A Content Fetching Layer over NDN (with demo)

Jun Bi

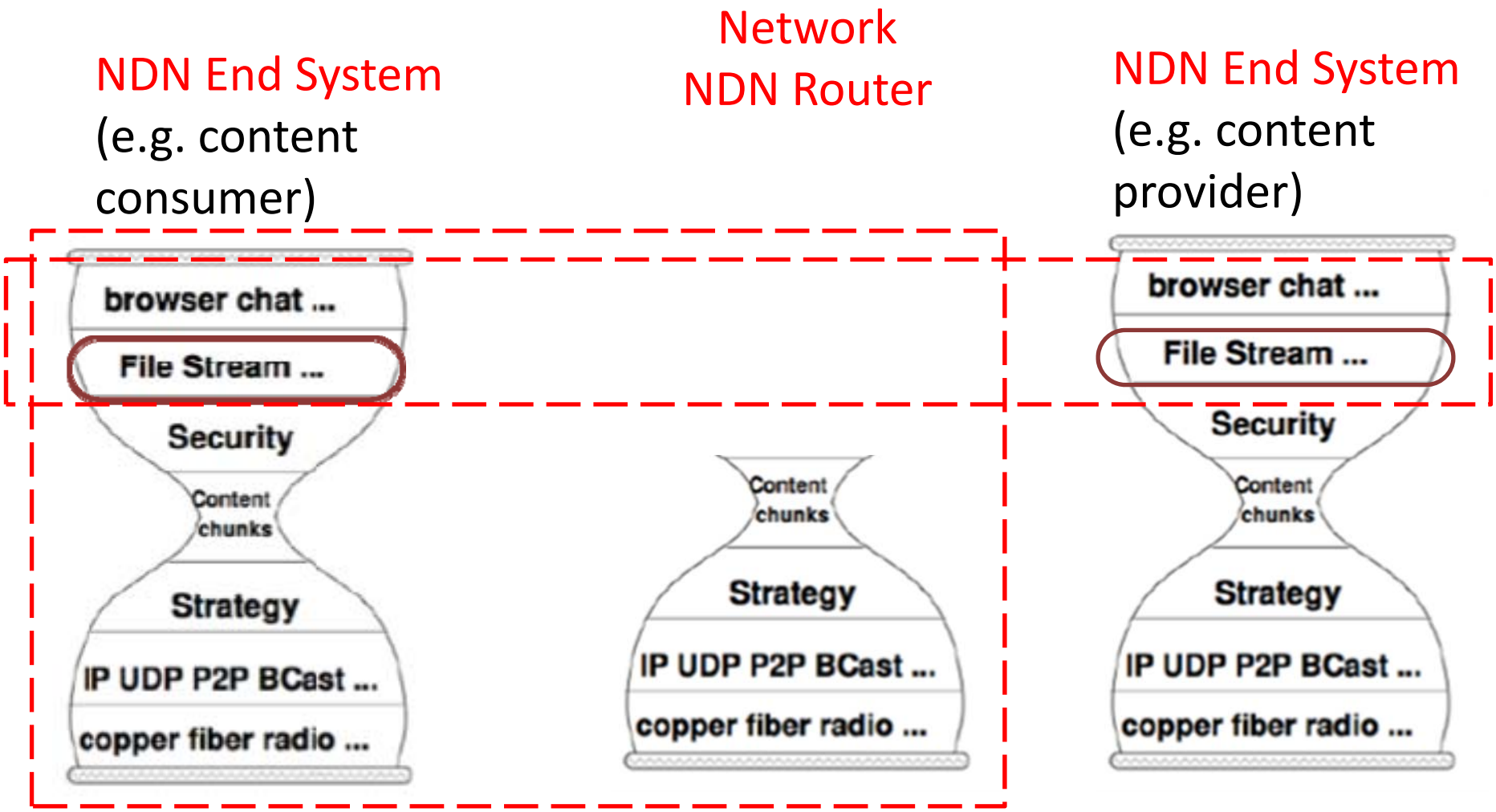
Xiaoke Jiang, Zhaogeng Li, Ao Guo

Tsinghua University/CERNET

Motivation

- In NDN application layer, applications need to get a sequence of chunks of “application content” (content concept at application layer)
  - E.g, File: error-free
  - Video streaming: real-time
- So we need a layer over NDN, to serve for applications to fetch “application content” with better performance
  - We call it *Content Fetch Layer*

