



101st IETF/IRTF, Mar. 2018, London

CCNinfo: Discovering Content and Network Information in Content-Centric Networks

draft-asaeda-icnrg-ccninfo-00

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Changes

- Replace Contrace with CCNinfo
 - Old: "Contrace: Traceroute Facility for Content-Centric Network"
 - https://tools.ietf.org/html/draft-asaeda-icnrg-contrace-04
 - New: "CCNinfo: Collecting Content and Network Information in Content-Centric Networks"
 - https://tools.ietf.org/html/draft-asaeda-icnrg-ccninfo-00
 - Several small changes from contrace-04

Brief introduction of CCNinfo

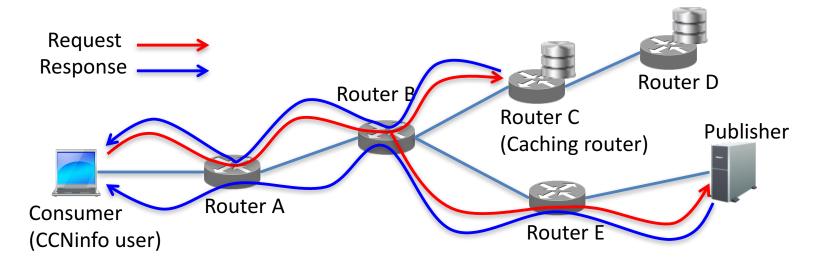
- What is CCNinfo?
 - CCNinfo aims for discovering/retrieving information about the network (incl. multi-path) topology and in-network cache in CCN
- CCNinfo collects the information
 - Caching router (per prefix)
 - Number of hops from consumer to caching router / publisher
 - RTT between consumer and caching router / publisher
 - Cache hit ratio
 - Number of the accesses (i.e., received Interests) per cache / chunk
 - Lifetime and expiration time per cache / chunk
 - Path stretch ("d / P")
 - where "d" is the hop count of the data and "P" is the hop count from the consumer to the publisher

Brief introduction of CCNinfo – cont'd

- Why is CCNinfo needed? What is the benefit?
 - For protocol analysis
 - Effectiveness, robustness, and cost of designed networks
 - Hop count / RTT for content retrieval, multipath, in-network caching algorithm
 - Characteristics of content
 - Content popularity, cache hit ratio
 - As an OAM tool
 - Operation
 - Cache lifetime or expiration time
 - Monitoring
 - CS capacity and usage at router, num. of interests per content
 - Trouble shooting
 - Availability of caching routers and publishers
 - Yet, policy-based information provisioning
 - Fine-grained access control and policy configuration for information disclosure highly required

CCNinfo Basic Behavior

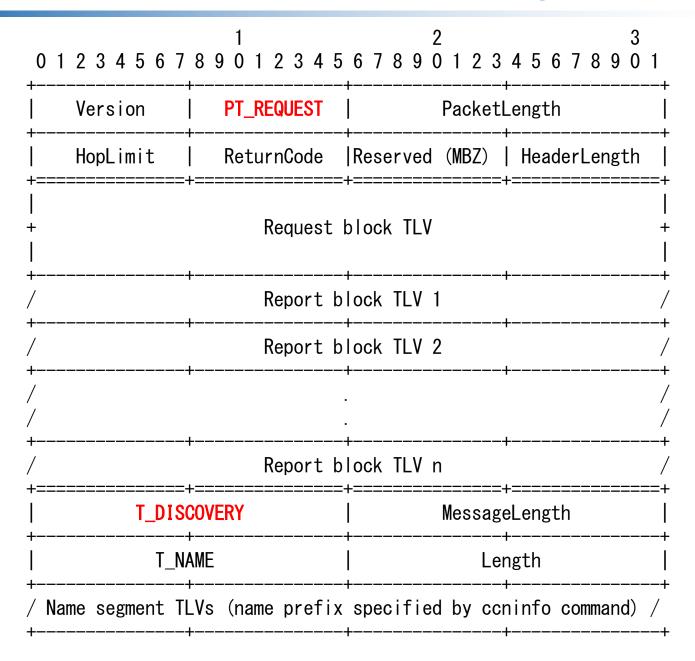
- Request message is initiated by CCNinfo user and forwarded toward caching router or publisher based on the FIB in a hopby-hop manner
- Request message includes Request block and Report block(s)
- Reply message is initiated by caching router or publisher and forwarded toward CCNinfo user based on the PIT entry
- Reply message includes Reply block and Reply sub-block(s)



