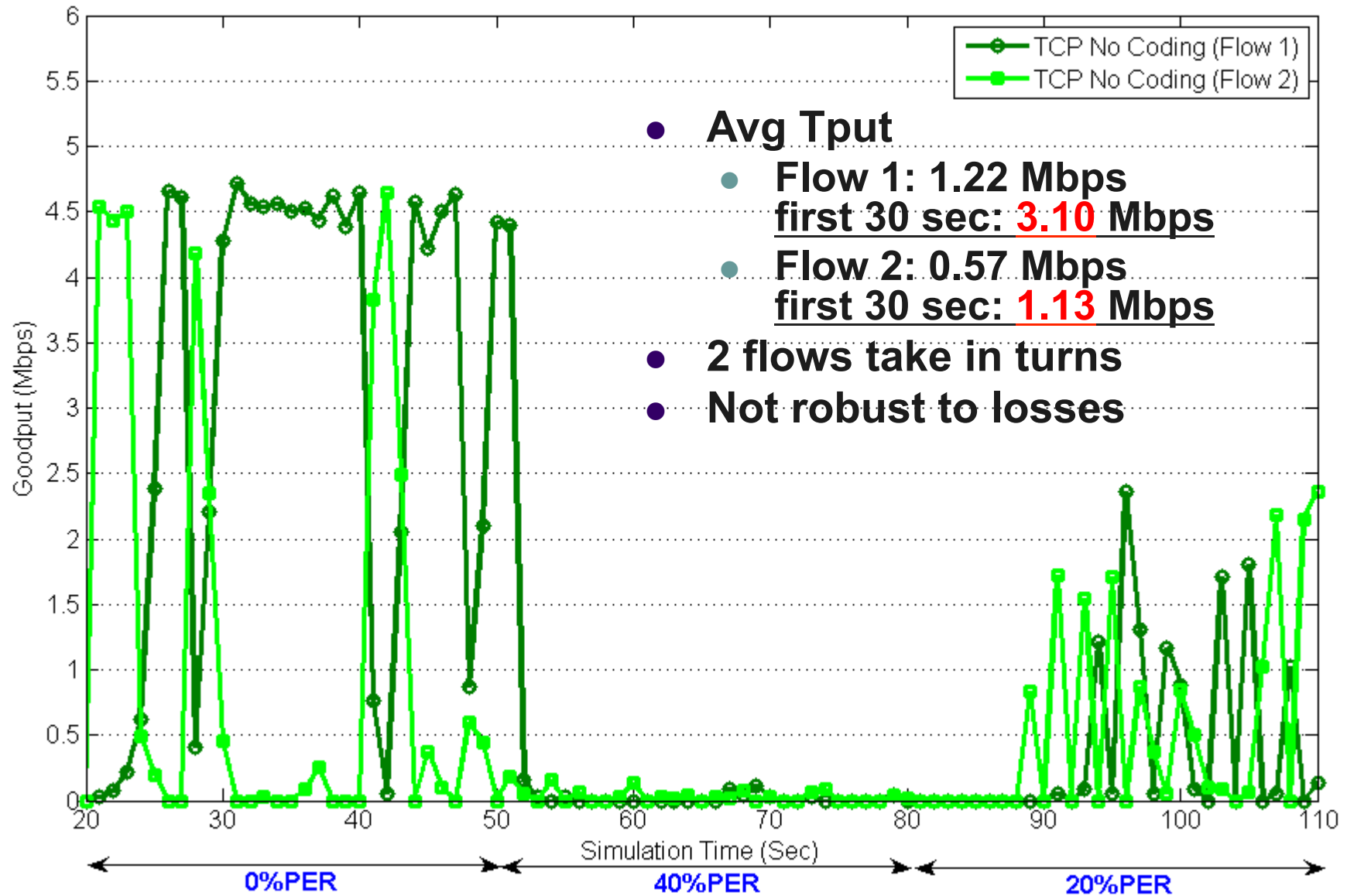
- 2 TCP Flows in 6 Runs

		NC	Redundancy
Run 1	TCP Flow 1	No	No
	TCP Flow 2	No	No
Run 2	TCP Flow 1	ComboCoding	Adaptive
	TCP Flow 2	ComboCoding	Adaptive
Run 3	TCP Flow 1	ComboCoding	Adaptive
	TCP Flow 2	No	No
Run 4	TCP Flow 1	ComboCoding	Adaptive
	TCP Flow 2	No	Adaptive
Run 5	TCP Flow 1	ComboCoding	No
	TCP Flow 2	No	No
Run 6	TCP Flow 1	No	Adaptive
	TCP Flow 2	No	Adaptive

