

Enhanced mobility anchoring wt

Teleconferences since IETF92

- ◆ May 22 9:30-10:30 CDT
- ◆ May 27 9:30-10:30 CDT
- ◆ Attended: Jong-Hyouk Lee, Seil Jeon, Carlos Bernardos, Fred Templin, H Anthony Chan

- ◆ draft-yhkim-dmm-enhanced-anchoring
 - AF split into AF-CP and AF-DP
- ◆ draft-chan-dmm-distributed-mobility-anchoring-02
 - session → flow
- ◆ draft-chan-dmm-distributed-mobility-anchoring-03
 - relate to examples of existing proposed solutions
 - IP address/prefix is anchoring to ...
- ◆ draft-chan-dmm-distributed-mobility-anchoring-04
 - improved security discussion
 - clarified FM-DP being distributed, FM-CP, LM each being distributed or centralized

Background: Functions involved in anchoring

- ◆ The following functions seem essential to anchoring:
 - Allocation of IP prefix/addr and Route advertisement.
- ◆ Other functions of anchor in different proposals: each function is not necessarily present in all the solutions:
 - (1) packets to/from the MN traverse through.
 - (2) indirection, e.g., set up tunnel
 - (3) information, e.g., binding HoA and CoA
 - (4) sends route update, e.g., using BGP

Description of IP address/prefix anchoring

- ◆ An IP address, i.e., Home Address (HoA), or prefix, i.e., Home Network Prefix (HNP) allocated to a mobile node is topologically anchored to a node when the anchor node is able to advertise a connected route into the routing infrastructure for the allocated IP prefix.

LM and FM in RFC7429

- ◆ Internetwork Location Management (LM) function: managing and keeping track of the internetwork location of an MN. The location information may be a binding of the IP advertised address/prefix, e.g., HoA or HNP, to the IP routing address of the MN or of a node that can forward packets destined to the MN. It is a control plane function.
 - LM may be a logical function at CPA and CPN
- ◆ Forwarding Management (FM) function: packet interception and forwarding to/from the IP address/prefix assigned to the MN, based on the internetwork location information, either to the destination or to some other network element that knows how to forward the packets to their destination.
 - FM-CP may be a logical function at CPA and CPN
 - FM-DP may be a logical function at DPA and DPN
- ◆ FM-DP may be distributed in distributed mobility management
- ◆ LM, FM-CP may each be centralized or distributed

Functions involved in anchoring

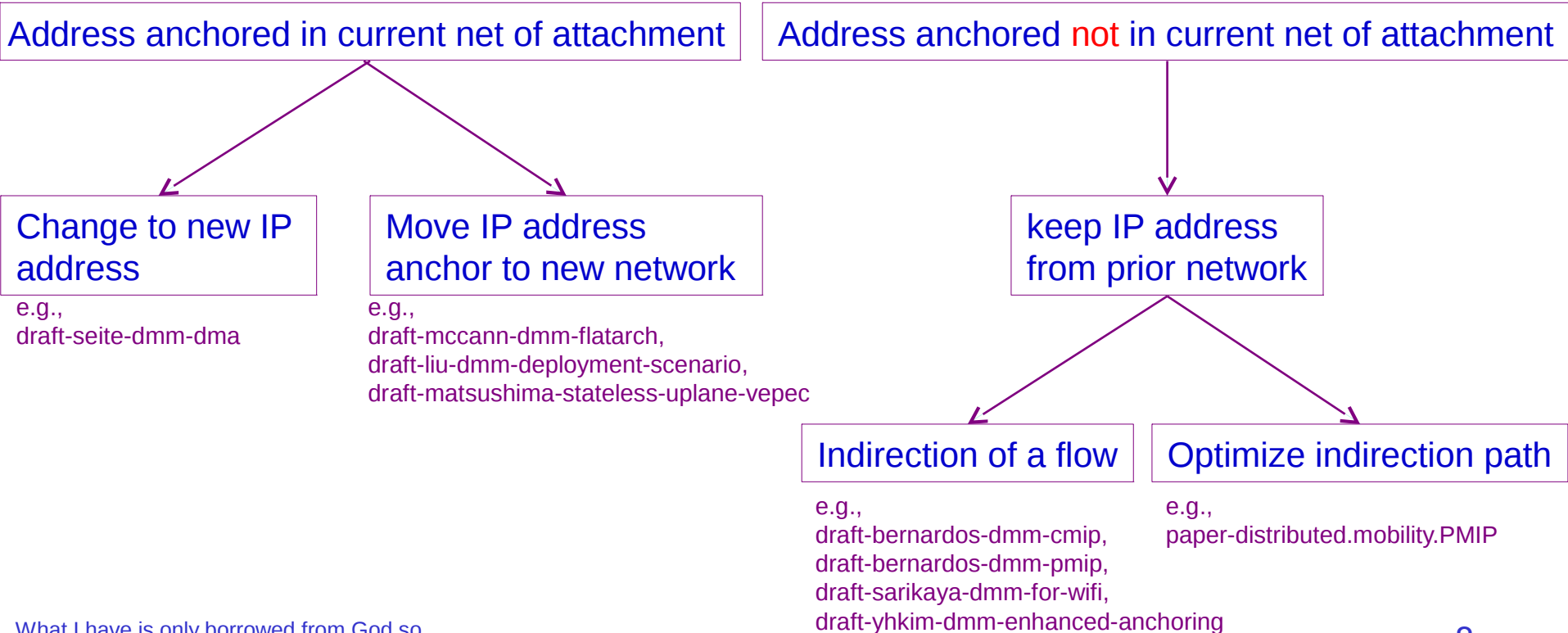
- ◆ The following functions seem essential to anchoring:
 - Allocation of IP prefix/addr and Route advertisement.

