ALTO Extension: Routing State Abstraction using Declarative Equivalence

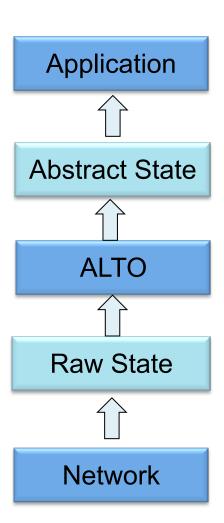
draft-gao-routing-state-abstraction-00 draft-yang-alto-topology-06

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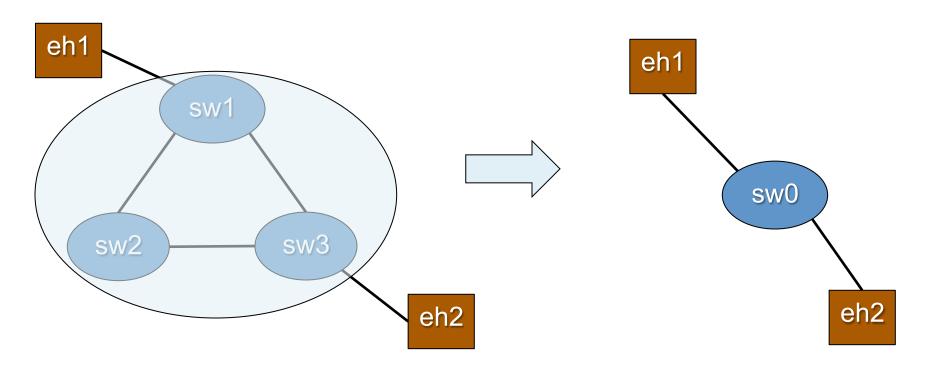
Bigger Picture

- A general objective of ALTO is to provide network state to applications for better traffic engineering
- It is important that ALTO provide abstract network states, to
 - protect information privacy
 - improve scalability

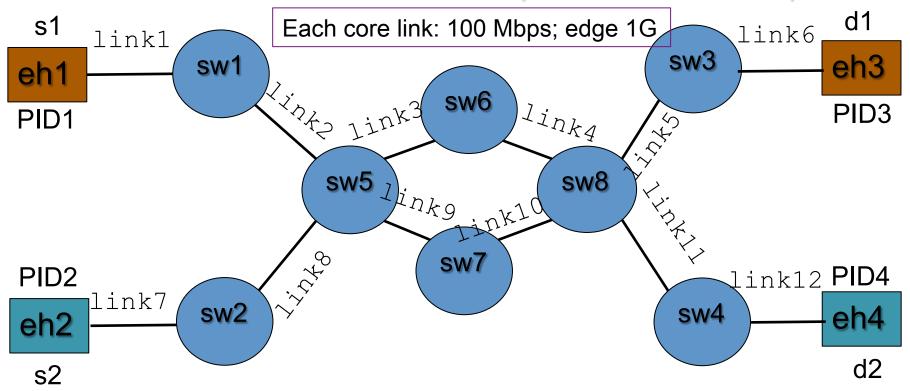


Key Question: How to Compute Abstract State

Basic approach: static template (e.g., single-node template)



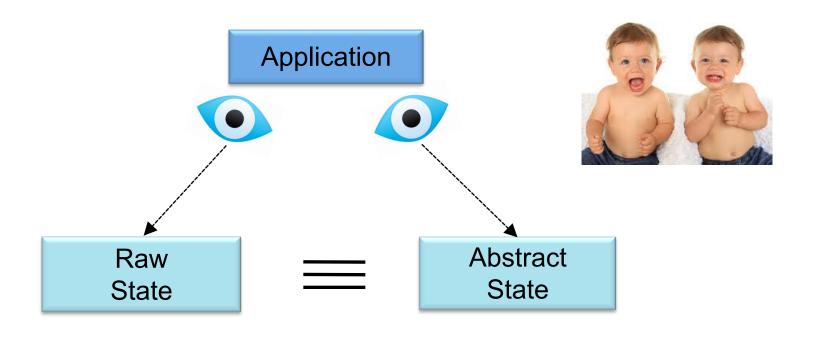
What we learned from the path-vector example



- App requests available bandwidth for two pairs:
 - PID1 (eh1) -> PID3 (eh3); PID2 (eh2) -> PID4 (eh4)
- Ambiguity of cost map based on single-switch abstraction:
 - Two disjoint paths (200 Mbps when concurrent), e.g.,
 - PID1 -> PID3: sw5 -> sw6 -> sw8;
 - PID2 -> PID4: sw5 -> sw7 -> sw8
 - Shared bottleneck (still 100 Mbps when concurrent)

