

RTP Topologies

Clue Interim June 2012

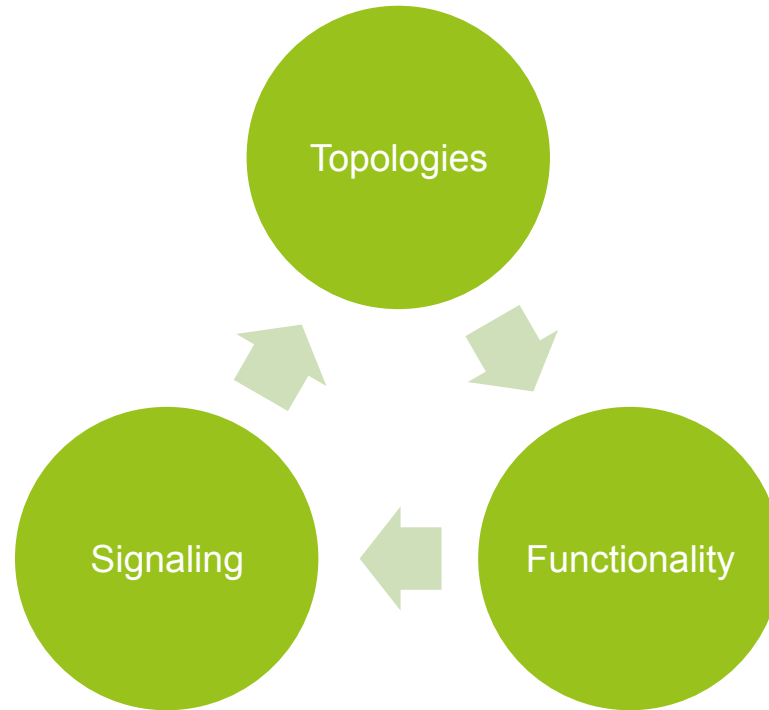
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Outline



- › The Big Issue
- › Goal
- › Evaluation Criteria
- › Topologies
- › Conclusions

The Big Issue



Signaling can restrict what topologies that are supported, and thus the functionality a CLUE system may have!

Goal



- › Present what **RTP functionality** a given topology enables
- › Start a discussion on what topologies and media plane functions CLUE wants to support
 - To ensure correct requirements on the signaling

Evaluation Criteria



› Security

- Key-management
- Nodes that are in the security context (who have the keys)?
 - › Trust Structures
- Source Authentication
 - › End-to-End verifiable
 - › Trust in central node

Evaluation Criteria



› Congestion Control

– Multi-hop

- › Need for information and requests to bridge across hops

– One or Multiple Receivers of the same RTP stream

- › Meet requirements from multiple end-points

– Transcoding

- › Enabling bit-rate adjustments
- › Breaking multi-hop control loops

– Media Aggregate adjustments

- › Prioritization between streams

Evaluation Criteria



› Source Identification

- Receiver must be able to determine source of media
 - › Reference in Meta-Information
 - › Identity for Control Requests
- Media mixing or compositions
 - › Multiple contributing sources
- Translation of source identification information
 - › Require additional layer of identification labels
 - › OR
 - › Force all end-point communication through node that translates
- Conference Wide common identity space required?

Evaluation Criteria



- › Bandwidth Consumption
 - Deliver most appropriate media properties
 - › Transcode
 - › Choice from Simulcast alternatives
 - › Source Codec Control
 - Select the N out of M streams most needed by the application
 - › Possibility to Prune unneeded streams
 - Mix or composite N streams into one
 - Translation transcode / re-encode
 - › Increases bandwidth usage to maintain quality

Evaluation Criteria



› Media Quality

- Transcoding / re-encoding

- › Quality reduced per spent bit

- Delay

- › Need to be kept low in to maintain interactivity

- › Inter continental communication

- The rate vs distortion relation is approximately logarithmic

- › Quality gain per bit will affect prioritization between:

- Increasing one stream's quality

- Allowing additional streams

- Difficult trade-off between quality, delay, bit-rate consumption and functionality

Evaluation Criteria



› Distribution of Complexity

– Various factors, e.g.

