Anycast vs. DDoS: Evaluating Nov. 2015 Root DNS event

*Giovane C. M. Moura*¹, Ricardo de O. Schmidt², John Heidemann³, Wouter B. de Vries², Moritz Müller¹, Lan Wei³, Cristian Hesselman¹

¹SIDN Labs ²University of Twente ³USC/ISI

2016-07-18

IEFT 96 – MAPRG/IRTF







IP Anycast

* It's simple: " making a particular Service Address available in multiple, discrete, autonomous locations" (RFC4786, 7094)

* Improves performance and resilience (1 IP → Many services, 1 down, others operate)

* Widely use in DNS (and also CDNs)



DDoS

- * Getting bigger (400Gbps +)
- * Getting cheaper (booters, few dollars)
- * Happening more often
- * Core idea: bring down services
- * Question: *How anycast behaves during a DDoS attack?*
- * Case study: Root DNS events Nov 2015







The Root DNS system

• List the records that points to all TLDs (.com, .nl, .net...)

letter	operator	sites (global, local)	architecture
Α	Verisign	5 (5, 0)	anycast
В	USC/ISI	1 (1, 0)	single site
С	Cogent	8 (8, 0)	anycast
D	U. Maryland	87 (18, 69)	anycast
Е	NASA	12 (1, 11)	anycast
F	ISC	59 (5, 54)	anycast
G	U.S. DoD	6 (6, 0)	anycast
Н	ARL	2 (2, 0)	primary/backup
I	Netnod	49 (49, 0)	anycast
J	Verisign	98 (66, 32)	anycast
K	RIPE	33 (15, 18)	anycast
L	ICANN	144 (144, 0)	anycast
М	WIDE	7 (6, 1)	anycast

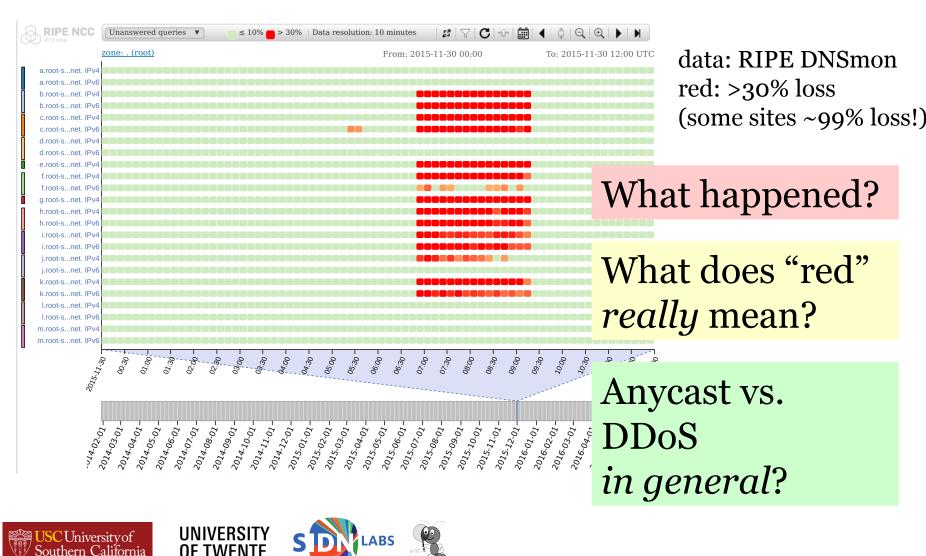
Table: The 13 Root Letters, each operating a separate DNS service, and their number of sites and architecture as of 2015-11-18.







A Bad Day at the Root...



Summary of the Events

Two events

- 2015-11-30t06:50 for 2h40m
- 2015-12-01t05:10 for 1h

affected 10 of 13 letters

about 5M q/s or 3.5Gb/s per affected letter

• aggregate: 155Gb/s

real DNS queries, common query names, from spoofed source IPs **implications:**

some letters had high loss

overall, though DNS worked fine

clients retried other letters (as designed)







data: A-Root had full view (Verisign presentation); RSSAC-002 reports

How Well Does Anycast Defend?