Our General Design Approach

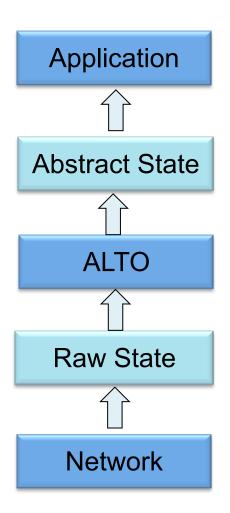
- ALTO server computes dynamic, minimal network state
 - Compute abstract state that is smaller but equivalent to application required



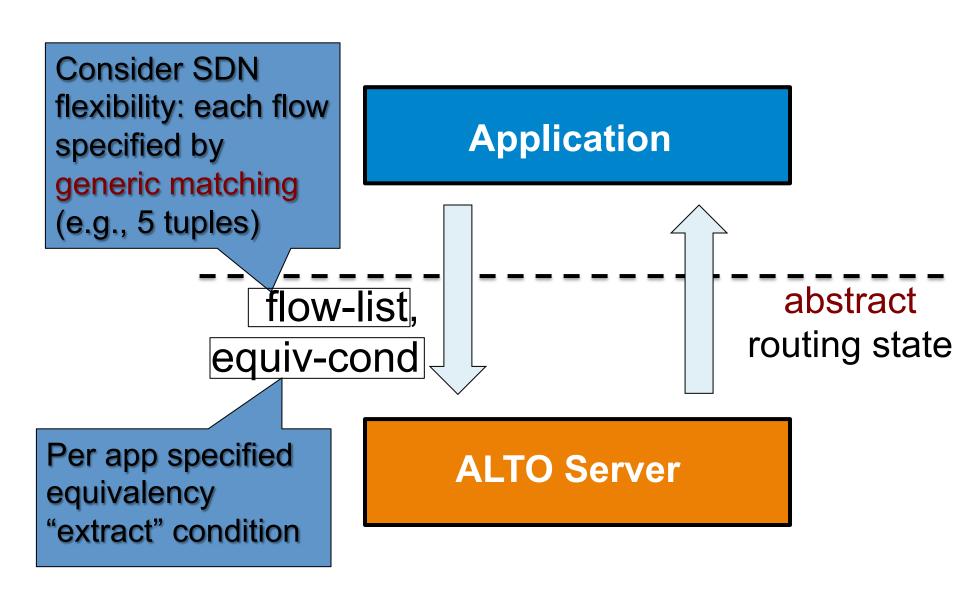
ALTO Extension: Routing State Abstraction Service based on Declarative Equivalence (RSA-DE)

- Why routing state
 - Routing state is basic
 - Current focus of ALTO (e.g., cost map,
 ECS) is mostly on routing state

 We may consider other types of network state in future extensions



Routing State Abstraction Service



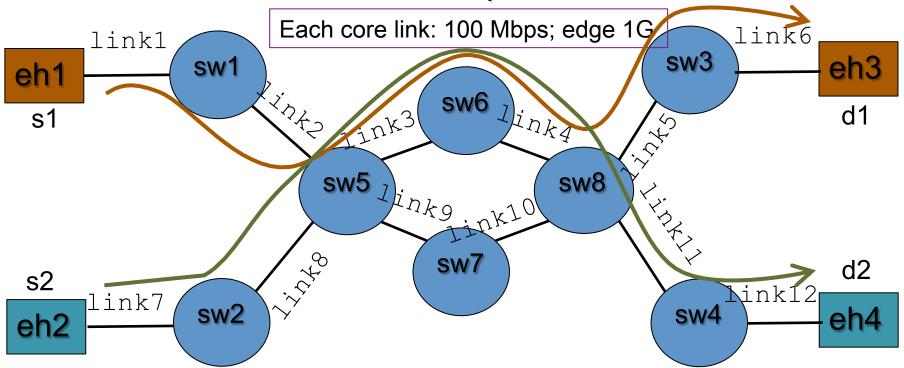
How to Specify flow-list

- This is relatively straightforward
 - Server announces matching capabilities in IRD
 - E.g., IP address only, allowing ports, ...
 - Incremental deployment: allows IP only
 - A flow-list is a list of flows, where each flow is specified by a matching condition (e.g., OpenFlow like) condition