- ◆ Other functions of anchor in different proposals: each function is not necessarily present in all the solutions:
 - (1) packets to/from the MN traverse through.(DPA)
 - (2) indirection, e.g., FM-CP at CPA and CPN to set up tunnel between DPA and DPN
 - (3) information, e.g., binding HoA and CoA (LM at CPA and CPN)
 - (4) sends route update, e.g., using BGP

Example solutions

- ◆ draft-bernardos-dmm-cmip
- ◆ draft-bernardos-dmm-pmip
- ◆ draft-liu-dmm-deployment-scenario
- ◆ draft-matsushima-stateless-uplan-vepc
- ◆ draft-mccann-dmm-flatarch
- ◆ draft-sarikaya-dmm-for-wifi
- ◆ draft-seite-dmm-dma
- ◆ draft-yhkim-dmm-enhanced-anchoring
- ◆ paper-distributed.mobility.PMIP

Categorization (unification) of different approaches



Address anchored (to AR2) in current network of attachment

Before moving:

MN in Net1 is allocated IP1 from AR1

After MN moves to Net2

MN in Net2 is allocated IP2 from AR2

MN starts a flow(IP2,IPCN, ...) .

IP2 is in current network of attachment

Net1

Net2

AR1:
⚓ IP1

AR2:
⚓ IP2
MN(IP2): flow(IP2,...)

Changing to new IP address address anchored (to AR2) in current network of attachment

Before moving:

MN in Net1 is allocated IP1 from AR1

MN and CN has flow using IP1 and IPCN,

After MN moves to Net2

MN in Net2 is allocated IP2 from AR2

MN and CN restarts a new flow(IP2,IPCN, ...) .

IP2 is in current network of attachment.

Data Plane Packet in flow(IPCN,IP2,...)
from CN to MN:

1. Packet from CN → AR3
2. Forward from AR3 → AR2
3. Forwarded from AR2 → MN

Net1

AR3:
⚓ IPCN
CN: flow(IPCN, IP2, ...)

②

Net2

AR1:
⚓ IP1
MN(IP1): flow(IP1,...)

e.g.,

draft-seite-dmm-dma



move

AR2:
⚓ IP2
MN(IP2): flow(IP2,...)

③

Moving the IP address anchor address anchored (to AR2) in current network of attachment

Before moving:

MN in Net1 is allocated IP1 from AR1

MN and CN has flow using IP1 and IPCN, ...

After MN moves to Net2

MN in Net2 is allocated IP2 from AR2

MN and CN needs to continue the flow

Control Plane:

AR1 delegates prefix IP1 to AR2

Net1

LM:

IP1 ↔ IPAR2

FM-CP: DHVPv6-PD

AR1:

⚓ IP1
MN(IP1): flow(IP1,...)

Net3

AR3:
⚓ IPCN
CN:
① ↑
flow(IPCN, IP1, ...)

DHCPv6-PD, BGP/SDN



e.g.,

draft-mccann-dmm-flat40,

draft-liu-dmm-deployment-scenario,

draft-matsushima-stateless-uplane-vepec

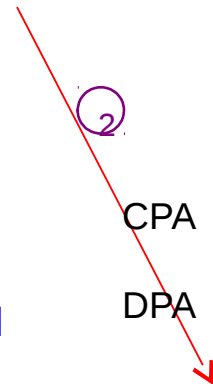
Data Plane Packet in flow(IPCN, IP1, ...)
from CN to MN:

1. Packet from CN → AR3
2. Forward from AR3 → AR2
3. Forwarded from AR2 → MN

②

CPA

DPA



Net2

LM:

IP1 ↔ IPAR2

AR2:

⚓ IP1, IP2
MN(IP2): flow(IP1, ...)

③ ↓

Address anchored to AR1 **not** in current network of attachment

Before moving:

MN in Net1 is allocated IP1 from AR1

After MN moves to Net2:

MN in Net2 is allocated IP2 from AR2

flow(IP1, ...) is using IP1 not in current net of attachment

(to be continued in next slide)

Net1

Net2

DPA

AR1:
⚓ IP1

AR2:
⚓ IP2
MN(IP2): flow(IP1,...)

Keeping the IP address address anchored to AR1 **not** in current network of attachment

Before moving:

MN in Net1 is allocated IP1 from AR1

MN and CN has flow using IP1 and IPCN,

Net3

After MN moves to Net2:

MN in Net2 is allocated IP2 from AR2

MN and CN need to continue the flow(IP1,IPCN,
...)