561 root DNS locations for **13 services** (in 2016-01)

is 561 *too few? too many?* what happens *under stress?*







Anycast in Good Times

anycast matches a **user** to a (hopefully) nearby **site**

X-SJC

you

anycast divides the Internet into catchement (often messy and

(some **sites** have

more capacity)

USC University of Southern California





your

friend

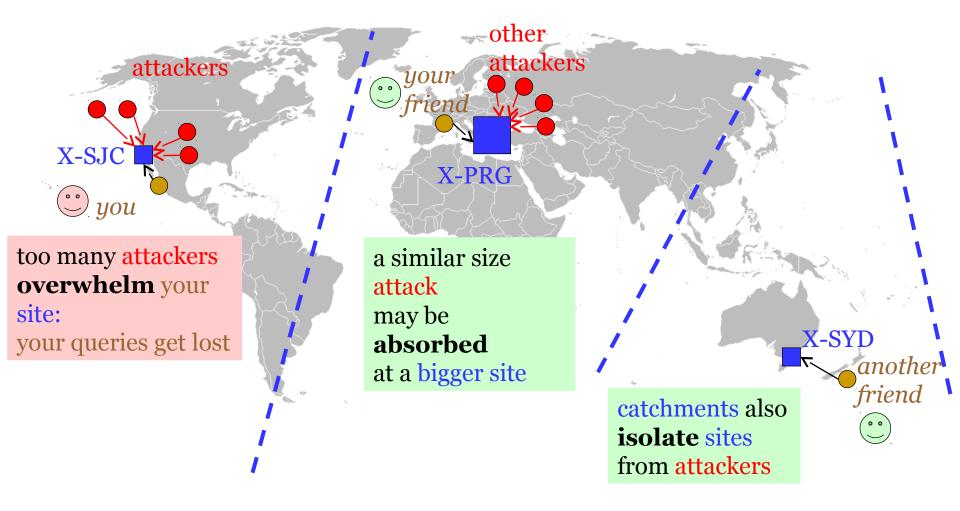
X-PRG

anothe

friend

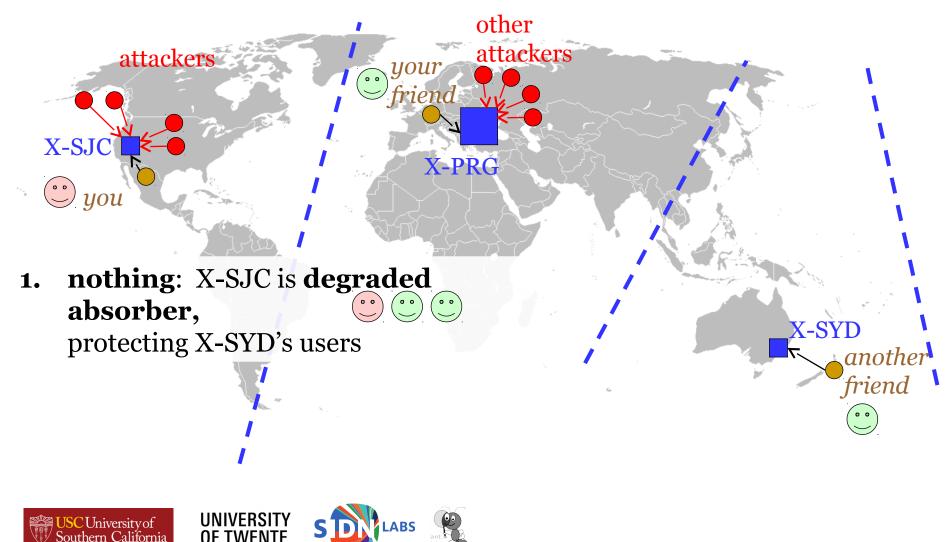
X-SYD

Anycast Under Stress





Anycast Under Stress (do nothing)



Anycast Under Stress (withdraw some routes)

uour

___other

X-PRG

attackers

1. nothing: X-SJC is degraded absorber,

attackers

X-SJC

vou you

2. withdraw routes from X-SJC; may shift attackers to big site



another

friend

X-SYD

Anycast Under Stress (withdraw other routes)

your

friend

X-PRG

other

attackers

1. nothing: X-SJC is degraded absorber,

attackers

- 2. withdraw routes from X-SJC; may shift attackers to big si
- **3. withdraw** wrong routes from X-SJC; may shift attackers to other sit



you you





another

friend

Best reaction to stress: you don' t know

don't know: attacker number of attackers location of attackers affects of routing change