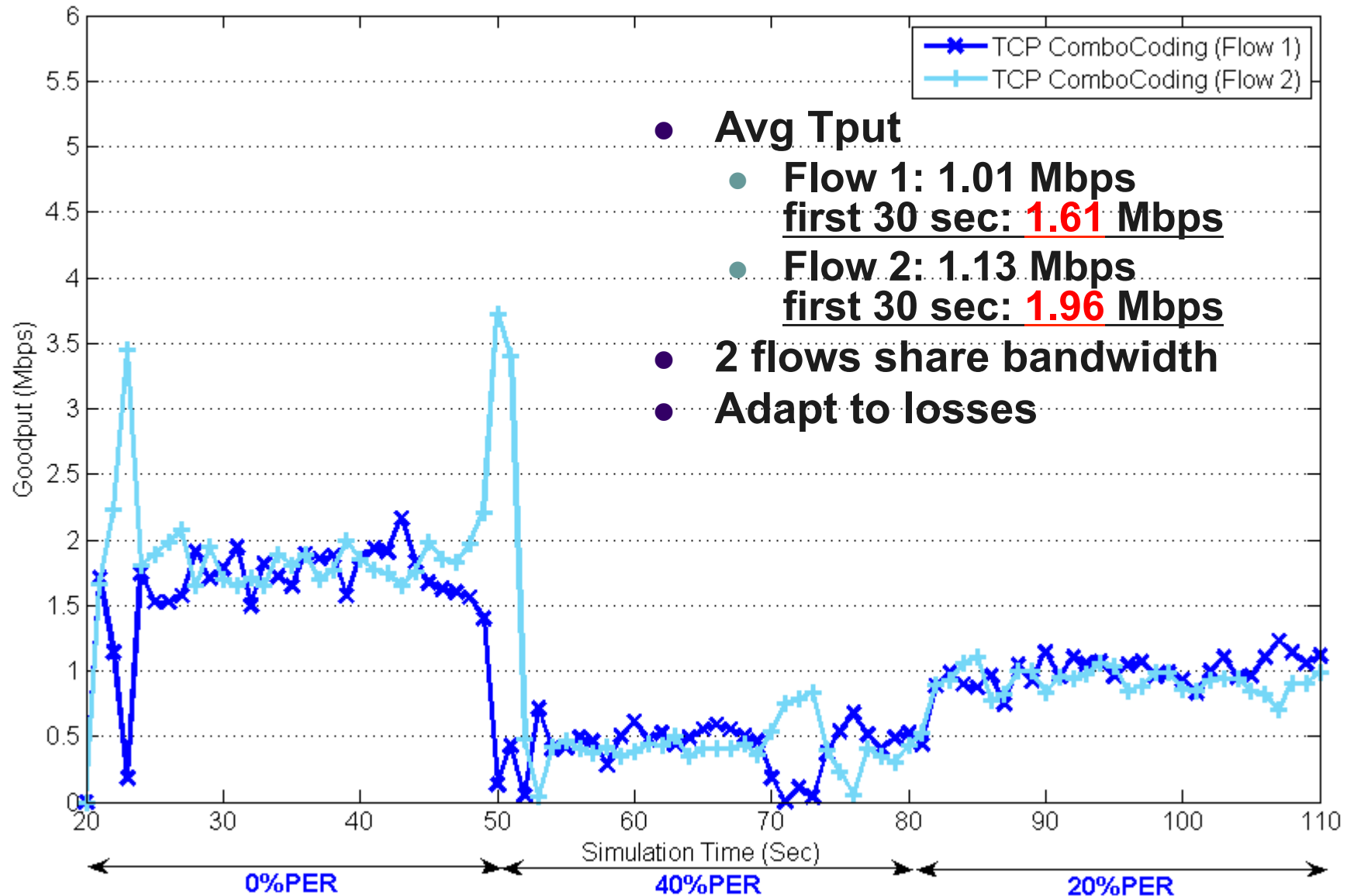
- Simulation time: 110 seconds (FTP starts at 20 sec)
- Vary per link Packet Error Rate over time
 - 20~50 sec: 0% PER
 - 50~80 sec: 40% PER
 - 80~110 sec: 20% PER

