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CCNX INTEREST AGGREGATION

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PRINCIPLES

- Minimize coupling between consumer and forwarders
- No timers on forwarders
- Provable properties

DIFFERENCE FROM ICNRG -03 DRAFT

- Spelled out algorithms in more detail
- Adopted Interest HopLimit decrement idea [1]

[1] J.J. Garcia-Luna-Aceves and M. Mirzazad-Barijough. Enabling Correct Interest Forwarding and Retransmissions in a Content Centric Network. In *Proc. ACM ANCS'15*, pages 135–146, Oakland, California, 2015.

PROPERTIES (1-3)

1. **Aggregation:** If there are no losses and no retransmissions, then for each link in the forwarding path exactly one Interest and at most one ContentObject traverse the link.
2. **Retransmissions:** If there are retransmissions by one or more consumers and those retransmissions do not pass a ContentObject response in flight, then at most one ContentObject traverses the reverse path.
3. **Interest In Flight:** If a retransmission passes a ContentObject in flight on a given link, then that Interest will propagate only as far as the first cached copy of the response.

PROPERTIES (4-6)

4. **ContentObject In Flight:** If a retransmission passes a ContentObject in flight on a given link, then there will be a duplicate ContentObject sent on that link, but it will not propagate further than that link.
5. **Cycle Termination:** An Interest that travels in a cycle will not repeat the same cycle.

PROPERTY 6: CYCLE TERMINATION

- The **Cycle Termination** property is significantly different than prior CCNx 0.x and NDN cycle termination. In those protocols, because of the nonce, an Interest that visits the same node twice – assuming the nonce is remembered long enough – is dropped.
- This means that such a node cannot continue using other paths, even if those will reach the destination.
- In CCNx 1.0, the cycle termination property means that an Interest will not repeat the same cycle, but it could be forwarded along another shorter path.
- Described in more detail in a few slides.

HOPLIMIT

- An Interest predecessor may be remote (i.e. another node) or local (i.e. an application connecting directly to the forwarder).
- The HopLimit indicates the number of allowed remote node hops.
- Forwarding to and from an application does not count as a hop.
- This semantic allows an implementation, if desired, to indicate “for me” in the FIB table via routes with a 0 HopCount, which is the notional model we use here. As this is internal behavior of a forwarder, it is not standardized.

Algorithm 1 Receive Interest

```
1: procedure RECEIVEINTEREST(Predecessor P, Interest I)
2:   if P is remote then
3:     if  $I.HopLimit = 0$  then
4:       Send InterestReturn (HopLimit)
5:       Drop Interest
6:     else
7:       Decrement  $I.HopLimit$ 
8:     end if
9:   end if
10:  if Satisfy I from ContentStore then
11:    Send ContentObject to P
12:  else
13:     $Verdict \leftarrow AGGREGATE(P, I)$ 
14:    if Verdict is Forward then
15:      if FORWARDINTEREST(P,I) is false then
16:        Drop Interest
17:      end if
18:    end if
19:  end if
20: end procedure
```

HopLimit must decrease each hop.

Only first Interest from a new predecessor for existing PIT entry is aggregated.

Algorithm 2 Interest Aggregation

```
1: function AGGREGATE(Predecessor P, Interest I)
2:   if does not exist a PIT entry for I then
3:     Create PIT entry with Predecessor P
4:     Return Forward
5:   else if PIT entry exists with Predecessor P then
6:     Retransmission from P
7:     Return Forward
8:   else
9:     → Add Predecessor P to PIT Entry
10:    Return Aggregate
11:   end if
12: end function
```

FIB hop count must not exceed the decremented HopLimit (Alg 1, Line 7)

Algorithm 3 FIB Lookup

```
1: function FIBLOOKUP(Predecessor P, Interest I)
2:   Exclude FIB entries that point to P
3:   NextHops ← LongestPrefixMatch(I)
4:   for N ∈ NextHops do
5:     → if I.HopLimit < N.HopCount then
6:       Remove N from NextHops
7:     end if
8:   end for
9:   Return NextHops
10: end function
```

Interest output has
HopLimit set to the FIB's
HopCount (which must
be less than the input
HopLimit, Alg 3 Line 5)