› Processing

› Memory

› Implementation cost

– Depending on Topology

– Some complexity can be moved between central nodes and end-points

› Impact on a central node can be different from an end-point for a given functionality

– Node Limitations must be taken into account

› Forces location of functionality

› Can cost quality

Topologies Outline



- › Point to Point
- › Distributed End-point
- › Multi-Unicast (MESH)
- › Mixers
 - Media Mixer
 - Media Switching
 - Source Projection
- › Relay (Transport Translator)
- › Selective Forwarding
- › End-point Forwarding
- › Any Source Multicast
- › Sender Source Multicast

Point to Point



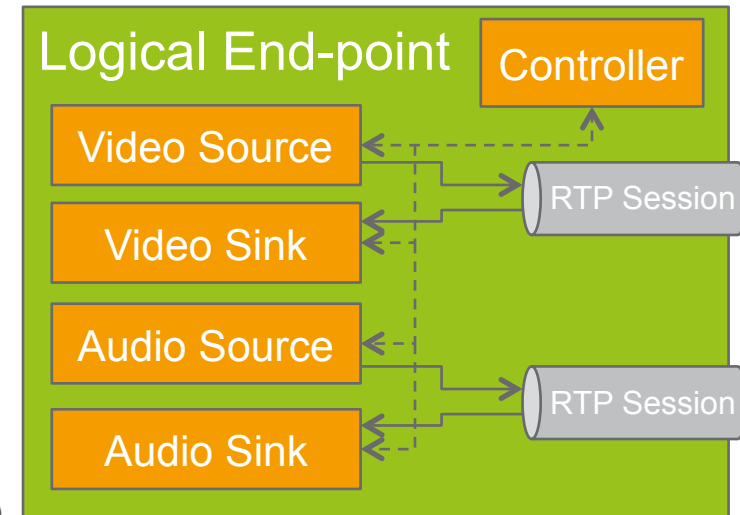
- › Description:
 - One peer communicate directly with a single peer over unicast
- › Security:
 - Authentication of Peer / User
 - Confidentiality and Integrity between peers
- › Congestion:
 - Receiver can report statistics or request direct adaptation from sender
- › Identification:
 - Senders sources map one to one with RTP media streams
- › Bandwidth:
 - Receiver can request media to be tailored to its needs
 - Action to increase or decrease bandwidth can depend on the current path capacity
 - › At high capacity add additional streams to provide additional functions
 - › At low reduce to single stream and focus at maximize quality for the most important content
- › Quality:
 - Optimal in relation to Path Capacity and Properties
- › Complexity:
 - All in the end-points
 - Limitations in end-point directly affect what sender can produce and receiver can accept



Distributed End-point



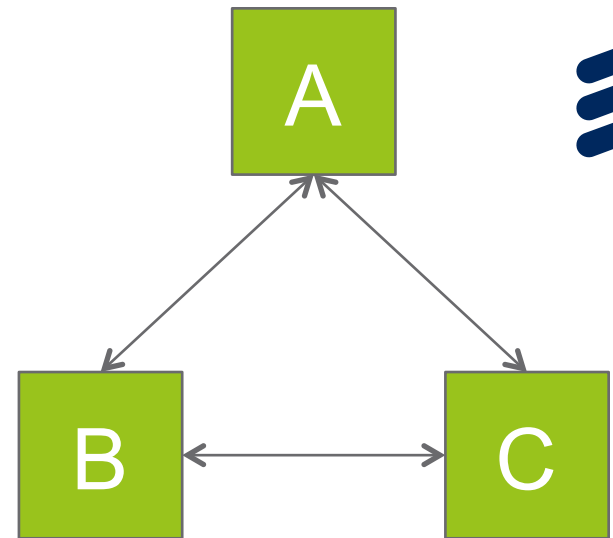
- › Description:
 - Distributed realization of a logical end-point
 - Different IP addresses for various components
 - › Camera (Video source)
 - › Display (Video sink)
 - › Microphone / Audio mixer (Audio source)
 - › Loudspeaker (Audio sink)
 - › Controller (Signaling end-point)
 - There can be multiple instances of one component type
- › Security:
 - **Each source or Sink must be keyed with the other end-points key(s)**
 - Controller responsible to provide logical end-point identity
- › Congestion:
 - Receiver component can report statistics or request direct adaptation from media sending component
 - **Prioritization between media streams in the aggregate are complicated by distribution**
 - Due to different source / destination addresses network load balancer may give different routes to different flows
- › Identification:
 - Senders sources map one to one with RTP media streams
 - A logical end-point may have multiple presences in an RTP session due separation of sources and sinks
 - **Multiple different IP addresses or hidden behind aggregation point**
- › Bandwidth:
 - **Trade-off between centralized control and distributed handling of adaptation and prioritization**
- › Quality:
 - May become sub-optimal in relation to Path Capacity and Properties due to control latencies
- › Complexity:
 - Additional complexities for control within the end-point