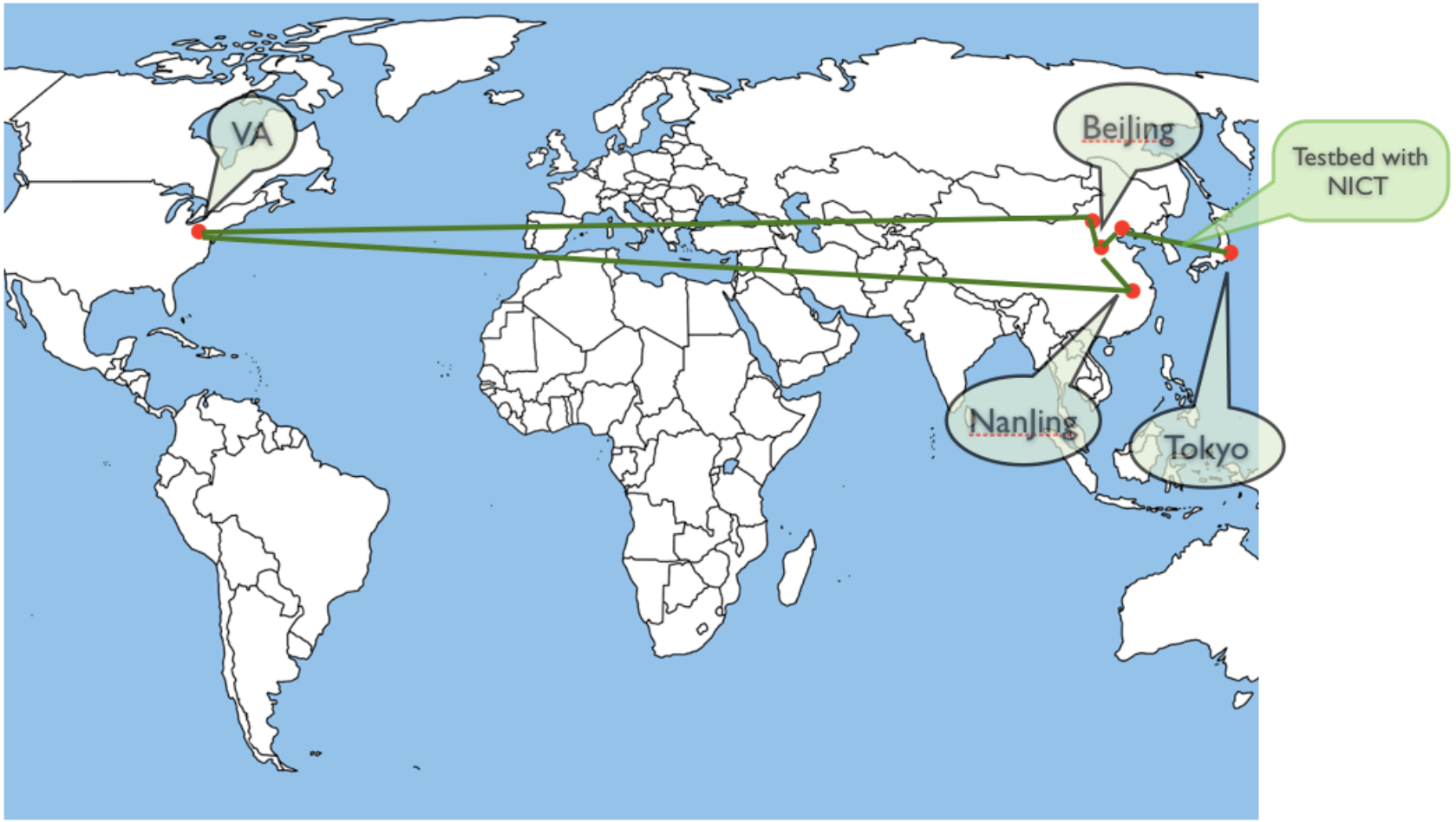
Motivation


NDN End System
(e.g. content consumer)
Network
NDN Router
NDN End System
(e.g. content provider)
browser chat ...
File Stream ...
Security
Content chunks
Strategy
IP UDP P2P BCast ...
copper fiber radio ...
Content chunks
Strategy
IP UDP P2P BCast ...
copper fiber radio ...
browser chat ...
File Stream ...
Security
Content chunks
Strategy
IP UDP P2P BCast ...
copper fiber radio ...