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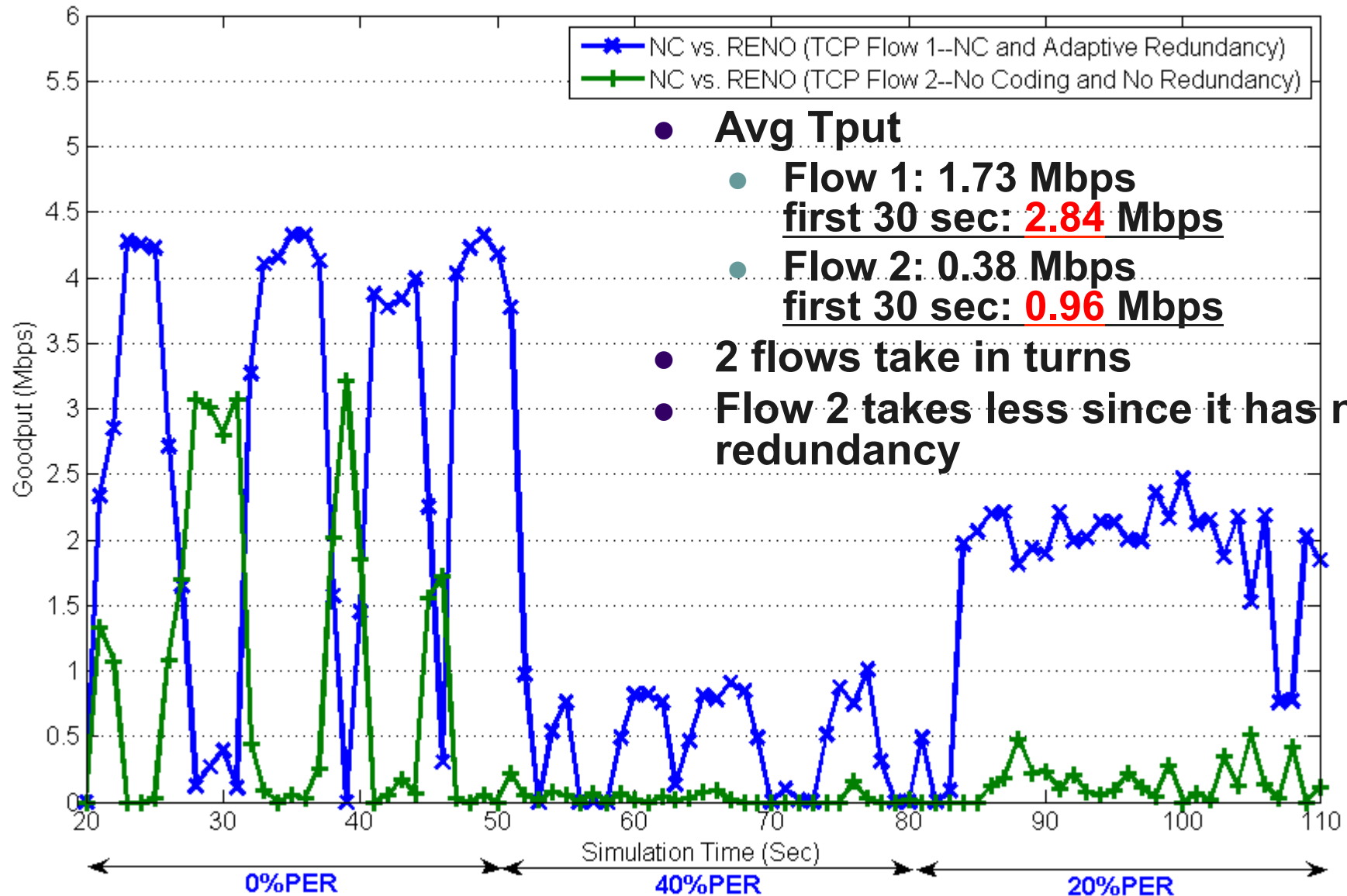
Run 1: No Coding No Redundancy



