
Alternative Security Solutions for MPTCP

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draft-paasch-mptcp-lowoverhead-00
draft-paasch-mptcp-ssl-00

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A low-overhead, low-security MPTCP for data centers

Motivation

- Need for low-overhead MPTCP handshake
 - Data centers
 - Non security-critical traffic

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Goals:

- Reduce chances of token collision
- Avoid extensive HMAC-calculations
- Allow stateless SYN+JOIN handling

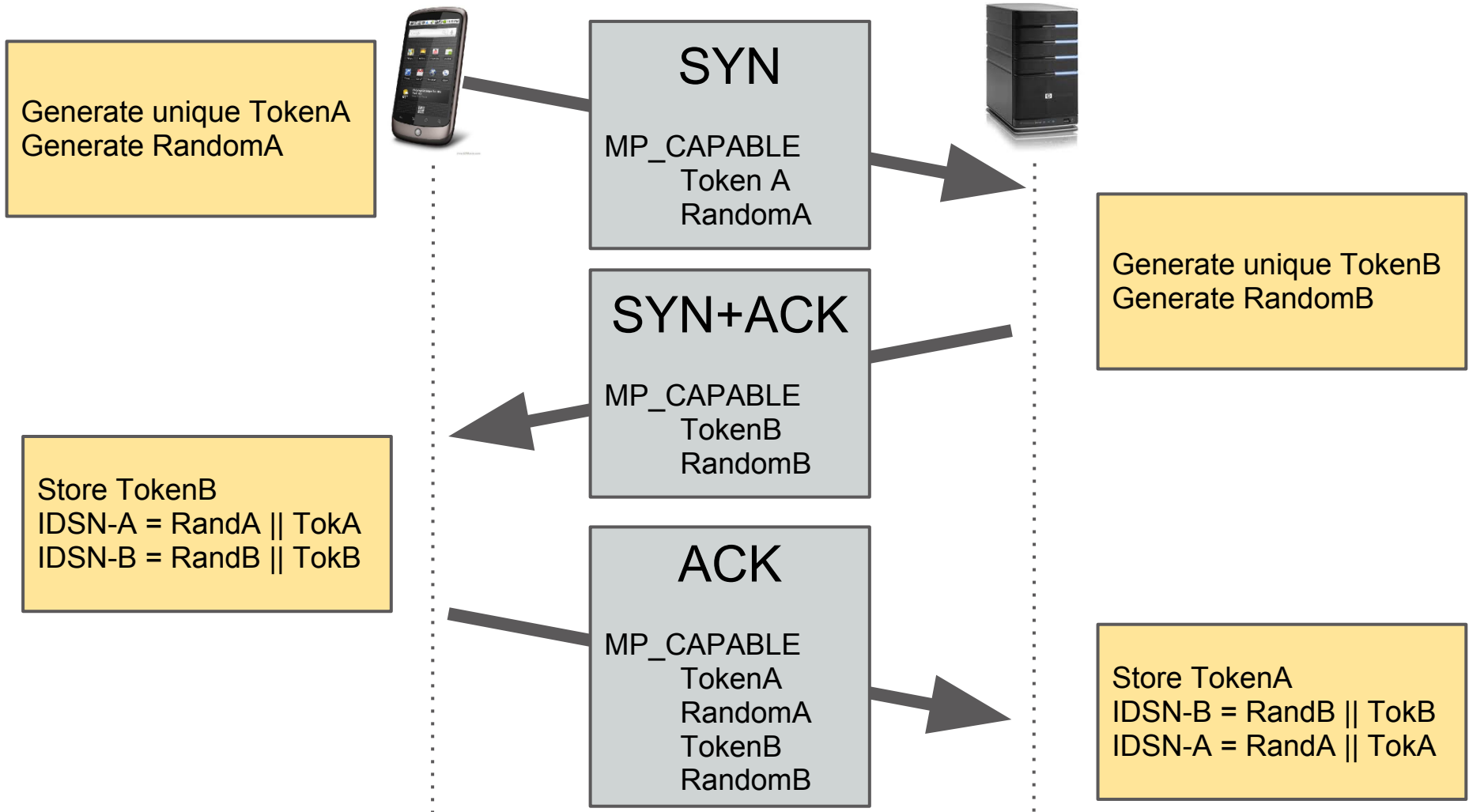
Non-goal:

- Authenticate new subflows
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Low overhead initial handshake

- Purposes of the initial handshake:
 - Detect whether the peer supports MPTCP
 - Exchange a connection identifier (Token)
 - Agree on the Initial Data Sequence Number
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Low overhead initial handshake



Low overhead initial handshake

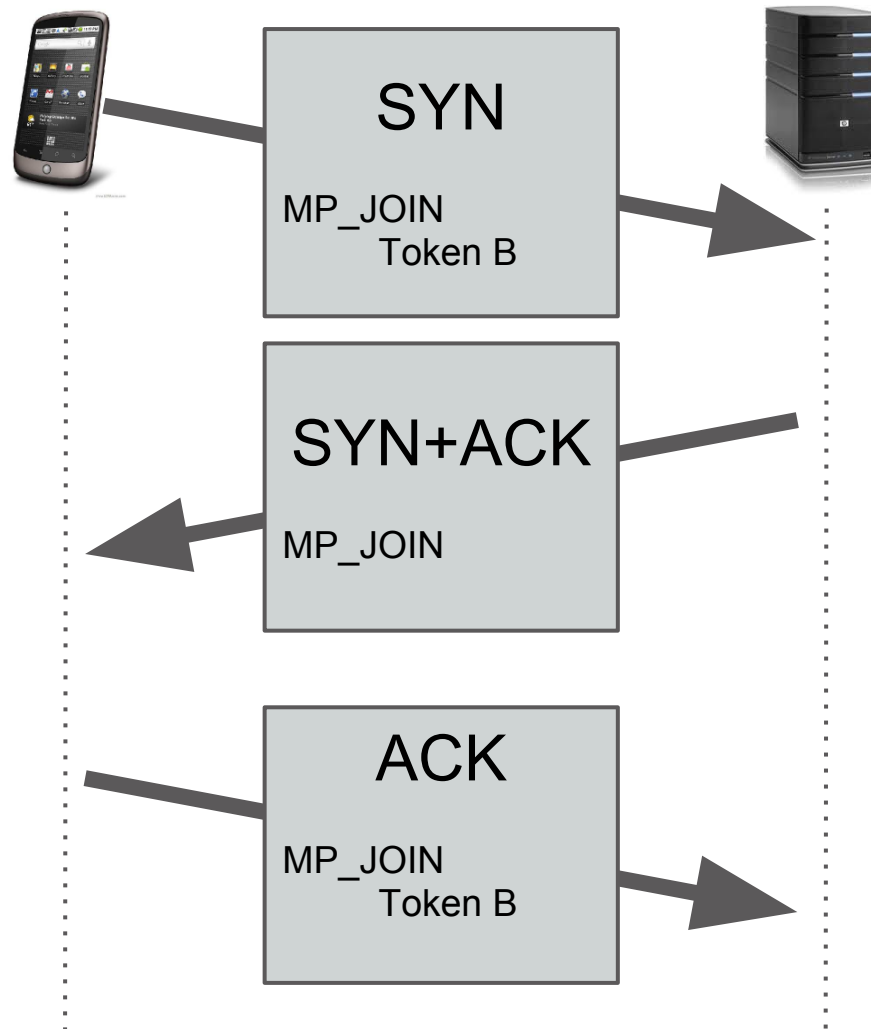
Generating the token:

- Could be a simple counter
- Better: block-cipher (RC5) of the counter plus a local secret

Benefits:

- Guarantees a high probability of uniqueness
 - No need to compute a SHA-1
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Low overhead additional subflows



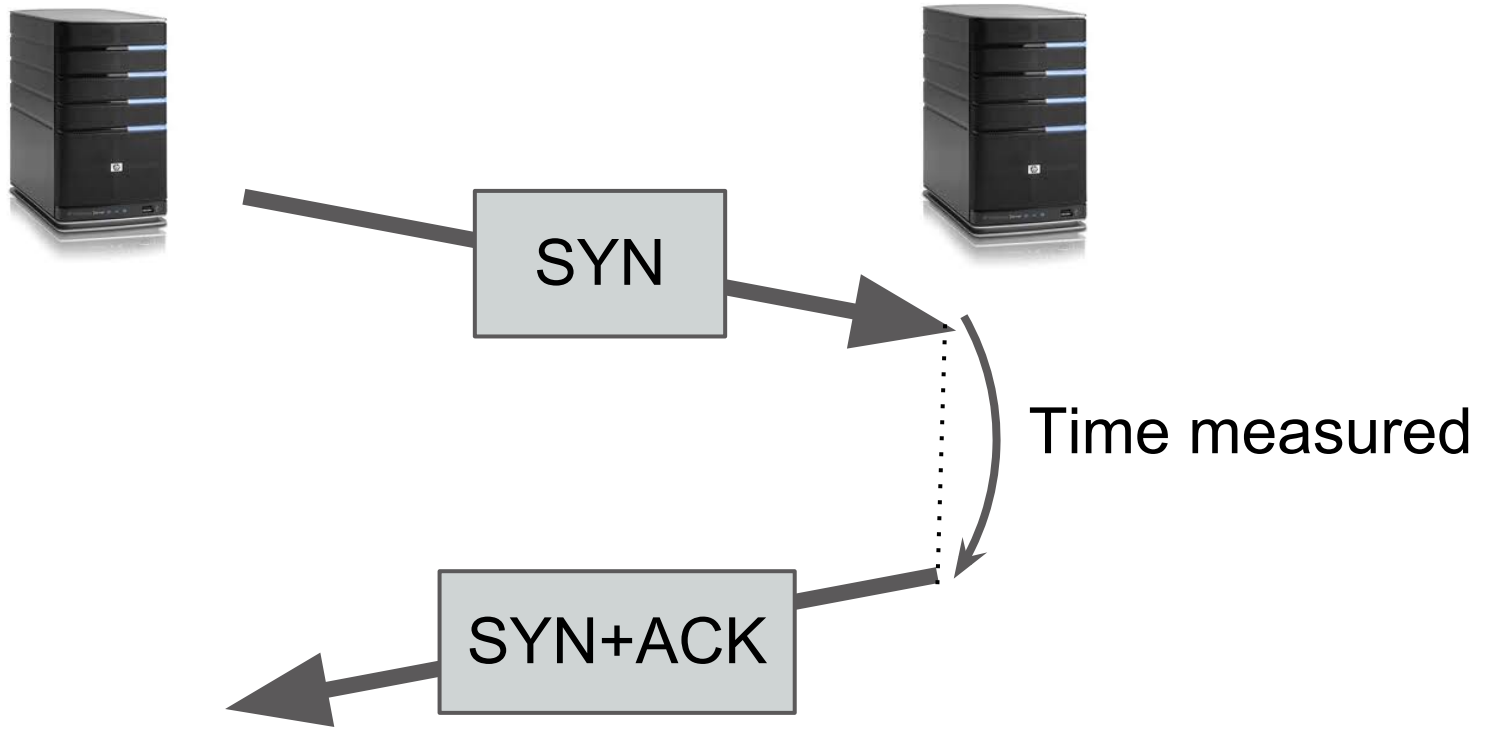
Low overhead additional subflows

Benefits:

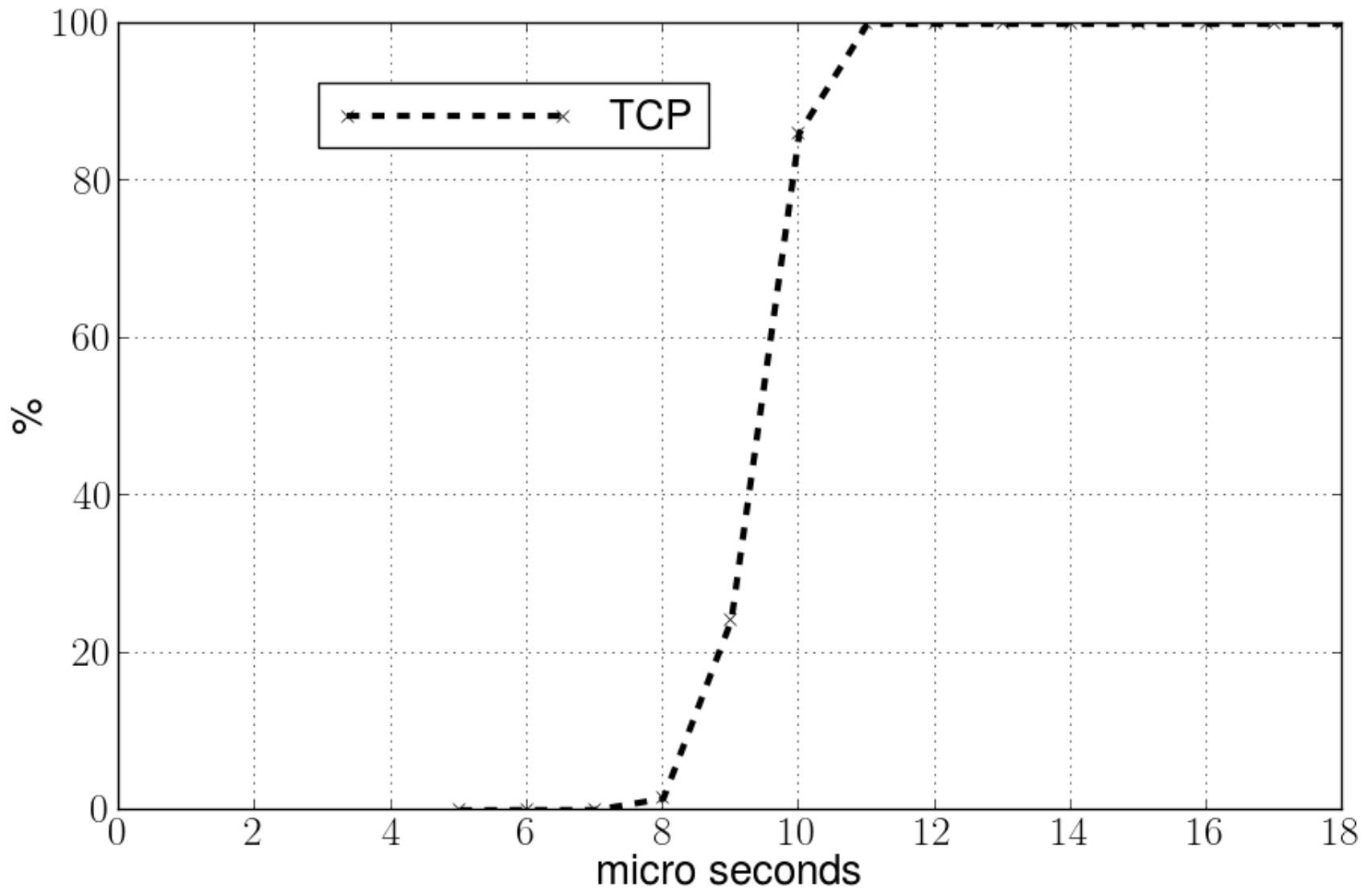
- No need to generate an HMAC
 - Allows stateless handling of SYN+JOIN
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Performance evaluation of MPTCP low overhead