Multi-Unicast (Mesh)



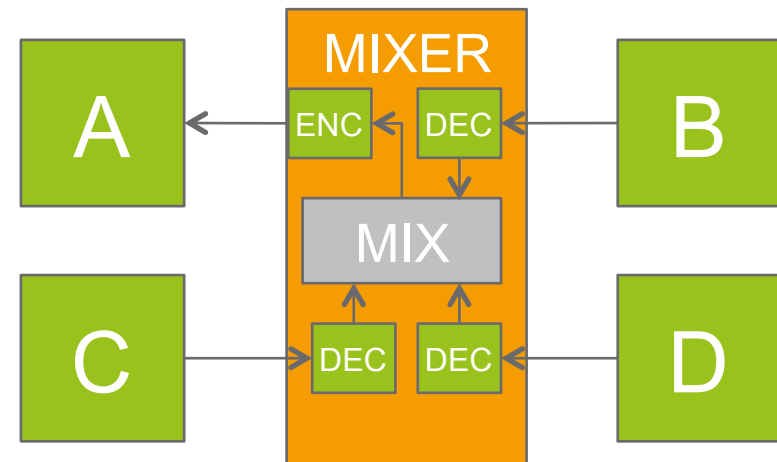
- › Description:
 - One peer communicate directly with multiple peers
 - Each peer to peer communication is independent unicast
 - **Each peer pair can have its own RTP session**
- › Security:
 - Individual Authentication of each Peer / User
 - Confidentiality and Integrity between pair of peers
- › Congestion:
 - Receiver can report statistics or request direct adaptation from sender
 - **All Peers will commonly share first hop/hops and the available capacity / bottleneck**
 - Sender can produce independently encoded media or produce one encoding sent to multiple peers.
- › Identification:
 - Sender's sources map one to one with RTP media streams within one RTP session
 - **Using multiple RTP sessions results in independent SSRC/CSRC spaces between the sessions**
 - › Could select to use unique values over multiple RTP sessions or use different layer
- › Bandwidth:
 - **Receiver can request media to be tailored to it's needs**
 - › **May be forced to accept a compromise based on other paths in case sender share media encoder**
 - Desirable to enable different trade-offs based on path capacity
- › Quality:
 - Can be optimal in relation to Path Capacity and Properties
 - **To reduce sender complexity in encoding less than optimal quality may be received**
- › Complexity:
 - All in the end-points
 - Limitations in end-point directly affect what sender can produce and receiver can accept
 - Trade-off in amount of complexity each pair of peers create can affect conference properties