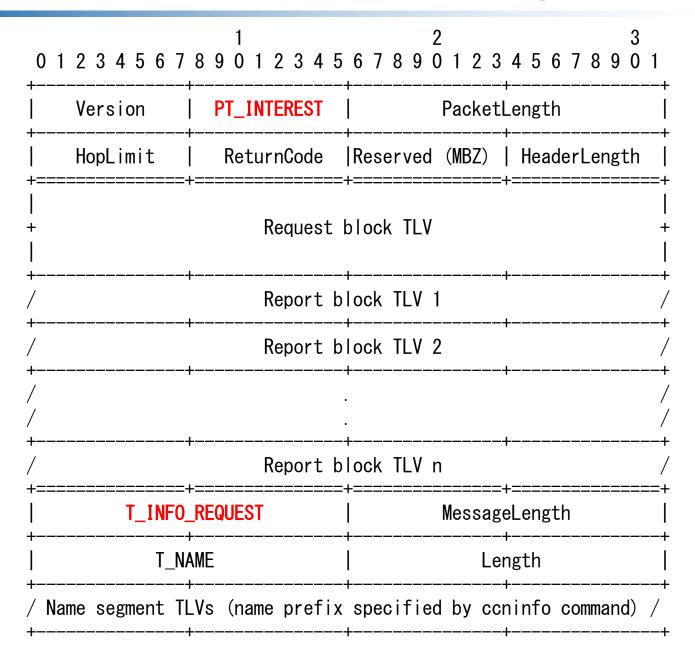
CCNinfo Request/Reply Messages

- Compatible with CCNx-1.0 TLV format
- CCNinfo Request Message
 - Request message consists of a fixed header, Request block TLV, Report block TLV(s), and Name TLV
- CCNinfo Reply Message
 - Reply message consists of a fixed header, <u>Request block TLV</u>, <u>Report block TLV(s)</u>, Name TLV, and and <u>Reply block/sub-block</u> <u>TLV(s)</u>
- Case 1 (defined in the current I-D): New Request type (PT_REQUEST) and Reply type (PT_REPLY) values in the fixed header are used
- Case 2: No new protocol type, but new message type values (T_INFO_REQUEST / T_INFO_REPLY) for interest / data are required
 - PT_INTEREST and PT_CONTENT are used

Case 1: Request Message



Case 2: Request Message



Case 1: Reply Message

0 1 2 3 4 5 6 7	8 9 0 1 2 3 4 5	6789012	3 4 5 6 7 8 9 0 1										
Version	PT_REPLY	PacketLength											
HopLimit	ReturnCode	Reserved (MBZ)	HeaderLength										
Request block TLV													
+													
T_DIS	COVERY	+=====================================											
T_N/	+ AME	Length											
/ Name segment TI	_Vs (name prefix	specified by c	cninfo command) /										
+ /	Reply b	lock TLV											
 -	Reply sub-h	olock TLV 1	- 										
+ / /	·												
/	Reply sub-l	olock TLV k	_+										

Case 2: Reply Message

	ı	Z		ა									
0 1 2 3 4 5 6 7	8 9 0 1 2 3 4 5	67890	1 2 3 	4 5 6 7 8 9 0 1									
Version	PT_CONTENT	PacketLength											
HopLimit	ReturnCode	Reserved (M	 MBZ) 	HeaderLength									
+ Request block TLV													
+													
T_INFO	D_REPLY												
T_N/	 AME	Length											
/ N ame segment TL	_Vs (name prefix	specified by coninfo command)											
/	Reply b	lock TLV											
 	Reply sub-b	plock TLV 1		,									
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Case 1 vs. Case 2

- Different operations from the regular Interest / Data are required (= less trouble, easy-to-understand implementation)
 - Routers SHOULD NOT remove the PIT entry created by the CCNinfo Request until timeout value expires.
 - CCNinfo Requests SHOULD NOT result in PIT aggregation in routers during the Request message transmission.
 - CCNinfo Replies MUST NOT be cached in routers upon the Reply message transmission.

Case 1 vs. Case 2 – cont'd

- Policy-based information provisioning: different and finegrained access control for authorized / unauthorized
 Requests are required, for example,
 - Router SHOULD forward the Request message to the upstream router toward the content forwarder
 - However, according to the "Node Identifier" in the Request block or "Signature" in the ValidationAlgorithm TLV, router SHOULD configure either of;
 - (1) All (all cache information is disclosed)
 - (2) Partial (cache information with the particular name prefix can (or cannot) be disclosed), or
 - (3) Deny (no cache information is disclosed)

A brief comparison with ICN traceroute

	CCNinfo	ICN traceroute
Objective	Collect comprehensive information to support diagnostic test and statistical analysis	Obtain limited information of specific route for debug and diagnosis
Methodology	Explore the tree structure from the consumer to the routers and the producer	Explore the basic information of specific route between the consumer and the producer
Algorithm	Interest-like requests, but waiting for the reply until time-out	Iterative algorithm similar with IP traceroute
Multipath	Works like ordinary interests	Possibly, but not specified
Protocol type	Case 1: New protocol types required Case 2: No new protocol type, but new type values for interest/data required	No new protocol type
Reserved name	No reserved name	Reserve "traceroute" suffix

Summary

 CCNinfo, which is compatible with CCNx-1.0 TLV format, is a powerful network tool providing various information

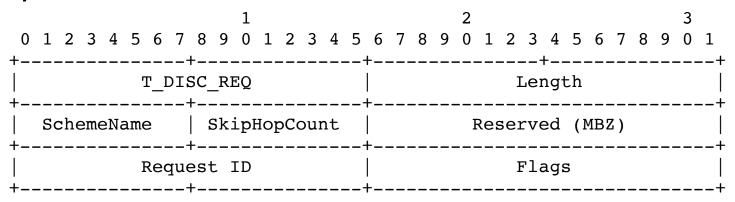
- RG draft?
 - We'll then start discussion more on new message types, etc.