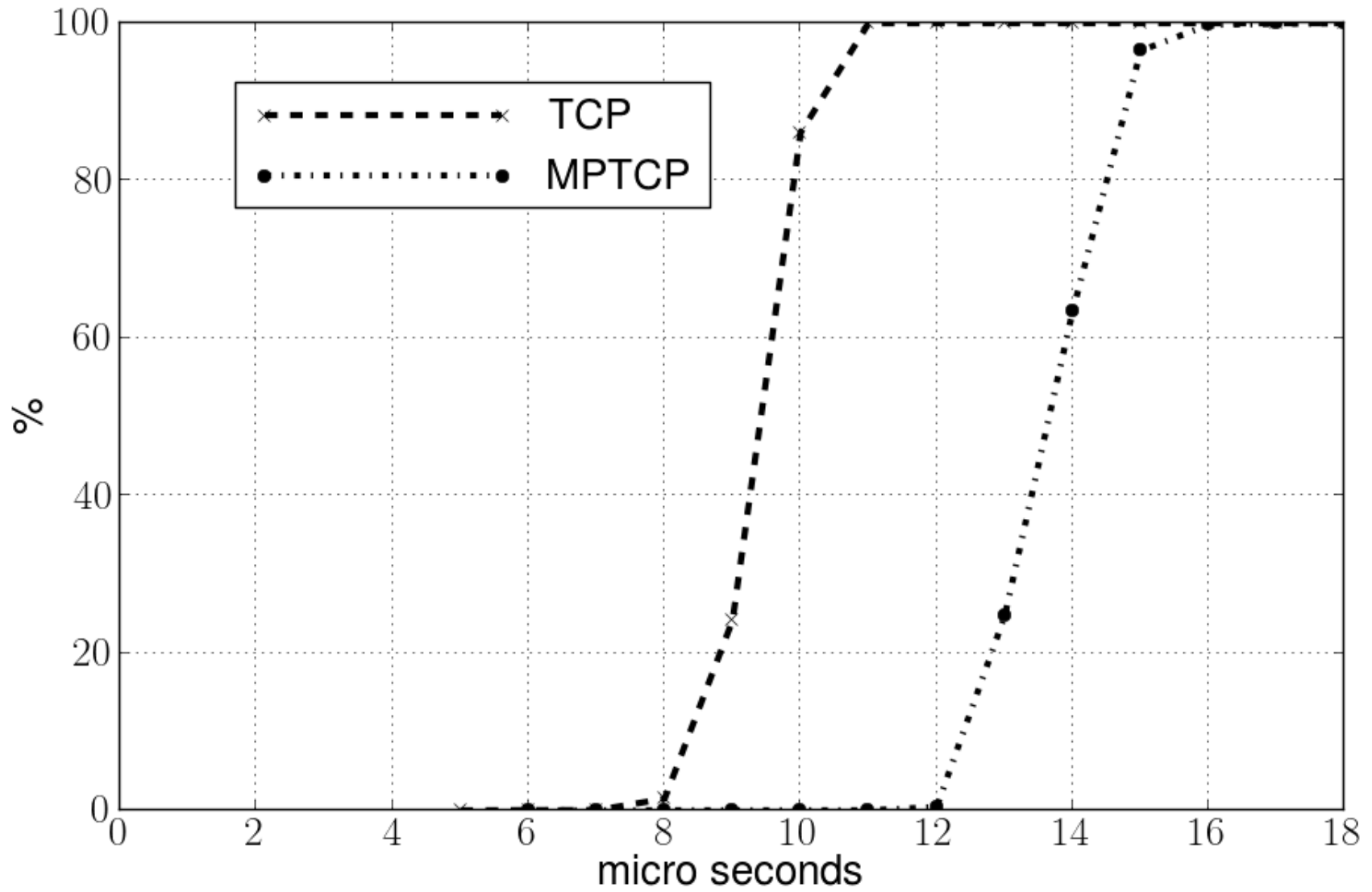
SYN/ACK with MP_CAPABLE



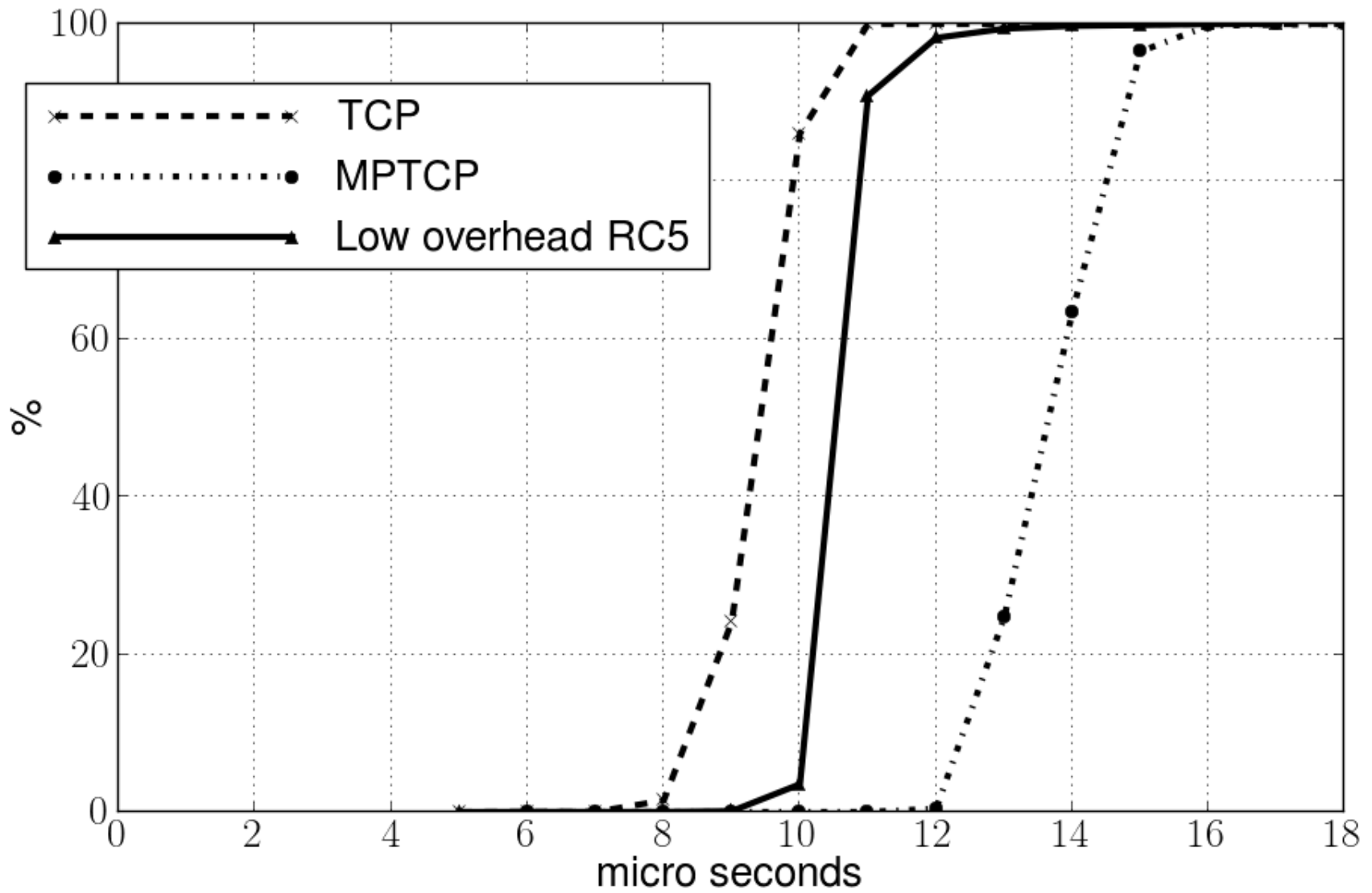
SYN/ACK with MP_CAPABLE



SYN/ACK with MP_CAPABLE



SYN/ACK with MP_CAPABLE



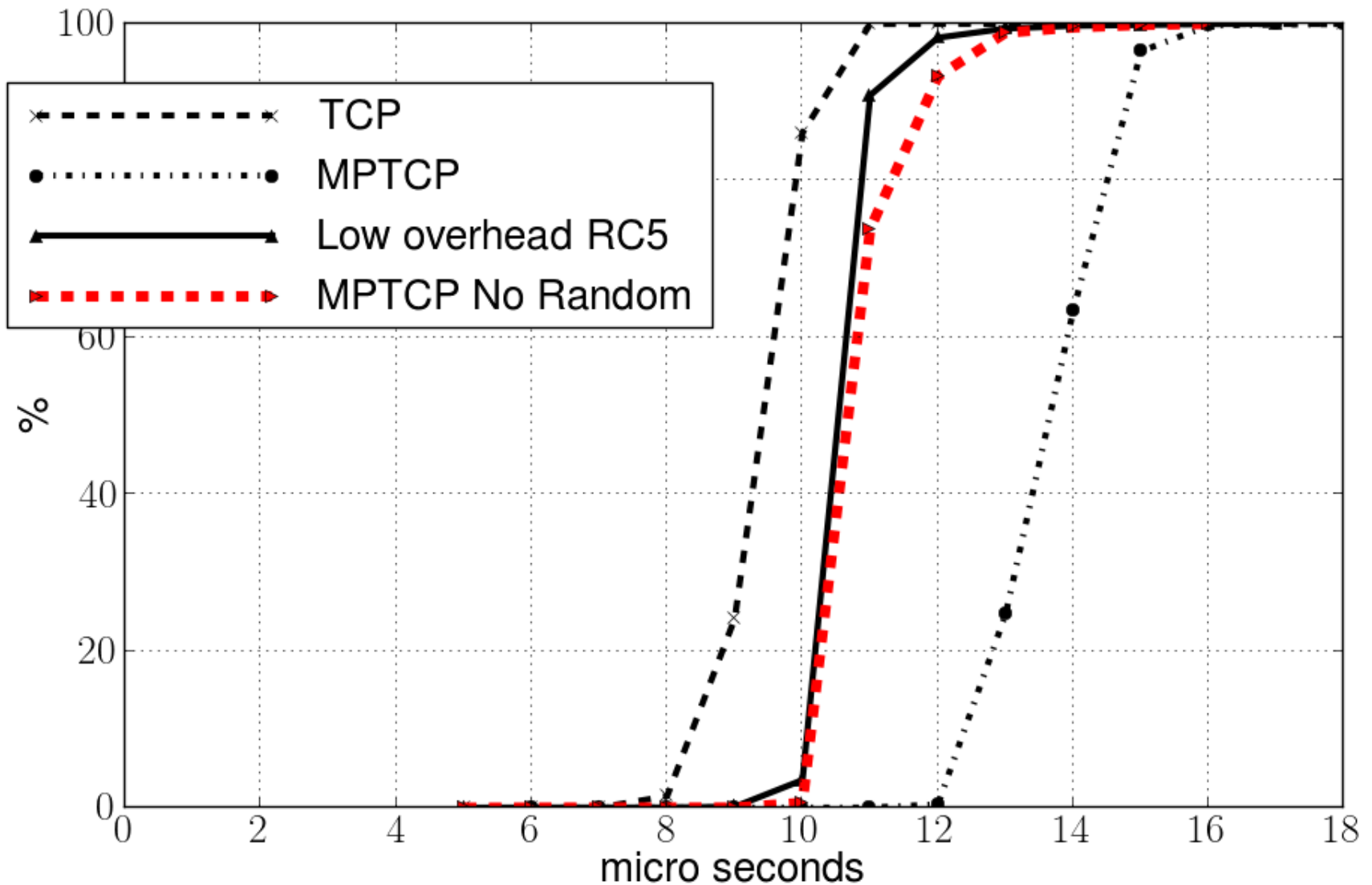
Random Number Generation

- MPTCP v0 generates random numbers:
 - Key (64-bits)
 - Random Number for MP_JOIN (32-bits)
 - Generating these random numbers accounts for ~20% of the time
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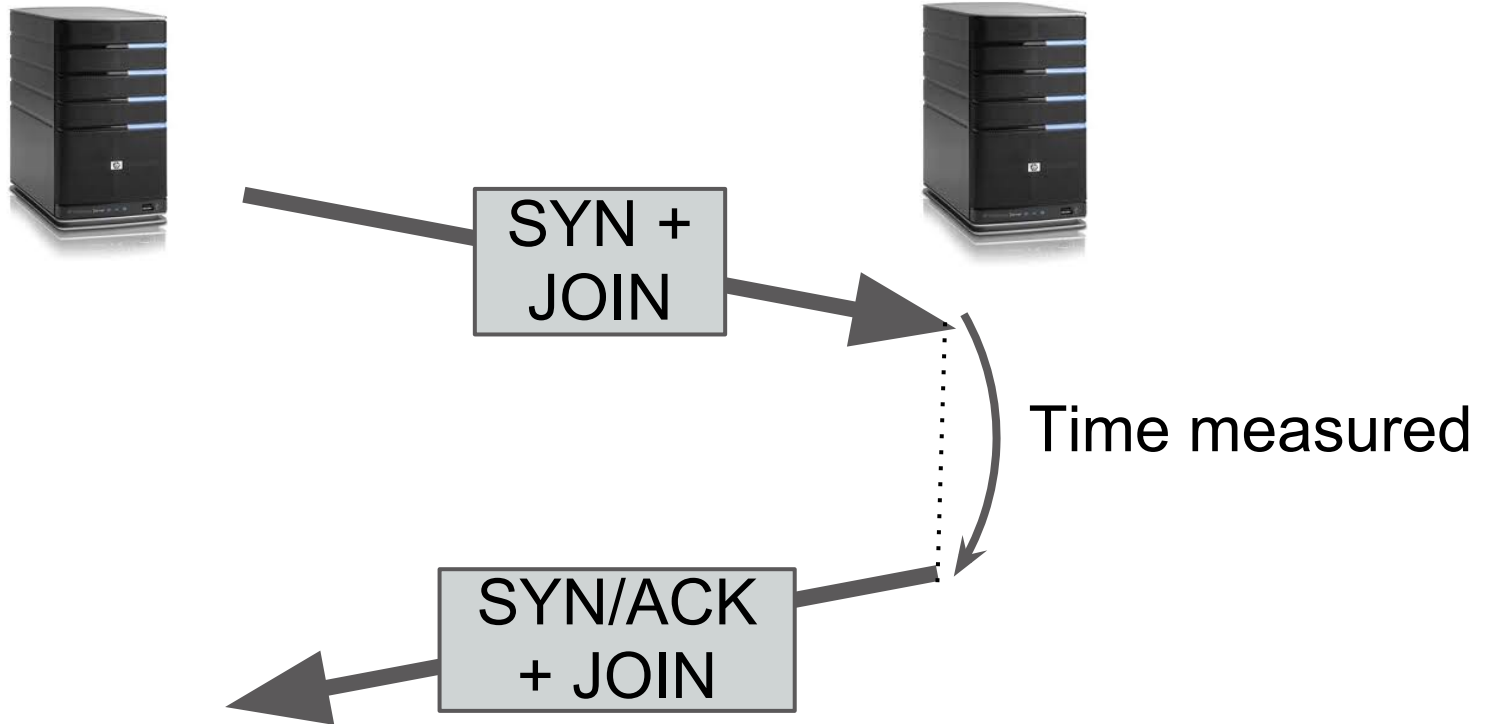
Random Number Generation

- RFC 6258: Generating the TCP sequence number:
 - $ISN = MD5(5\text{-tuple} + \text{secret})$
 - MPTCP Proposal:
 - $Key = Hash(5\text{-tuple} + \text{secret} + \text{counter})$
 - $Join\text{-Nonce} = Hash(5\text{-tuple} + \text{secret} + seqNo)$
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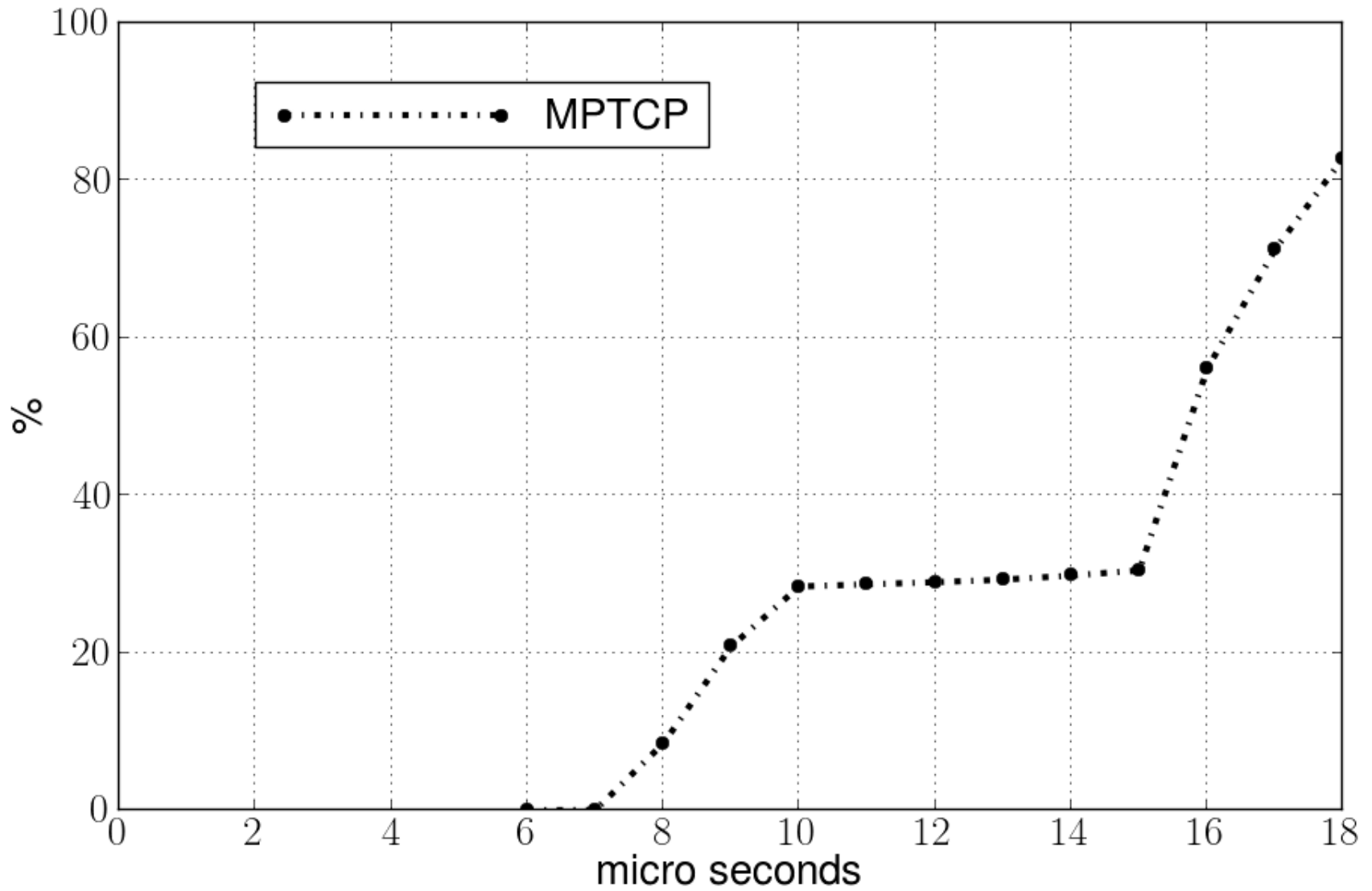
SYN/ACK with MP_CAPABLE