Media Mixer



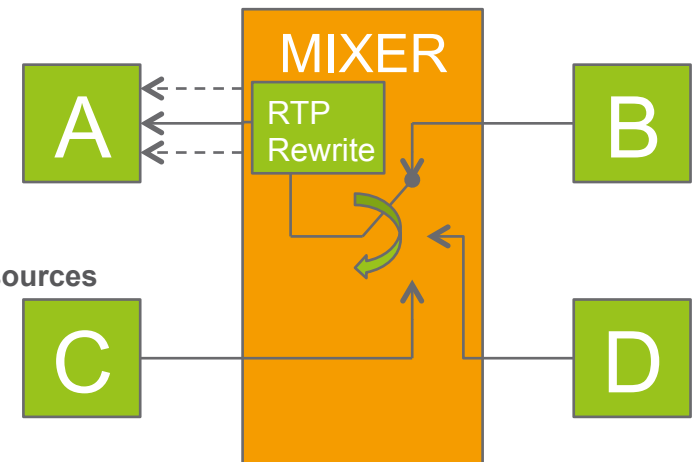
- › Description:
 - One peer communicates only with the Mixer
 - Each peer to mixer communication is independent unicast
 - **Mixer provides a mixed or composited media source based on the media streams from the other participants**
 - Each communication can have its own RTP session, or Mixer can create a conference-wide RTP session by sharing SSRC / CSRC
- › Security:
 - Mixer handles Authentication of each Peer / User
 - **Mixer is trusted entity and enforcer of some security functions**
 - Confidentiality and Integrity between peer and Mixer
- › Congestion:
 - Receiver (Mixer or peer) can report statistics or request adaptation from sender (Mixer or peer) on their path
 - Mixer can choose to forward report / request information (unaltered or aggregated) between paths
 - **Mixer typically produce independently encoded media to each peer, but may re-use some media between receiving peers**
- › Identification:
 - **Sender's sources are only visible as contributing sources in Mixer's RTP media streams**
 - Using multiple RTP sessions results in independent SSRC/CSRC spaces between the sessions
 - › Could select to use unique values over multiple RTP sessions or use different layer
- › Bandwidth:
 - **Mixer can reduce the number of concurrent media streams to a single per media type**
 - Receiver (also Mixer) can request media to be tailored to its needs
- › Quality:
 - Maximum Quality limited by participant to mixer path capacity
 - **Quality loss and delay increase in decoding encoding cycle**
- › Complexity:
 - **Mixer has one end-point complexity per end-point in the conference, plus media composition and some Mixer-specific logic**
 - Mixer proxies limitations in end-point affecting what sender can produce and receiver can accept, but may add further limits



Media Switching Mixer



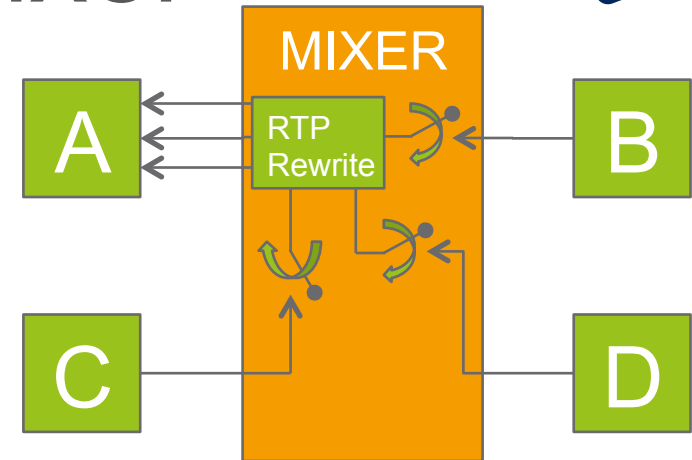
- › Description:
 - One peer communicate only with Mixer
 - Each peer to mixer communication is unicast with mixer feedback
 - **Mixer provides one or more conceptual sources selecting original sources**
 - Mixer creates a conference-wide RTP session by sharing SSRC / CSRC
- › Security:
 - Mixer handles Authentication of each Peer / User
 - Mixer is trusted entity and enforcer of some security functions
 - Confidentiality and Integrity between peer and Mixer
- › Congestion:
 - Receiver (Mixer or peer) can report statistics or request adaptation from sender (Mixer or peer) on their path
 - **Mixer needs to aggregate and forward report / request information between paths, based on some policy**
 - **Mixer distributes encoded media to multiple peers, making single receiver limitation affect more receivers**
 - **Mixer can make use of simulcast or scalable media encoders from senders to adapt to a peer**
- › Identification:
 - Sender's sources are only visible as contributing sources in Mixer's RTP media streams
 - Mixer can have multiple SSRCs representing different conceptual media sources
- › Bandwidth:
 - Receiver (also Mixer) can request media to be tailored to its needs, but will typically also affect other receivers
 - **Desirable to limit the amount of trade-off based on path capacity**
 - Simulcast and scalability can be used to meet different bandwidth needs or requirements
- › Quality:
 - Trade-off of end-to-end Path Capacity and Properties between receivers sharing media from the same sender
 - **Avoids transcoding and its quality reduction and delay penalty**
- › Complexity:
 - **Mixer has no end-point complexity per end-point in the conference, only switching and some Mixer-specific logic**
 - Mixer proxies limitations in end-point affecting what sender can produce and receiver can accept