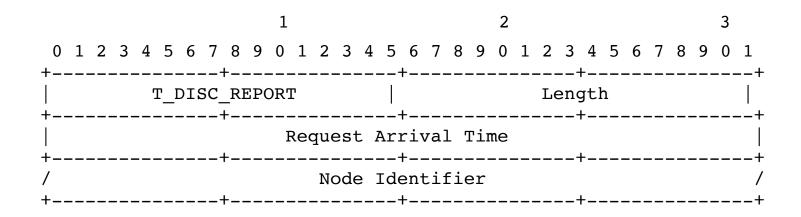
Backup

Request Block and Report Block

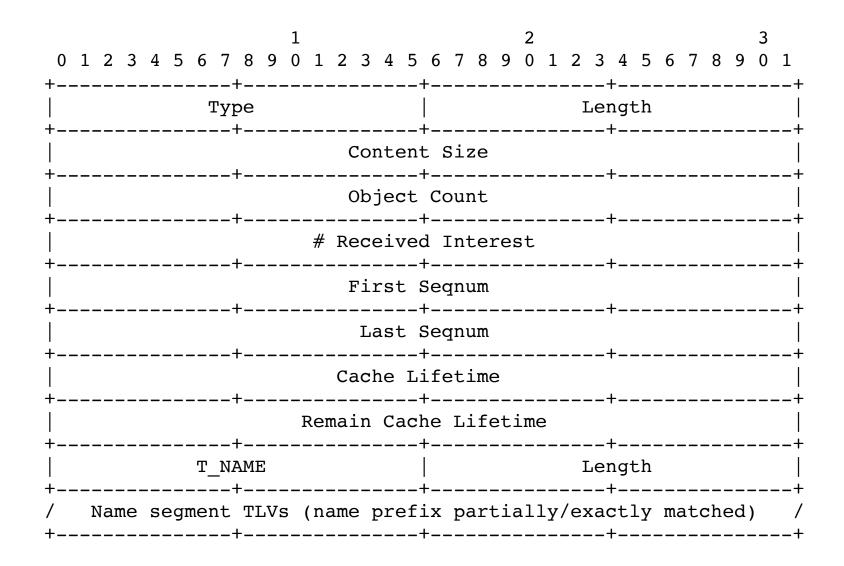
Request block



Report block



Reply Sub-Block for T_DISC_CONTENT and T_DISC_CONENT_OWNER



Reply Sub-Block for T_DISC_GATEWAY

		1										2												3							
0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1
+-								+							1	-							+	 -							+
	Туре									Length																					
+-	+										-								 -							+					
	Scheme Name								Reserved (MBZ)																						
+_	+											-								-							+				

Security Considerations

- Policy-Based Information Provisioning for Request
 - Routers can reply the CCNinfo Reply with the ADMIN_PROHIB return code without appending any Reply (sub-)block TLV
 - Permission, whether (1) All (all cache information is disclosed), (2) Partial (cache information with the particular name prefix can (or cannot) be disclosed), or (3) Deny (no cache information is disclosed), can be defined at routers
- Filtering of CCNinfo Users Located in Invalid Networks
 - Routers may support an access control mechanism to filter out Requests from invalid CCNinfo users located in invalid networks

Security Considerations – cont'd

- Topology Discovery
 - If a network topology is a secret, CCNinfo Requests may be restricted at the border of the domain, using the ADMIN_PROHIB return code
- Characteristics of Content
 - If publisher or content information is secret, CCNinfo Requests may be restricted at the border of the domain, using the ADMIN_PROHIB return code
- Shortening CCNinfo Reply Timeout
 - Routers may configure the shorter timeout value to time out the Request
- Limiting Request Rates
 - Routers may rate-limit CCNinfo Requests by ignoring some of the consecutive messages.

Security Considerations – cont'd

- Limiting Reply Rates
 - Routers in the traced path may rate-limit CCNinfo Replies.
- Adjacency Verification
 - Forwarding CCNinfo messages given from nonadjacent neighbor nodes/routers must be prohibited
 - Defining the secure way to verify the adjacency cannot rely on the way specified in CCNx message format or semantics; therefore a new TLV for adjacency verification using hop-by-hop TLV header will be defined in a separate document.