How to Specify equiv-cond: Intuition

- General structure of a network application
 - Has a set of flows
 - Has a set x of variables (e.g., rate of each flow)
 - Has a goal: min/satisfy obj(x)s.t.x satisfies constraints
 - obj(x) or the constraints may use network information, e.g.,
 - obj = x[1] * routingcost(f1) + x[2] * routingcost(f2)
 - any link e: sum(x[i]: fi uses link e) ≤ bw(e)
- equiv-cond is to convey the usage of network information in the application

Example



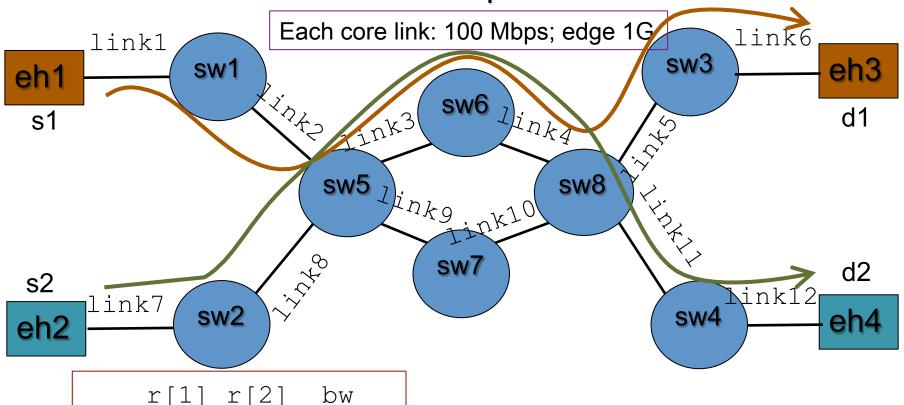
raw state

r[1]	r[2]	bw
link1	1	0	1G
link2	1	0	100M
link3	1	1	100M
Link4	1	1	100M
link5	1	0	100M
link6	1	0	1G
•••			

Given routing for given flows, the network raw state can be considered as a set of vectors, where each vector's dimension is the number of edges (links):

- r[i][e]: a vector representing if flow i uses link
 e. For example, r[2] shows that the route of flow 2 uses links 3/4, ..., not links 1/2/5/6, ...
- attr[e]: is the value of attr for link e. For example, bw[e] is the available bandwidth of link e.





1G

100M

100M

100M

100M

1G

0

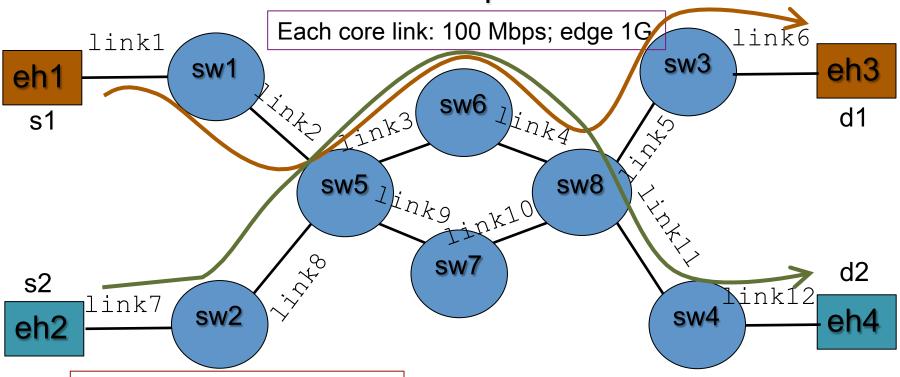
raw state | link2 | 1 | 0 | 1 | 1 | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | 1 | | | 1 | | | | 1 | | | 1 | |

link1

link6

A naïve approach is to return the whole raw network state





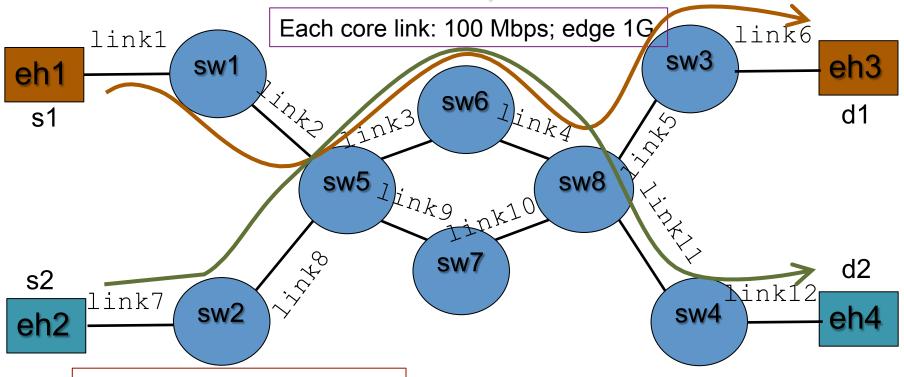
raw state

```
r[1]
          r[2]
                 bw
link1
                 1G
                 100M
link2
                 100M
link3
Link4
                 100M
link5
                 100M
link6
              0
                 1G
```