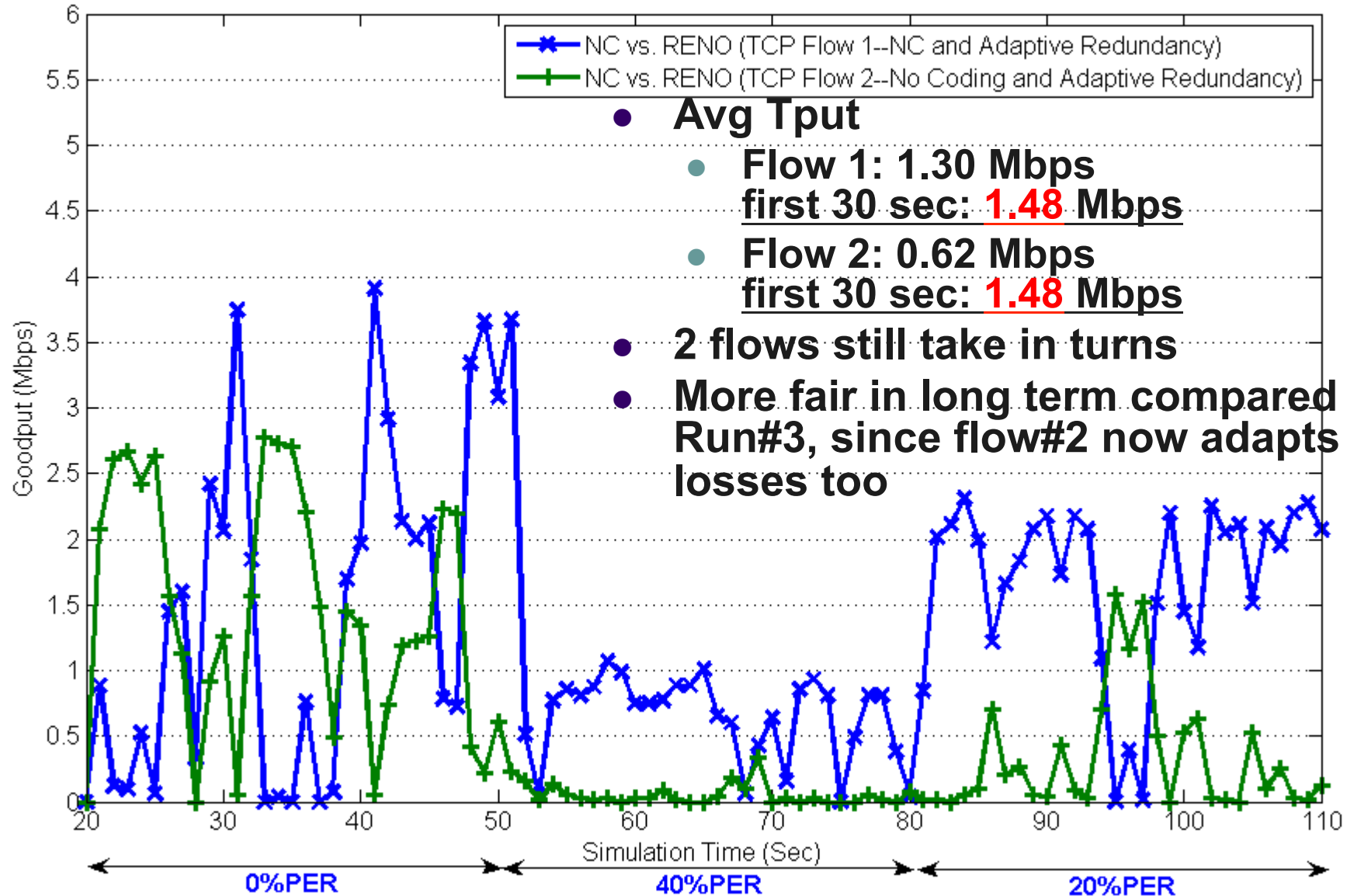
Run 2: ComboCoding with Adaptive Redundancy