Source Projection Mixer



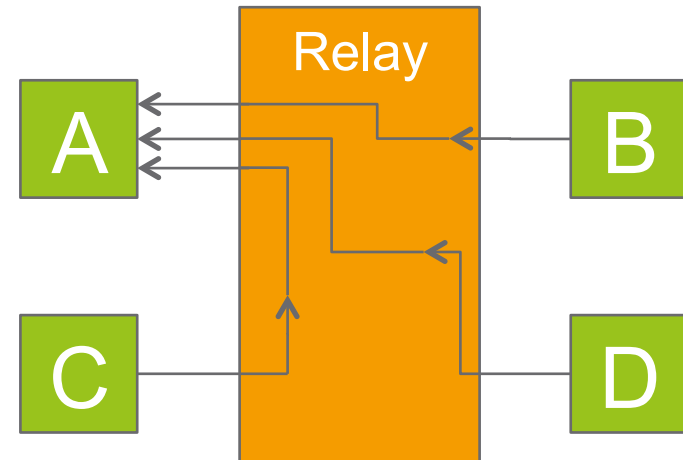
- › Description:
 - One peer communicate only with Mixer
 - Each peer to mixer communication is unicast with mixer feedback
 - **Each participant have its own RTP session with Mixer**
 - **Each conference media source is projected into each RTP session**
- › Security:
 - Mixer handles Authentication of each Peer / User
 - Mixer is trusted entity and enforcer of some security functions
 - Confidentiality and Integrity between peer and Mixer
- › Congestion:
 - Receiver (Mixer or peer) can report statistics or request adaptation from sender (Mixer or peer) on the closest link
 - **Mixer needs to aggregate and forward report / request information between links, based on some policy**
 - **Mixer distributes encoded media to multiple peers, making single receiver limitation affect more receivers**
 - Mixer can make use of simulcast or scalable media encoders from senders to adapt to a peer
- › Identification:
 - **Each media source is one to one mapped to a SSRC in Participants RTP session**
 - Sender's SSRC may be renumbered by Mixer, thus requiring RTP-external identification for E2E identity
- › Bandwidth:
 - Receiver (also Mixer) can request media to be tailored to its needs, but will typically also affect other receivers
 - Desirable to **limit** the amount of trade-off based on path capacity
 - Simulcast and scalable encoding can be used to meet different bandwidth needs or requirements
- › Quality:
 - Trade-off of end-to-end Path Capacity and Properties between receivers sharing media from the same sender
 - **Avoids transcoding and its quality reduction and delay penalty**
- › Complexity:
 - **Mixer has no end-point complexity per end-point in the conference, only switching and some Mixer-specific logic**
 - Mixer proxies limitations in end-point affecting what sender can produce and receiver can accept



Relay (Transport Translator)