(to be continued in next slide)

AR3:

⚓ IPCN
CN:

flow(IPCN, **IP1**,
...)

Net2

Net1

DPA

AR1:

⚓ **IP1**
MN(IP1): flow(IP1,...)



AR2:

⚓ **IP2**
MN(IP2): flow(**IP1**,...)

Indirection of a flow (keeping the IP address)

Control Plane:

LM at AR1(CPA) and AR2(CPN) has binding information: $IP1 \leftarrow \rightarrow IPAR2$

FM-CP at AR1(CPA) and AR2(CPN) sets up indirection path between AR1(DPA) and AR2(DPN)

Net1

LM:
 $IP1 \leftarrow \rightarrow IPAR2$

FM-CP: sets up indirection

FM-DP: $AR1 \leftarrow \rightarrow AR2$

AR1:

$\uparrow IP1$
MN(IP1): $flow(IP1, \dots)$

Net3

AR3:
 $\uparrow IPCN$
CN:
 $flow(IPCN, IP1, \dots)$

Data Plane Packet in flow(IPCN, IP1, ...) from CN to MN:

1. from CN \rightarrow AR3
2. Forwarded from AR3 \rightarrow AR1 in Net1
3. FM-DP: Indirection AR1(DPA) \rightarrow AR2(DPN)
4. Forwarded from AR2 \rightarrow MN

Net2

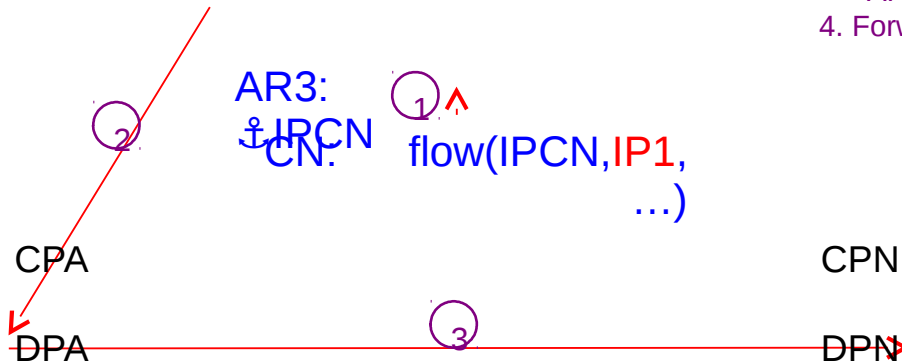
LM:
 $IP1 \leftarrow \rightarrow IPAR2$

FM-CP: sets up indirection

FM-DP: $AR1 \leftarrow \rightarrow AR2$

AR2:

$\uparrow IP2$
MN(IP2): $flow(IP1, \dots)$



e.g.,

draft-bernardos-dmm-cmp,
draft-bernardos-dmm-pmip, draft-sarikaya-dmm-for-wifi,
draft-yhkim-dmm-enhanced-anchoring

Changing indirection path (keeping the IP address)

Control Plane:

LM at AR1(CPA) and AR2(CPN) has the binding information: $IP1 \leftarrow \rightarrow IPAR2$;
LM information copied to AR3

FM-CP at AR3(CPA') and AR2(CPN) sets up indirection path between AR3(DPA') and AR2(DPN)

Net1

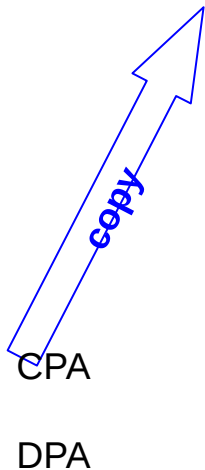
LM:
 $IP1 \leftarrow \rightarrow IPAR2$

FM-CP: sets up indirection

FM-DP: $AR1 \leftarrow \rightarrow AR2$

AR1:

$\uparrow IP1$
MN(IP1): $flow(IP1, \dots)$



Net3

LM:
 $IP1 \leftarrow \rightarrow IPAR2$

FM-CP: sets up indirection

FM-DP: $AR3 \leftarrow \rightarrow AR2$

AR3:

$\uparrow IP1$
CN: $flow(IPCN, IP1, \dots)$



e.g.,

paper-distributed.mobility



Data Plane Packet in $flow(IPCN, IP1, \dots)$
from CN to MN:

CPA'

1. from CN \rightarrow AR3

DPA'

2. FM-DP: Indirection from AR3(DPA') \rightarrow AR2(DPN)
3. Forwarded from AR2 \rightarrow MN



CPN

DPN

Net2

LM:
 $IP1 \leftarrow \rightarrow IPAR2$

FM-CP: sets up indirection

FM-DP: $AR3 \leftarrow \rightarrow AR2$

AR2:

$\uparrow IP2$
MN(IP2): $flow(IP1, \dots)$



Straw poll?

- ◆ Informational draft to explain enhanced/distributed anchoring and the different dmm approaches:
- ◆ draft-chan-dmm-distributed-mobility-anchoring-04 as informational draft
- ◆ Separate drafts on new features to standardize