App decision variables: x[1], x[2].

equiv-cond: Lambda specification of app usage of network state in constraints:





raw state

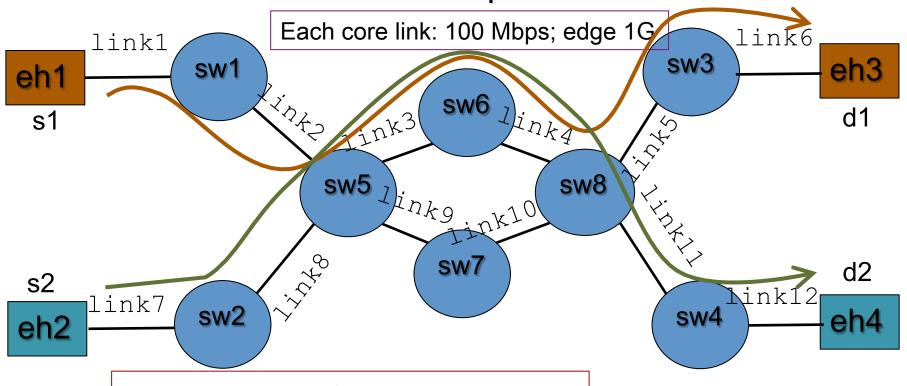
r[1]	r[2]	bw
link1	1	0	1G
link2	1	0	100M
link3	1	1	100M
Link4	1	1	100M
link5	1	0	100M
link6	1	0	1G
•••			

Insight: If the attributes of a link e do not appear in a tight (independent) constraint, the link does not need to be known to the app.

Implementation

- ALTO will not define the implementation, the discussion of implementation is to show feasibility and help with understanding
- Steps to implement RSA-DE
 - ALTO server looks up routing for each flow i to obtain r[i]
 - E.g., looks up in Flow Rule Manager (FRM) in ODL/ONOS
 - ALTO server applies redundancy elimination to find the minimal set of independent links

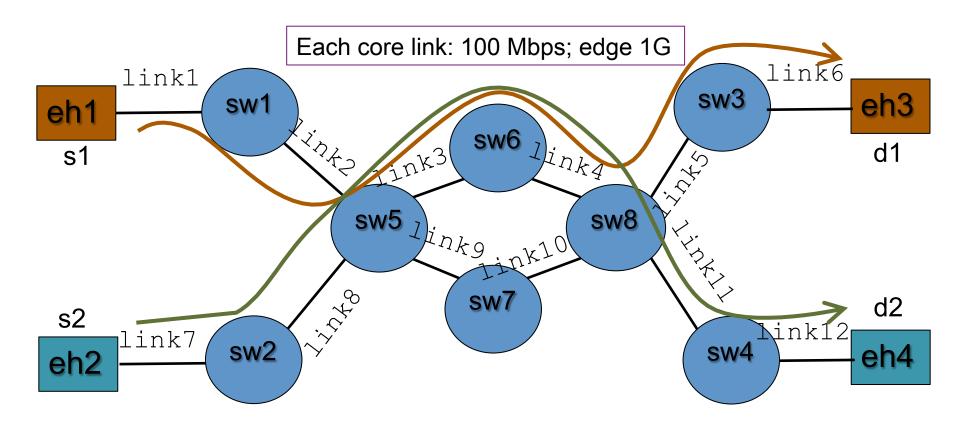




	r[1]	r[2]	var	bw
link1	1	0	x[1]	1G
link2	1	0	x[2]	100M
link3	1	1		100M
Link4	1	1		100M
link5	1	0		100M
link6	1	0		1G
•••				

Fast (redundancy elimination) algorithm to compute links providing non-redundant constraints.

