# BATS: An efficient network coding solution for packet loss networks

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#### March 6, 2014 Presented at the IRTF Network Coding Research Group

Joint work with Shenghao Yang (IIIS, Tsinghua U)



#### Patents Related to BATS code

- 1 US patent pending (US Patent App. 13/112,589)
- PCT application in China and some European countries



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Intermediate	End-to-End	Maximum Rate		
forwarding	retransmission	0.64		
forwarding	fountain codes	0.64		



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forwarding	retransmission	0.64
forwarding	fountain codes	0.64
network coding	random linear codes	0.8

#### Achievable Rates: *n* hops



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	Intermediate Operation	Maximum Rate	
-	forwarding	$0.8^n \rightarrow 0, \ n \rightarrow \infty$	
	network coding	0.8	

#### Complexity of Linear Network Coding

- T: length of a packet; K: number of packets
- Encoding:  $\mathcal{O}(TK)$  per packet.
- Decoding:  $\mathcal{O}(K^2 + TK)$  per packet.
- Network coding:  $\mathcal{O}(TK)$  per packet. Buffer K packets.



### Batched Sparse (BATS) Codes



[YY11] S. Yang and R. W. Yeung. Coding for a network coded fountain. ISIT 2011, Saint Petersburg, Russia, 2011.

#### Encoding of BATS Code: Outer Code

- Apply a "matrix fountain code" at the source node:
  - **1** Obtain a degree d by sampling a degree distribution  $\Psi$ .
  - Pick d distinct input packets randomly.
  - **③** Generate a batch of *M* coded packets using the *d* packets.
- Transmit the batches sequentially.



 $\mathbf{X}_i = \begin{bmatrix} b_{i1} & b_{i2} & \cdots & b_{id_i} \end{bmatrix} \mathbf{G}_i = \mathbf{B}_i \mathbf{G}_i.$ 

- The batches traverse the network.
- Encoding at the intermediate nodes forms the inner code.
- Linear network coding is applied in a causal manner within a batch.



#### Belief Propagation Decoding

- Find a check node *i* with degree<sub>*i*</sub> = rank( $\mathbf{G}_i \mathbf{H}_i$ ).
- 2 Decode the *i*th batch.
- Opdate the decoding graph. Repeat 1).



The linear equation associated with a check node:  $\mathbf{Y}_i = \mathbf{B}_i \mathbf{G}_i \mathbf{H}_i$ .

#### Precoding

- Precoding by a fixed-rate erasure correction code.
- The BATS code recovers  $(1 \eta)$  of its input packets.



[Shokr06] A. Shokrollahi, Raptor codes, IEEE Trans. Inform. Theory, vol. 52, no. 6, pp. 25512567, Jun. 2006.

Source node	encoding	$\mathcal{O}(TM)$ per packet	
Destination node decoding		$\mathcal{O}(M^2 + TM)$ per packet	
Intermediate Node	buffer	$\mathcal{O}(TM)$	
	network coding	$\mathcal{O}(TM)$ per packet	

- T: length of a packet
- K: number of packets
- M: batch size

#### BATS codes with M = 32 and q = 256.

ĸ	coding overhead		inactivation no.			
Λ	average	max	min	average	max	min
1600	2.04	16	0	94.0	119	72
8000	6.30	77	0	215.5	268	179
16000	26.58	1089	0	352.2	379	302

- M = 1: BATS codes degenerate to Raptor codes.
  - Low complexity
  - No benefit of network coding
- M = K and degree  $\equiv K$ : BATS codes becomes RLNC.
  - High complexity
  - Full benefit of network coding.
- Exist parameters with moderate values that give very good performance



- Packet loss rate 0.2.
- Node *s* encodes *K* packets using a BATS code.
- Node *u* caches only one batch.
- Node *t* sends one feedback after decoding.

## Experiment setting



## Experiment setting



- Sender/receiver: a laptop with open source Atheros wireless drivers.
- Relay: one wireless router with Atheros chipset running OpenWrt (about 150HKD/20USD)
- WiFi 802.11 b/g/n at 2.4GHz
- Sender's rate is set to 1 Mb/s to reduce the effect of the router's low computation power.

	Average rate (Kb/s)
BATS w/ recoding	592.86
BATS w/o recoding	530.65
TCP (normal 802.11)	420.33

#### Application: vehicular ad-hoc network



#### Application: mobile ad-hoc network



- BATS codes provide a digital fountain solution with linear network coding:
  - Outer code at the source node is a matrix fountain code.
  - Linear network coding at the intermediate nodes forms the inner code.
  - Prevents BOTH packet loss and delay from accumulating along the way.
- The more hops between the source node and the sink node, the larger the benefit.