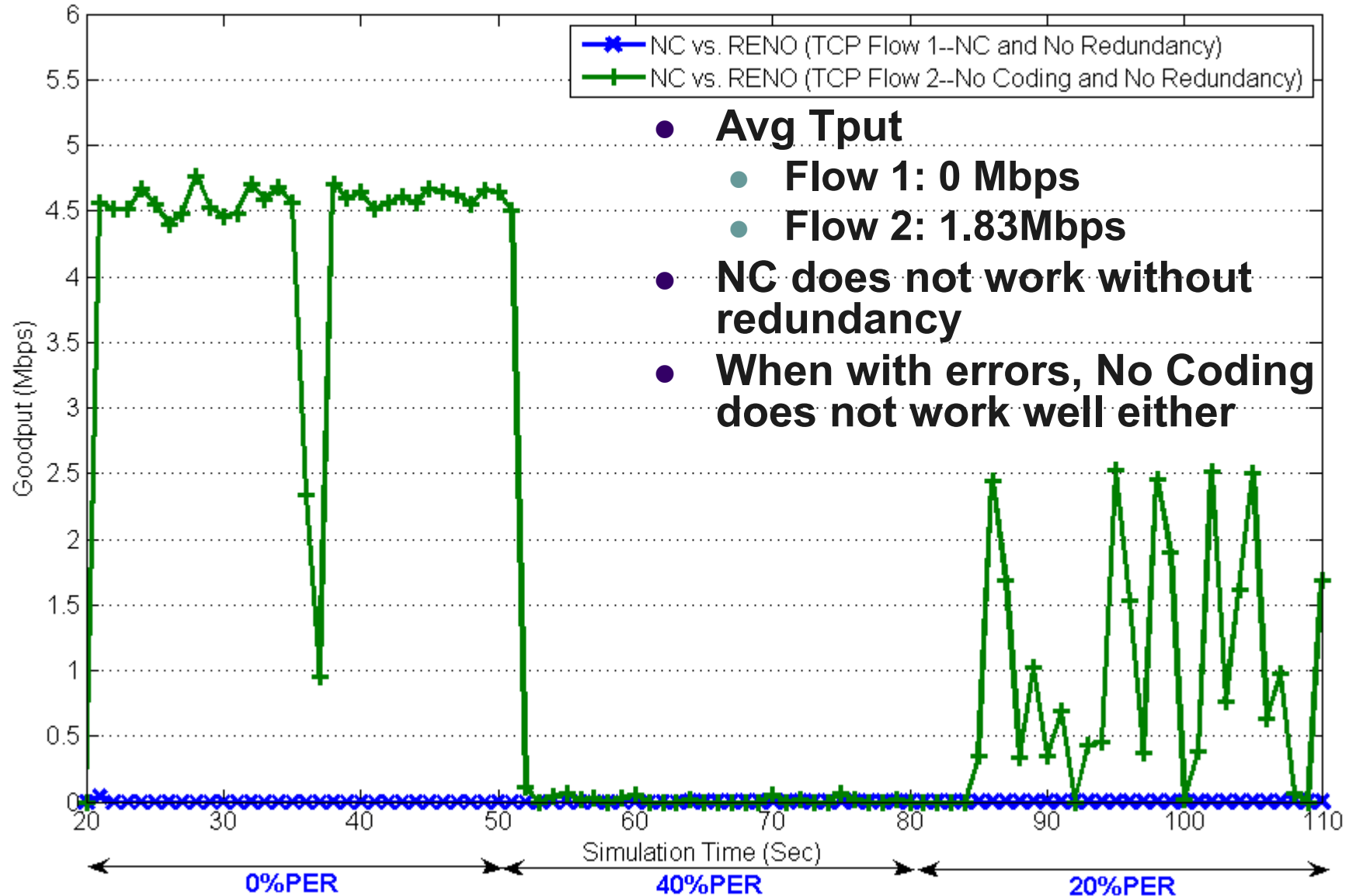
Run 3: No Coding vs. ComboCoding



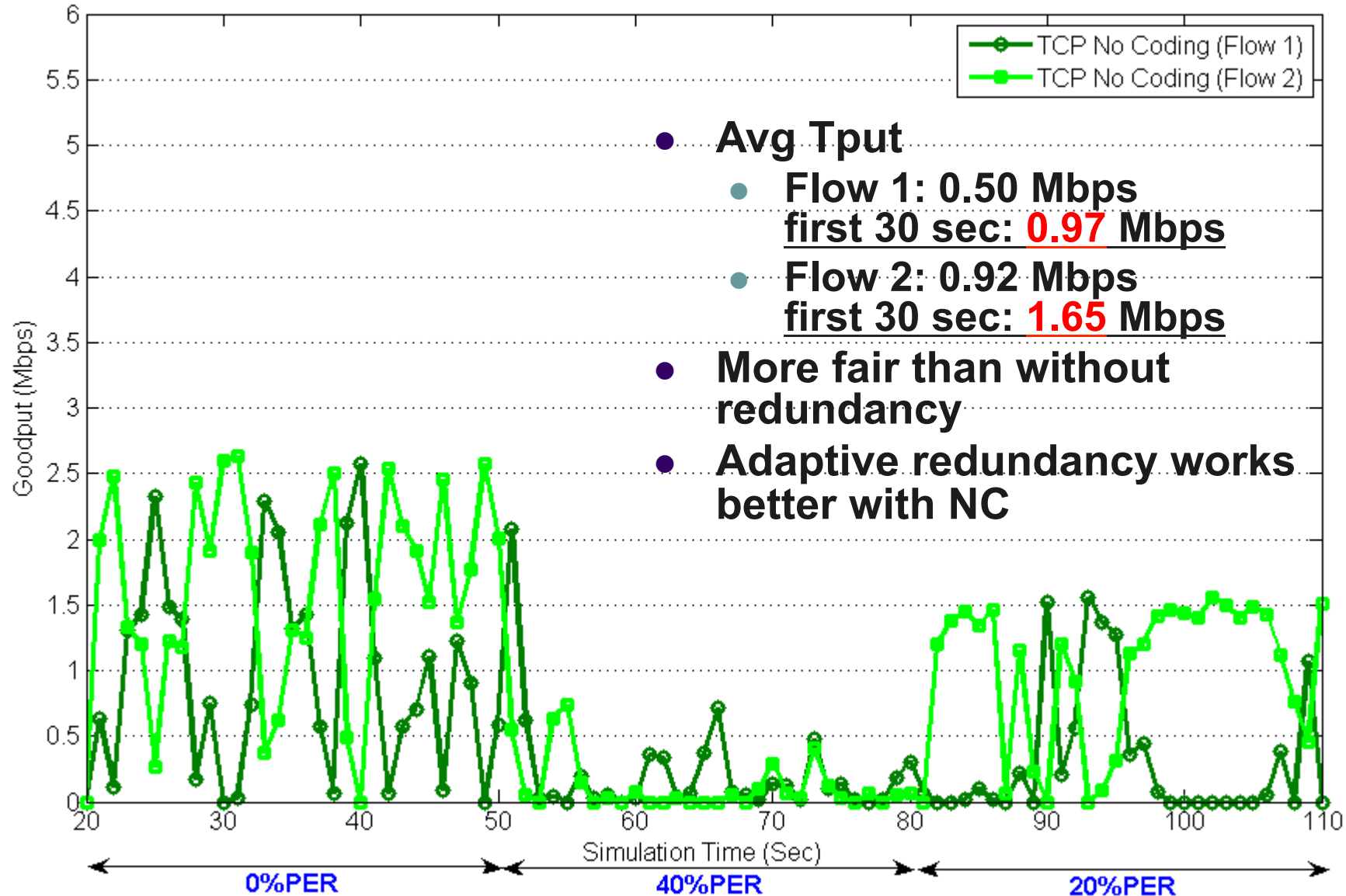
Run 4: Adaptive Reno vs. ComboCoding



Run 5: No Coding vs. NC No Redundancy



Run 6: No Coding with Adaptive Redundancy



Conclusion

- TCP/NetCode works much better than TCP for random loss channels
- TCP/NetCode is intra-fair and more stable than TCP
- TCP/NC can coexist with TCP (although there is some unfriendliness due to the adaptive redundancy)
- Current work
 - Adaptive redundancy control should employ loss discrimination techniques to have a more accurate random error estimate

Related Work



- "On the impact of random losses on TCP performance in coded wireless mesh networks," INFOCOM 2010 by Prof. Ros et al.
 - Study applies only to the single hop, opportunistic coding between different flows (ie inter-flow coding), like COPE
 - Interflow -coding improves throughput, but is known to be vulnerable to random errors
 - Two links correct reception requirement and packet-loss synchron across TCP flows are unique of Interflow-coding
- We use standard TCP NR and only Intraflow Coding:
 - Interflow coding applied only to DATA and ACKs



THANK YOU 😊