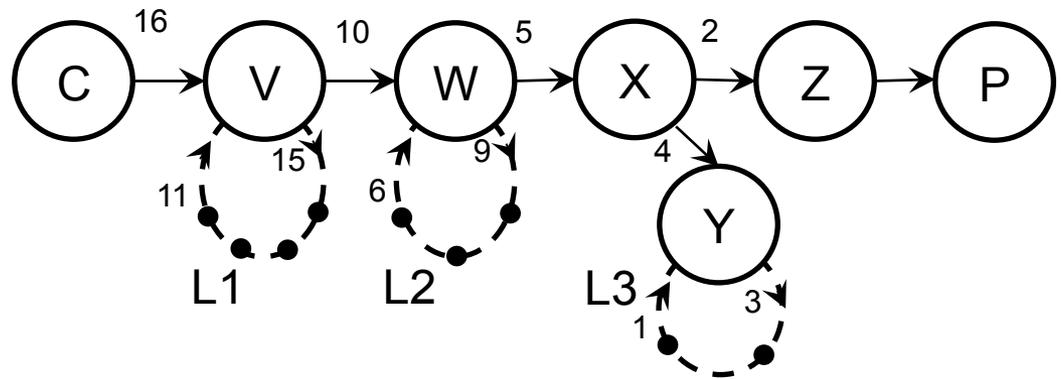
Algorithm 4 Forward Interest

```
1: function FORWARDINTEREST(Predecessor P, Interest I)
2:   NextHops  $\leftarrow$  FIBLOOKUP(P, I)
3:   if NextHops is not empty then
4:     for  $N \in \textit{NextHops}$  do
5:        $I.HopLimit \leftarrow N.HopCount$ 
6:       Send I to N
7:     end for
8:     Return true
9:   else
10:    Send InterestReturn (NoRoute)
11:    Return false
12:  end if
13: end function
```

PROPERTY PROOFS

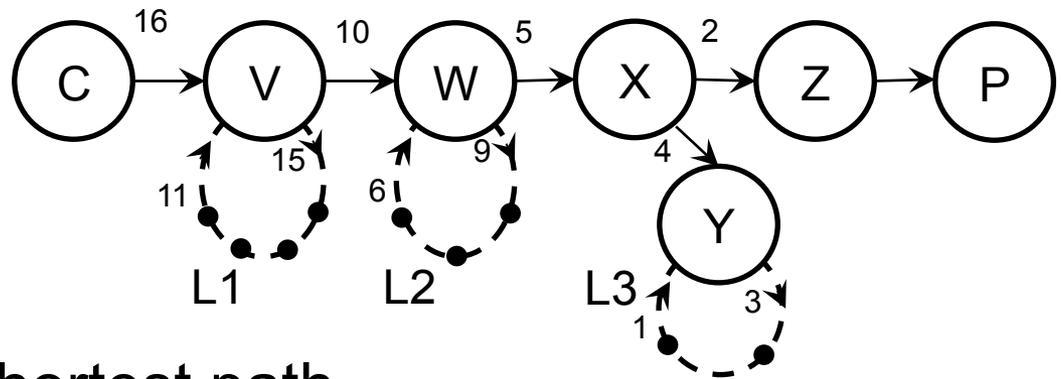
- Not presented here, written up in paper submission.
- Can sketch out the proofs in Q&A if you have specific questions.

CYCLES



- Use this topology for example.
- Assumes each forwarder uses unequal cost multipath.
- Numbers indicate the FIB HopCount of each successor.
- Interest sent from consumer C towards producer P.

CYCLES



- **C-V-W-X-Z-P**: The shortest path.
- **C-V-L1-V-W-X-Z-P**: The cycle L1 is permissible because the HopLimit will be 15 on entering the loop and 11 on exiting the loop, so the successor W is feasible.
- **C-V-W-L2-W-X-Z-P**: The cycle L2 is permissible because the HopLimit will be 9 on entering the loop and 6 on exiting the loop, so the successor X is feasible.
- **C-V-L1-V-W-L2-W-X-Z-P**: As in prior two cases.
- **C-V-W-X-Y-L3-Y-{}**: The path via Y will terminate after the first cycle because the HopLimit on exiting L3 will be 1 and Y has no successor with a feasible HopCount.

CONCLUSION

- CCNx 1.0 Interest aggregation algorithm
 - Does not use timers.
 - Does not depend on the ARQ mechanism (nor does the ARQ mechanism depend on the Interest aggregation scheme).
- Has desirable properties
 - At most 1 interest forwarder if no retransmissions.
 - At most 1 ContentObject downstream per link, even with retransmissions (ignoring in-flight misses).
 - Interest will not travel a cycle more than once.