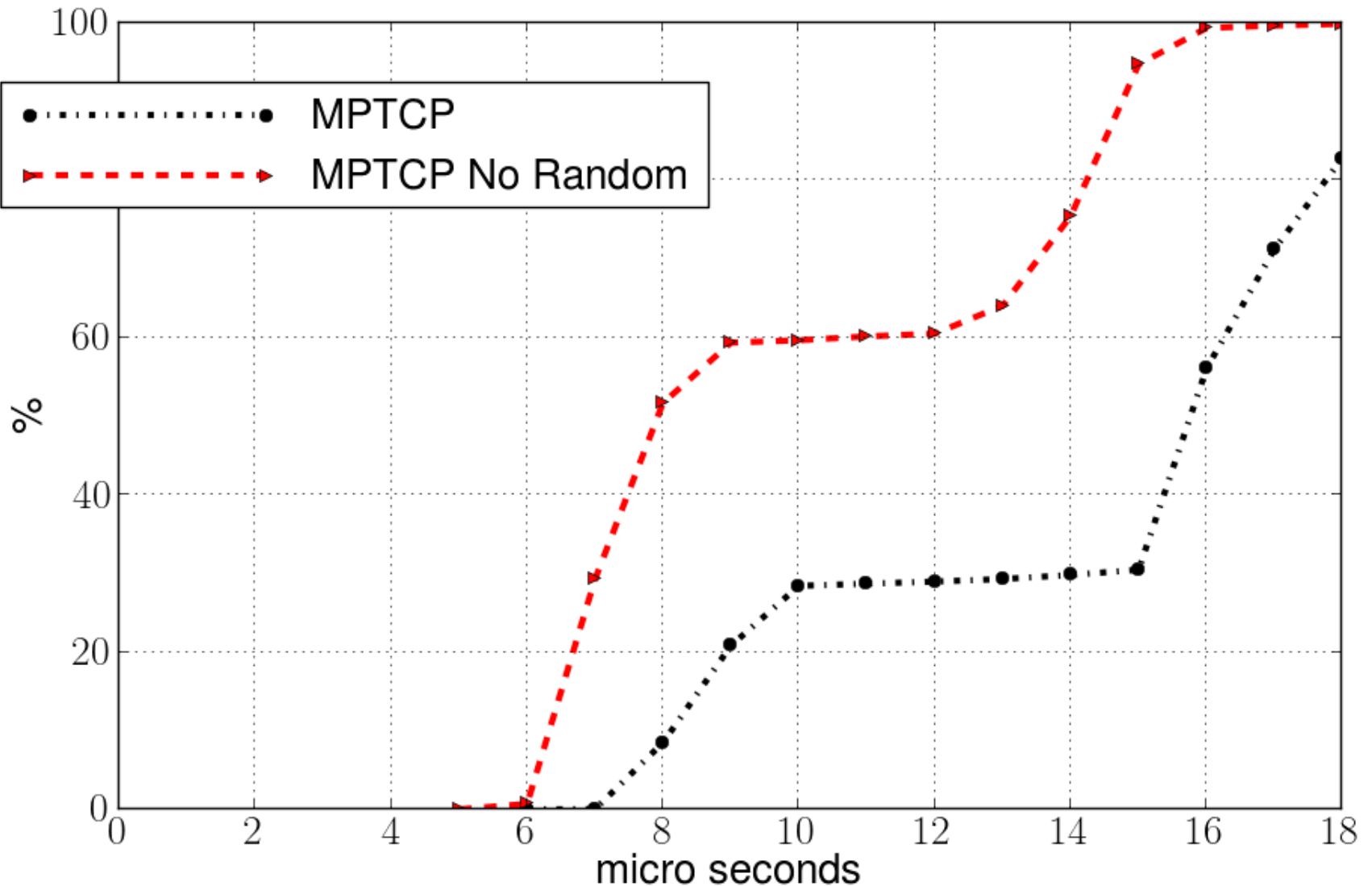
SYN/ACK with MP_JOIN



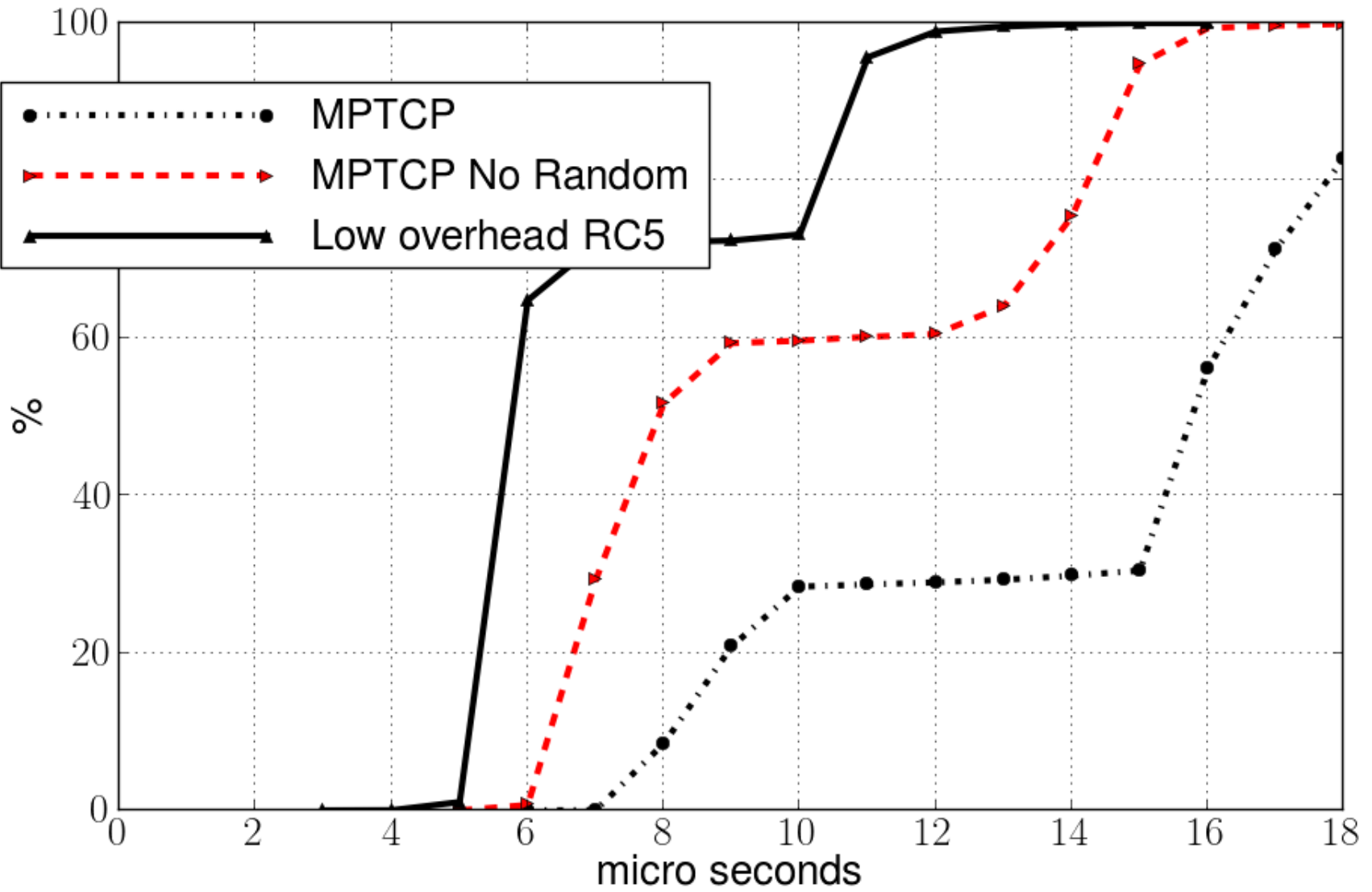
SYN/ACK with MP_JOIN



SYN/ACK with MP_JOIN



SYN/ACK with MP_JOIN

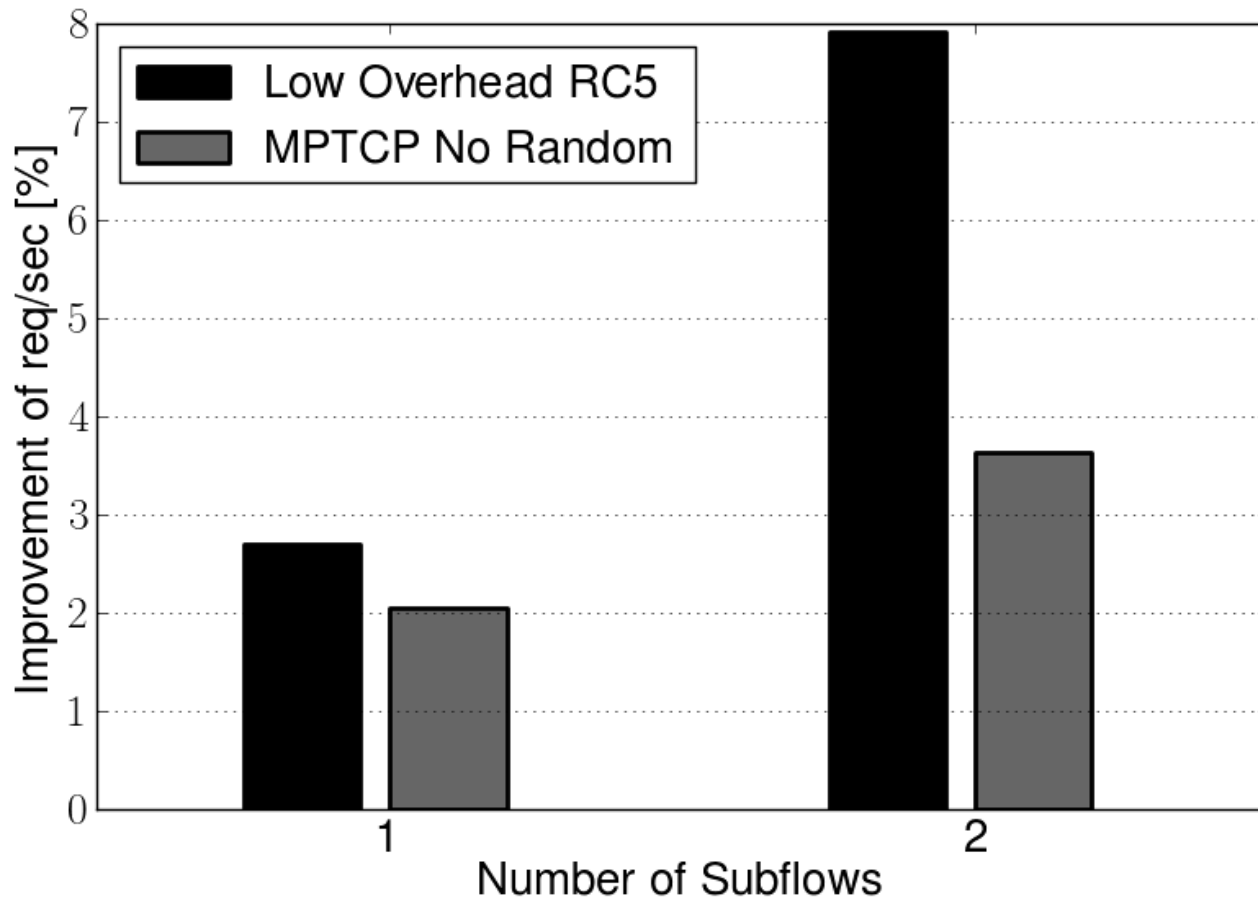


Impact on "real" traffic

- 250 simultaneous HTTP-requests
 - File-size: 1KB
 - Increasing number of MPTCP-subflows

 - Measuring Requests/Second
 - Baseline: Regular MPTCP
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Impact on "real" traffic



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Securing the MultiPath TCP handshake with external keys

Motivation

