BF-based chunk availability compression for PPSP-02

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Outline

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- Background
- Analysis
- Proposal
- Discussion
- Example

Motivation

- There are frequent bitmap exchanges in PPSP.
  - Uncompressed bitmap is relatively big (of several KBs).
  - They are exchanged frequently (less than several seconds).
  - It sets a limit to the system's efficiency and scalability
- There are efficiency requirements in PPSP PS
  - PPSP.TP.REQ-3: The tracker protocol MUST take the frequency of messages and efficient use of bandwidth into consideration, when communicating chunk availability information.
  - PPSP.PP.REQ-7: The peer protocol MUST take the frequency of messages and efficient use of bandwidth into consideration, when communicating chunk information.

Background

- Original bitmap Scheme uses a bit-array to represent the chunk set, and mark those bits corresponding to the locally available chunks.
- Con: Its length grows as the number of chunks (n) increases.
- Chunk range scheme uses a array of starting+ending chunk index pairs to represent continuous intervals.
- Byte range scheme uses an array of starting+ending byte index pairs to represent continuous intervals.
- Bin number scheme uses a universally assigned bin number (integer) for any given continuous interval.

Scheme type	Orig-bitmap	Chunk range	Byte-range	Bin number
Bit length	n	2klogn	2kmlogn	klogn

Variables: n - # of chunks; k - # of intervals; m - # of bytes per chunk.

A Further Look into Processing/Storage

Scheme type	Orig-bitmap	Chunk range	Byte-range	Bin number	???
Bit length	n	2klogn	2kmlogn	klogn	1
Formation	n	n	n	nlogn	1
Single Inquiry	1	2klogn	2kmlogn	logn	1
Group Inquiry(g)	g	2klogn	2kmlogn	glogn	1
Partial Update(p)	p	2klogn	2kmlogn	Tlogn	1
Delta bit len(d)	n	2d	2d	dlogn	1

Variables: n - # of chunks; k - # of intervals; m - # of bytes per chunk.

g - # of chunks in a compacted inquiry;

p/d - # of chunks in a partial chunk bitmap update.

Proposal: BF compression Scheme

```
-------------------------------------------------
BF(set S, integer m, hash set H)
1 filter=allocate m bits initialized to 0;
2 for each element xi in S do
3   for each hash functions hi in H do
4     filter[hi(xi)]=1;
5 return filter;
-------------------------------------------------
MT(element elm, BF filter, integar m, hash set H)
1 for each hash functions hi in H do
2   if (filter[hi(elm)]!=1)
3     return false;
4 return true;
-------------------------------------------------
ST(BF query, BF filter)
1 temp=query OR filter;
2 if (temp!=filter)
3     return false;
4 return true;
-------------------------------------------------
```

Figure 1 Basic algorithms for BF-bitmaps.

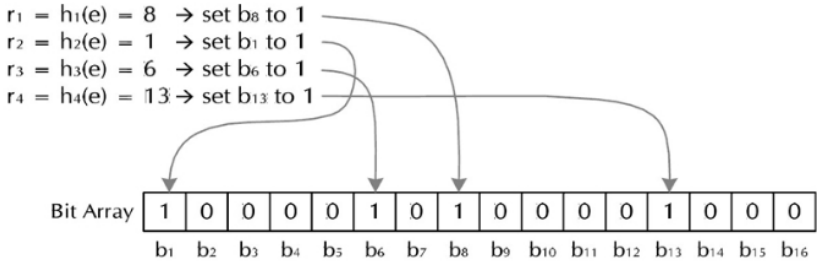


Figure 2 Bloom Filters: an example.

High Efficiency

Storage/transmission: Bit length: constant.

Processing: Formation/Inquiry/Update: constant.