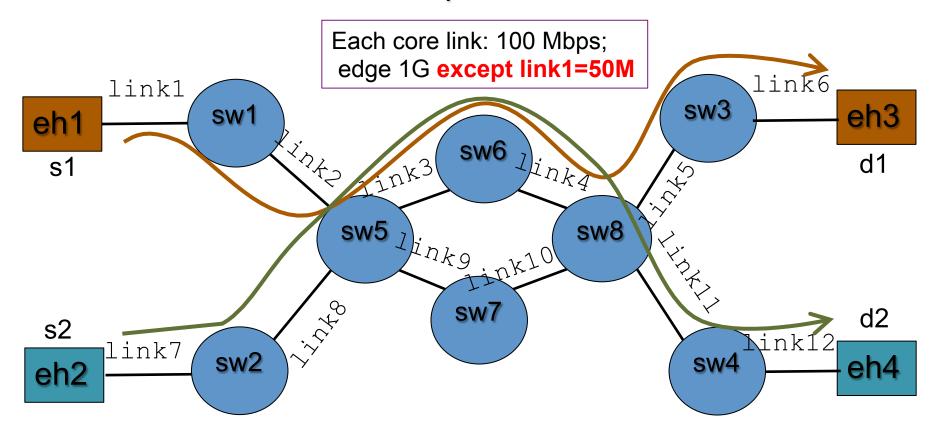
Example: Result



```
s_1 \rightarrow d_1: { ane<sub>1</sub>(\leq 100M) }

s_2 \rightarrow d_2: { ane<sub>1</sub>(\leq 100M) }
```

Example: Result



```
s_1 -> d_1: { ane<sub>1</sub>(\leq 100M), ane<sub>2</sub>(\leq 50M) }
s_2 -> d_2: { ane<sub>1</sub>(\leq 100M)}
```

Summary

- RSA-DE provides a powerful, general interface to allow applications to obtain network state
 - It is a generalization of the previous path-vector design
 - It works in the new SDN setting
 - More details on RSA-DE see backup slides

 Interest in the WG to pursue this direction as an ALTO extension?

Backup Slides

How to Specify equiv-cond

- Link Properties: announced in IRD capabilities
 - r[i] must be supported
 - bw, delay, loss, routingcost, ...
- equiv-cond
 - Variables: A list of opaque variables
 - Constraints: A list of lambda inequality expression for a given link, e.g., any e in E:

```
r[1][e]*x[1] + r[2][e]*x[2] <= bw[e]
```

- To simplify the representation:
 - [e] can be ignored since it is implied.
 - Allow meaningful variable names
 - Use underscore to represent array instead of brackets

- The order to parse a constraint
 - Symbols in the variable-list:
 - As variables
 - Symbols announced as capabilities
 - As link-property-names
 - Symbols for mathematical constants
 - As constants
 - Unknown symbols
 - As undefined if the variable-list is provided
 - As variables if intelligent parsing is enabled

```
variable-list
variable := variable | variable, variable-list
variable := [a-zA-Z][0-9a-zA-Z-]*

rs-query := flow-list constraint-list [variable-list]
flow-list := flow | flow, flow-list
flow := generic-match-condition
```

```
"query": {
    "flows": [
            "dst": "ipv4:192.168.1.22",
            "src": "ipv4:192.168.0.11",
            "tcp-dstport": "80"
        },
            "dst": "ipv4:192.168.3.44",
            "src": "ipv4:192.168.2.33",
            "tcp-dstport": "23"
    "conditions": [
        "r 0*f0 + r 1*f1 \le bandwidth",
        "r 0*f0 + r 1*f1 <= u*capacity"
    "variables": [
        "f0",
        "f1",
        "11"
```

An example where the user creates two flows with three variables.

- Variables £0 and £1 represent the bandwidth for each flow.
- Variable u represents the maximum link utilization.

Routing State Abstraction Service: Query Input

```
rs-query := flow-list equiv-cond
flow-list := flow [flow-list]
flow := generic-match-condition
```

- App provides two input parameters
 - flow-list: A set of flows
 - Consider the flexibility of SDN, which allows routing based on generic matching (e.g., 5 tuples) => each flow is specified by generic matching condition
 - equiv-cond: app declared equivalency condition

Path Vector: Example

```
HTTP/1.1 200 OK
Content-Length: TDB
Content-Type:
application/alto-costmap+json
{ "meta" : {
   "dependent-vtags" : [
    { "resource-id": "my-default-network-map",
      "tag": "3ee2cb7e8d63d9fab71b9b34cbf764436315542e" },
     {"resource-id": "my-topology-map",
      "tag": "4xee2cb7e8d63d9fab71b9b34cbf76443631554de"
   "cost-type": {"cost-metric": "bw", "cost-mode": "path-vector"},
   "cost-map" : {
     "PID1": { "PID1":[], "PID2":["ne56", "ne67"], "PID3":[], "PID4":["ne57"]
      "PID2": { "PID1":["ne75"], "PID2":[], "PID3":["ne75"], "PID4":[]
      }, ...
```