- MPTCP v0 sends the keys in clear
 - Attacker who sees the initial handshake can hijack an MPTCP session
 - Integrate application-level security into MPTCP to benefit from the app-security

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Goals:

- Do not send MPTCP's key in clear
- Application-level protocols already do negotiate a key (cfr. SSL)

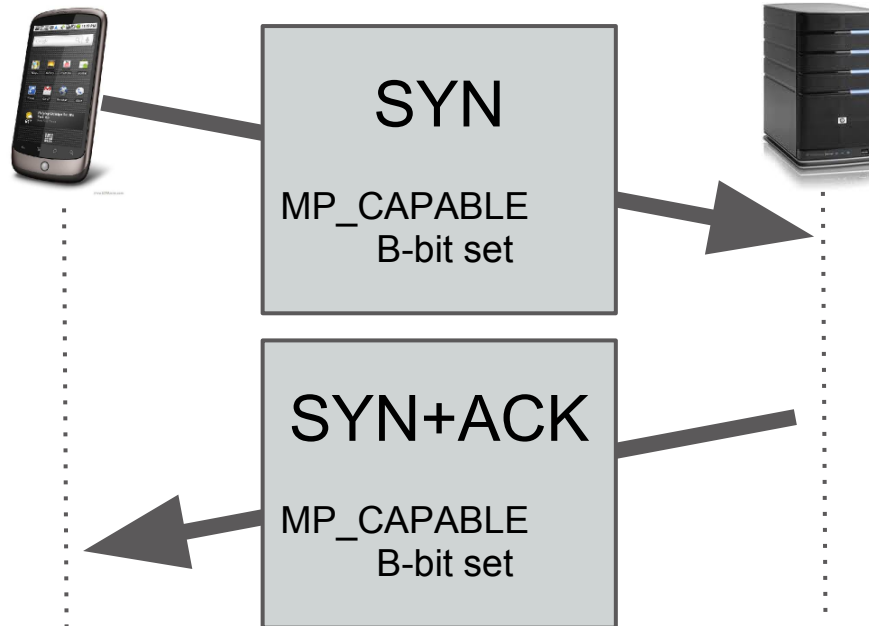
We should use these keys!

- Extend the socket-API to allow keys from the application
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SSL initial handshake

setsockopt(MPTCP_ENABLE_APP_KEY)

setsockopt(MPTCP_ENABLE_APP_KEY)