1. nothing: X-SJC is degrade absorber, \mathbf{x} $\mathbf{\alpha}$

attackers

don't fully control routing and

other

- 2. withdraw routes from X-SJ catchments may shift attackers to big si 🙂 🙂 🙂
- hard to make withdraw wrong routes from X-SJC; 3. informed choices may shift attackers to other sit

101r

X-PRG

X-SJC

🕐 you





another

friend

-SYD

What Actually Happens?

studying Nov. 30

we see **withdrawals** and **degraded absorbers** some clients loose service

• results vary by anycast deployment



Data About Nov. 30

RIPE Atlas

- 9000 vantage points (RIPE Atlas probes)
- try every letter every 4 minutes
 - except A-root, at this time, was every 30 minutes
- data-plane queries
- global, but heavily biased to Europe
- RSSAC-002 reports
- self-reports from letters
- not guaranteed when under stress
- **BGPmon** routing
- control plane



How About the Letters?

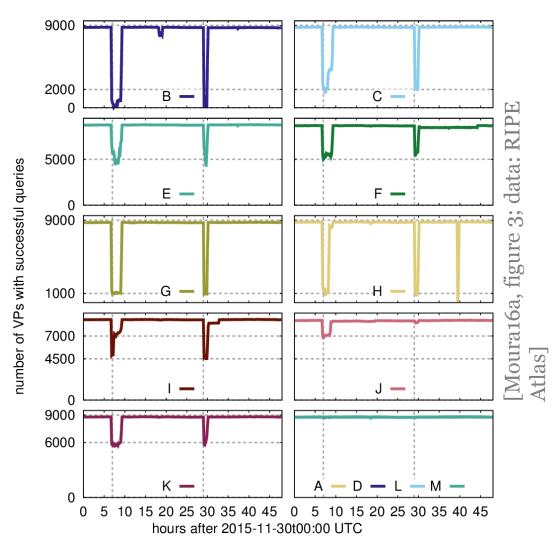
some did great: D, L, M: not attacked A: no visible loss

most suffered: a bit (E, F, I, J, K) or a lot (B, C, G, H)

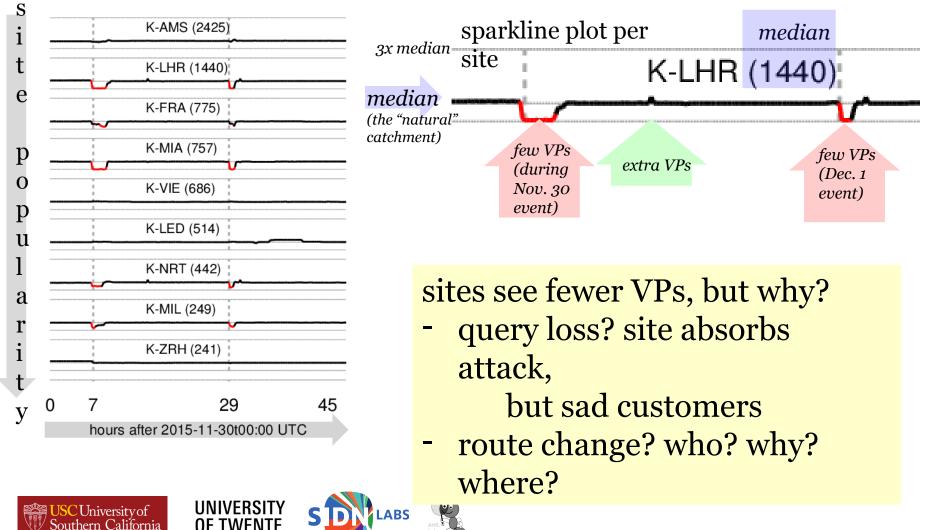
but does "x%" measure what *users actually see?*