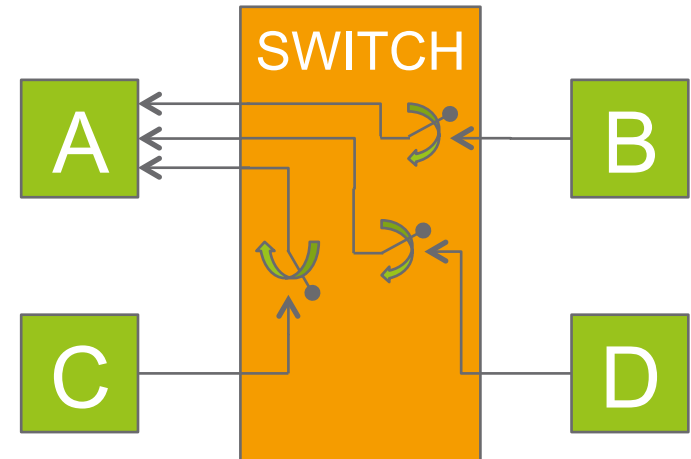
- › Description:
 - **One peer transmits only to the Relay, which forwards to multiple peers**
 - Each peer to Relay communication is unicast
 - **Relay creates a conference-wide RTP session**
- › Security:
 - **SRTP's regular source authentication can't authenticate peers**
 - › **For cryptographic verification TESLA or similar is needed**
 - Confidentiality and Integrity shared with all end-points
 - **Switch need not be trusted with media content**
 - Additional Keying mechanisms likely needed
- › Congestion:
 - **Each sender must aggregate receiver statistics reports or requests from all receivers**
 - **All Peers will share available capacity on all paths**
 - Any encoding changes due to congestion will affect all peers
- › Identification:
 - Sender's sources map one to one with RTP media streams
- › Bandwidth:
 - **Receiver bandwidth will always be the lowest common denominator from all paths**
 - Bandwidth optimizations must occur over whole conference not for individual paths
- › Quality:
 - Will be the lowest common denominator based on Capacity and Properties for all Paths
- › Complexity:
 - All in the end-points
 - Limitations in end-point directly affect what sender can produce and receiver can accept
 - Conference properties decided by lowest common denominator of peers



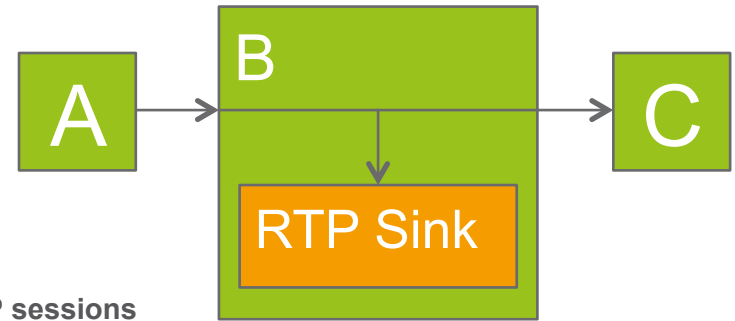
Selective Forwarding Switch



- › Description:
 - One peer communicate only with Switch
 - Each peer to Switch communication is unicast
 - **Switch creates a conference-wide RTP session**
 - **Switch turns individual source on and off based on some policy**
 - **Not supported by today's RTP!**
- › Security:
 - SRTP's regular source authentication can't authenticate individual peers
 - › For cryptographic verification TESLA is needed
 - Confidentiality and Integrity shared with all end-points
 - Switch need not be trusted with media content
 - Additional Keying mechanisms likely needed
 - Switching a source off and later on can break SRTP Roll over Counter
- › Congestion:
 - Each sender must aggregate receiver statistics reports or requests from all receivers
 - All Peers will share available capacity on all paths
 - Any encoding changes due to congestion will affect all peers
 - **Reporting and thus congestion detection will be confused by disappearing and reappearing sources**
- › Identification:
 - Sender's sources map one to one with RTP media streams
- › Bandwidth:
 - Receiver can request media to be tailored to its needs, but will typically also affect other receivers
 - **Which media streams an end-point receives can be individually tailored**
 - Desirable to limit the amount of trade-off based on path capacity
 - Simulcast and scalable encoders can be used to meet different bandwidth needs or requirements
- › Quality:
 - Trade-off of end-to-end Path Capacity and Properties between receivers sharing media from the same sender
- › Complexity:
 - **Switch has *no* end-point complexity per end-point in the conference, only forwarding logic and tables**