Endurable Lose of accuracy

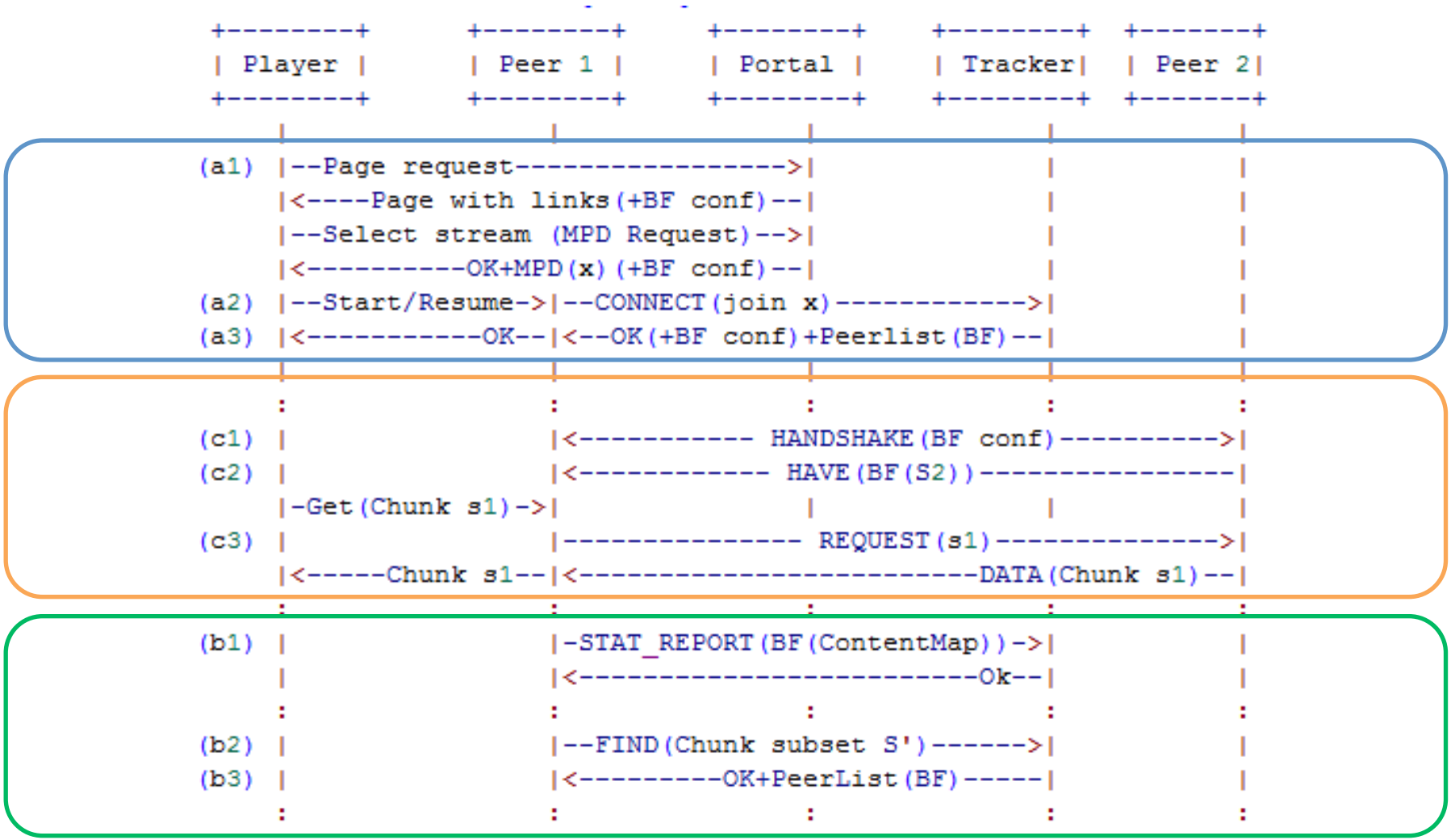
Be controlled by the system configuration of the bit array's length, choice and # of hash functions.

Example: a 2GB movie file, divided into 2MB chunks, whose a 1024-bit original chunk bitmap, can be represented by a 128-bit BF-bitmap (using 4 hashes), with only 3% mis-hits rate.

Suggestions for Discussion

- RECOMMENDED for PPSP-TP-base/extended
  - Strictly controllable cost for a central entity
    - constant bit-length irrelevant of the chunk-set
    - only replacement or simple bitwise operations needed
  - Certain mis-hits rate COULD be tolerable
    - Tracker serves as an initial broker for neighboring peers
- OPTIONAL for PPSPP
  - Peers willing to trade accuracy with cost-efficiency
    - Peers with limited computation/memory resources
    - Peers with huge number of concurrent links, e.g. SNs
  - Certain mis-hits rate MAY be tolerable
    - REQUEST and DATA SHOULD use the original chunk id.

Example: PPSP with BF-bitmaps



THANK YOU!