University of thern California

UNIVERSITY



Reachability at K-sites

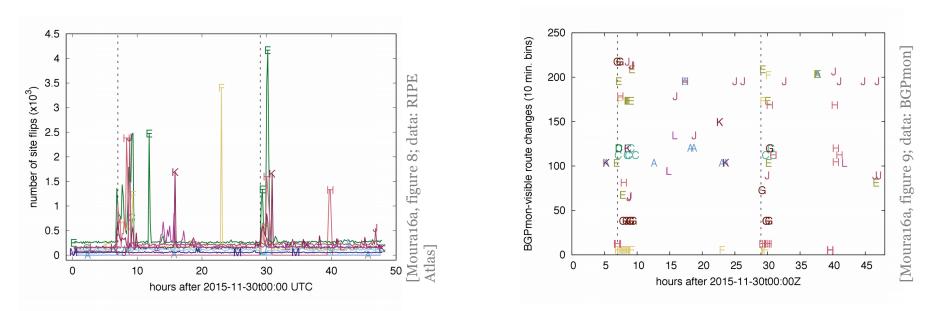


Site *Flips* from Routing Changes V a 360 minutes (in 4 minute bins) n Nov. 30 event t a stay at K-LHR; g sad during event e flip to K-AMS; (less) sad during P yellow: Kevent; 0 back to K-LHR after LHR flip to K-other blue: K-AMS and stay there n white: K-other flip to K-AMS S black: failed query [Moura16a, figure 11b; data: RIPE Atlas] UNIVERSITY C University of

Southern California

OF TWENTE

Confirming flips in BGP



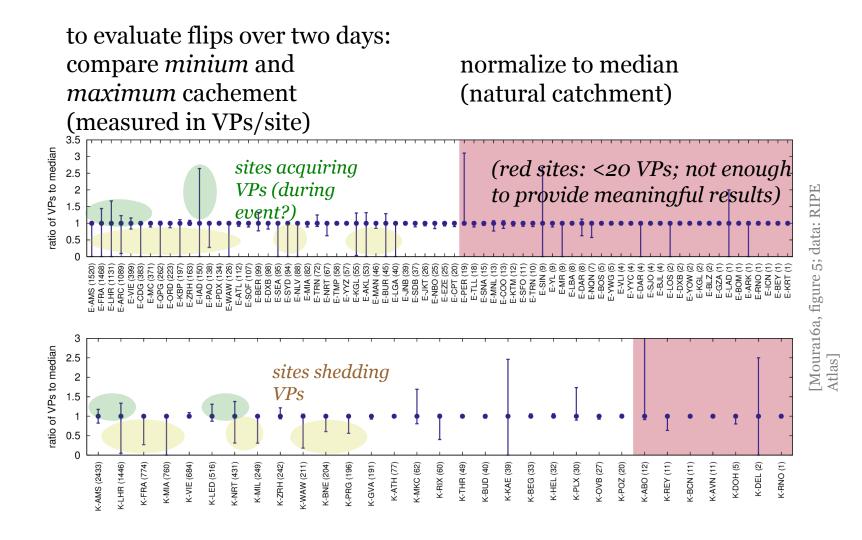
flips common during events for most letters

J<mark>SC</mark>University of outhern California UNIVERSITY OF TWENTE.





Flips Across Letters: E and K











Flips Implications

some ISPs are "sticky" and won't flip

• will suffer if their site is overloaded

some ISPs will flip

• but new site may not be much better

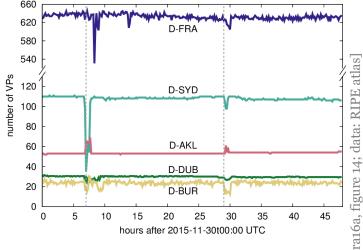
result depends on many factors

- *actions* taken by root operator
- routing choices by operator *and peer*
 - and perhaps *peer's peers*, depending on congestion location implementation choices
- DNS, routing



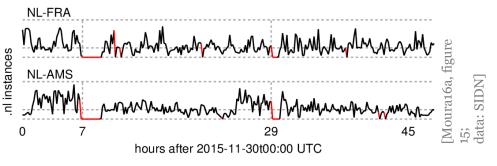
Aside: Collateral Damage

can an event hurt non-targets?
yes! ...a risk of shared datacenters



D-FRA and D-SYD: less traffic (even though D was not directly attacked)

University of



.NL-FRA and .NL-AMS: no traffic

In other attacks, B-Root's ISP saw loss to other customers

Conclusions

anycast under stress is complicated

- some users will see persistent loss
- "x% loss" is not complete picture

reactions depend on design and implementation choices

• many not under operator control

more info:

paper: <u>http://www.isi.edu/~johnh/PAPERS/Moura16a/</u> data: <u>https://ant.isi.edu/anycast/</u>