End-point Forwarding

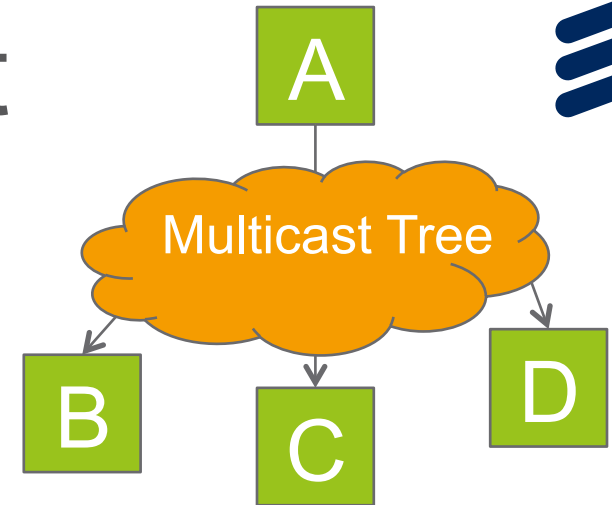


- › Description:
 - **One peer communicate only with peer, forwarding to other peers**
 - Each peer to peer communication is unicast
 - **If only forwarding RTP then a common RTP session is created**
 - **If B implements transcoding / RTP mixer functionality multiple RTP sessions can be created (Not further discussed, see Mixers)**
- › Security:
 - SRTP's regular source authentication can't authenticate individual peers
 - › For cryptographic verification TESLA is needed
 - Confidentiality and Integrity shared with all end-points
 - Additional Keying mechanisms could be used to avoid decryption / encryption cycle in B
- › Congestion (from A's perspective):
 - Sender must aggregate receiver statistics reports or requests from all receivers
 - All Peers will share available capacity on shared paths
 - Any encoding changes due to congestion will affect all peers
- › Identification:
 - Sender's sources map one to one with RTP media streams
- › Bandwidth:
 - Receiver bandwidth will always be the lowest common denominator from all paths
 - Bandwidth optimizations must occur over whole conference not for individual paths
- › Quality:
 - Will be the lowest common denominator based on Capacity and Properties for all Paths
- › Complexity:
 - All in the end-points, with some added complexity in B
 - Limitations in end-point directly affect what sender can produce and receiver can accept
 - Conference properties decided by lowest common denominator of peers

Any Source Multicast



- › Description:
 - **One peer communicate with all multicast group members**
 - **Multicast group is a conference-wide RTP session**
- › Security:
 - SRTP's regular source authentication can't authenticate individual peers
 - › For cryptographic verification TESLA is needed
 - Confidentiality and Integrity shared with all peers
- › Congestion:
 - **Each sender must aggregate receiver statistics reports or requests from all receivers**
 - **All Peers will share (single copy) available capacity on all links**
 - Any encoding changes due to congestion will affect all peers
- › Identification:
 - Sender's sources map one to one with RTP media streams
- › Bandwidth:
 - Receiver bandwidth will always be the lowest common denominator from all paths
 - **Bandwidth optimizations for a single multicast group must occur over whole conference not for individual paths**
 - Bandwidth adaptation can be achieved using multiple multicast groups and simulcast or scalability
- › Quality:
 - Will be the lowest common denominator based on Capacity and Properties for all Paths
- › Complexity:
 - All in the end-points
 - Limitations in end-point directly affect what sender can produce and receiver can accept
 - Conference properties decided by lowest common denominator of peers



Source Specific Multicast



› Description:

- A SSM tree enables media delivery to a number of receivers from aggregation point
- Media sources may be mixed, switched, selected etc. to generate media streams sent over SSM
- A receiver of the SSM media provides feedback (RTCP) over unicast
- If a receiver likes to send media it must be sent to media aggregator using separated unicast traffic

› Implications are left as an exercise ;-)

Conclusions



- › There are many topologies
 - Most, if not all are valid implementation choices for CLUE systems
- › Difficult to select trade-offs to optimize conference
- › Do we need to select supported topologies?
- › Does CLUE signaling need to take all into consideration?