SSL initial handshake



SSL exchange

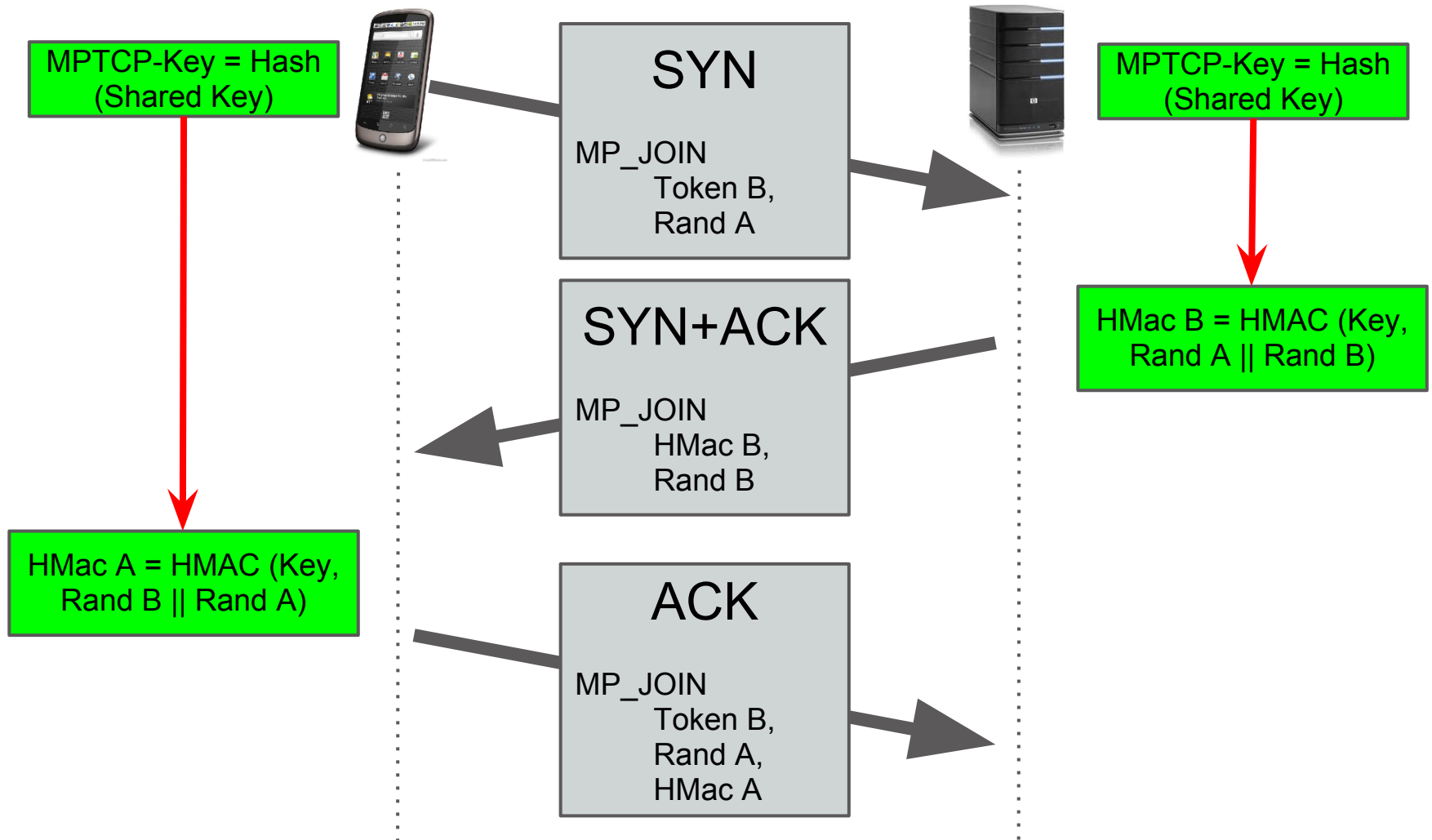
Hash(Shared Key)

Hash(Shared Key)

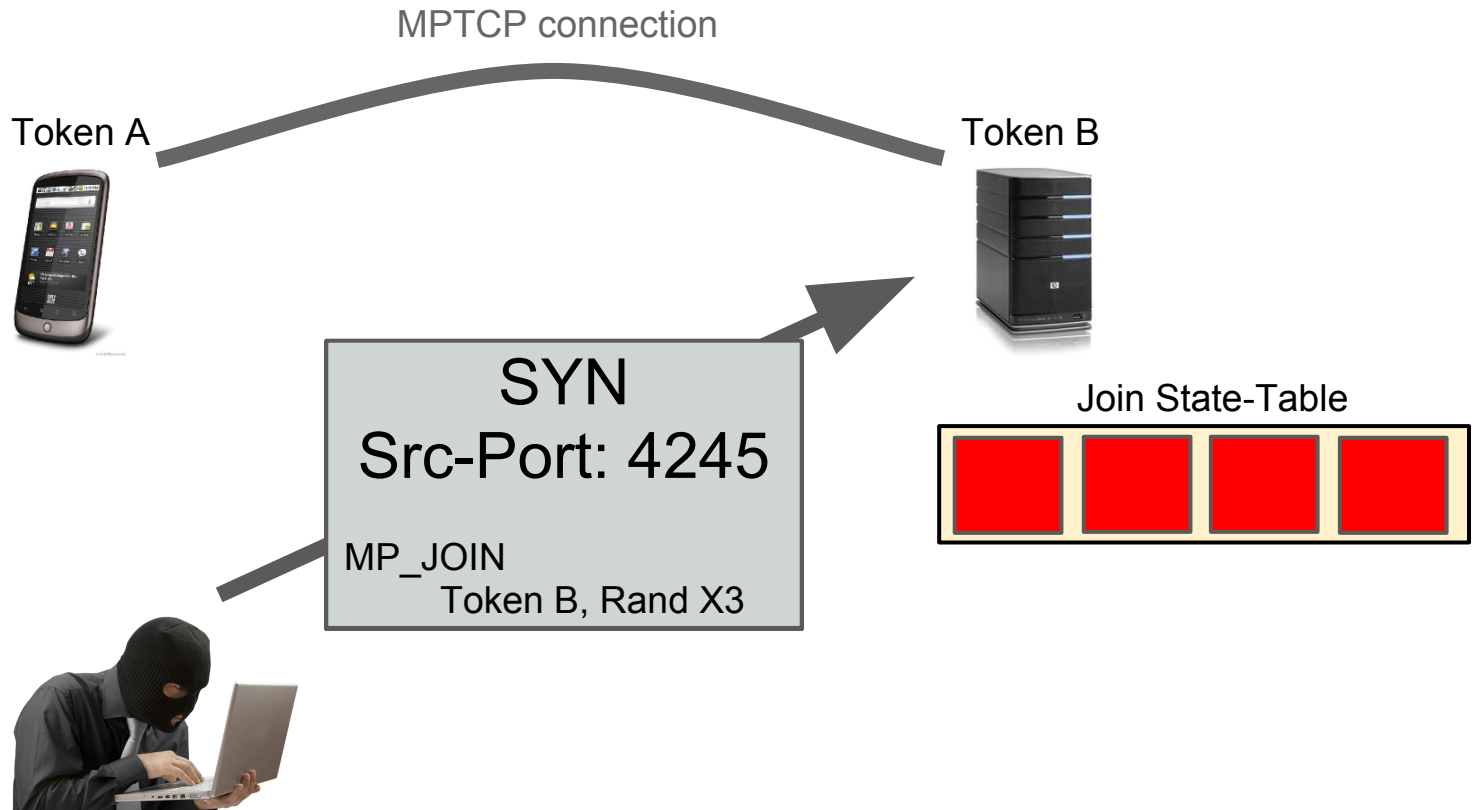
setsockopt(MPTCP_KEY)

setsockopt(MPTCP_KEY)

SSL additional subflow



JOIN-flooding Attack



Stateless JOIN-handling

- Prevents JOIN-flooding attacks
- Based on the mechanism of TCP SYN-Cookies

Stateless JOIN

Possible thanks to the modified JOIN-format:

MP_JOIN
HMAC - truncated to
128-bits
Random A
Token B

Server must generate his random number in a verifiable fashion:

$\text{Rand B} = \text{Hash}(5\text{-tuple} + \text{secret})$

Conclusion

MPTCP low overhead

- Performance improvement

MPTCP external keys

- Keys are no more exchanged in clear
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