# ICNRG Fairness Discussion

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#### Motivation

- Blind use of AIMD may cause fairness issues when in-network caching is used
  - Less bandwidth available for less popular contents
- What is fairness in an Information-centric network?
  - How to measure this fairness (i.e., metric)?

### Fairness in term of what?

- Bandwidth?
- Memory?
- Delay?
- ...?
- In the rest of the presentation, resource is simply written r

#### Flow fairness

- $r_f$ , resources of flow f:
  - r / # flows (e.g.,  $r_f = r/5$ )



### Content fairness

- $r_c$ , total resources for flows of content c
  - $r / # contents (e.g., r_c = r/2)$



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# Customer fairness

•  $r_u$ , total resources for flows of customer u:

• r / # customers (e.g.,  $r_u = r/3$ )



# Link/interface fairness

- *r<sub>i</sub>*, total resources for interface *i*:
  - r / # incoming interfaces (e.g.,  $r_i = r/3$ )



# Network fairness

- $r_f$ , resources of flow f on path p:
  - r / # flows on p



Flow (one color per path)

# Two options to enforce fairness

- End-to-end
  - consumers deal with rate control
- Hop-by-hop
  - routers deal with rate control
- How to cope with demand aggregation that hides consumers?

# Summary

- What is a flow?
- What is a resource?
- What fairness do we want to reach and how to measure it?