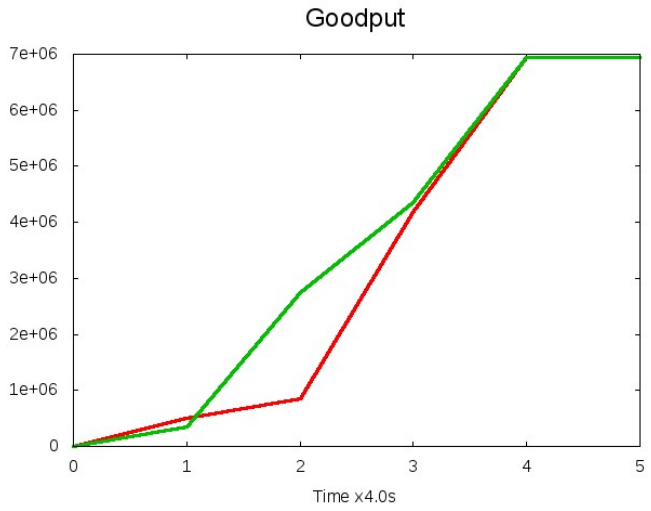
Content Fetching Layer

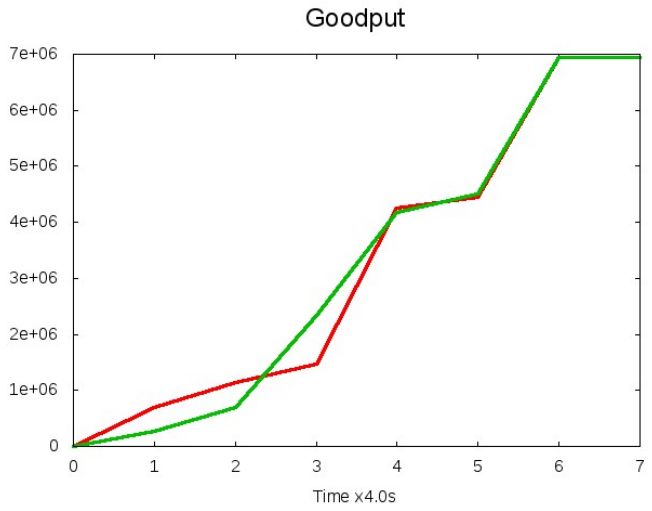
- To improve performance for retrieving a sequence of chunks for the same “application content”
  - consumer-to-network mode (for most applications)
    - Adaptive Congestion Window
      - Slow Start, AIMD; or new mechanisms
    - Adaptive Interests’ lifetime (RTO: Retransmission Timeout)
      - RTO = α*SRTT+β*RTTVar
      - α makes NDN routers have more time to try alternative paths for interests, β makes the RTO adaptive to network condition (faster re-retransmission request)
  - end-to-end mode (only for point-point applications) and (might) provider-to-network mode (e..g online video provider)
    - Adaptive Chunk Size
      - optimized by keep tracking chunk loss rate

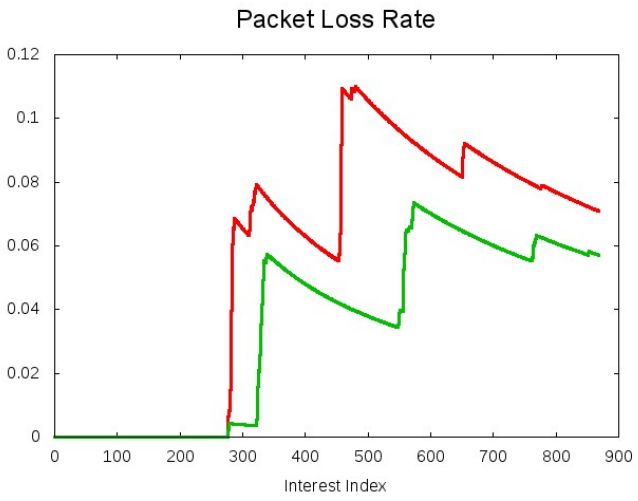
Experiment Topology (Implemented with NDNx)

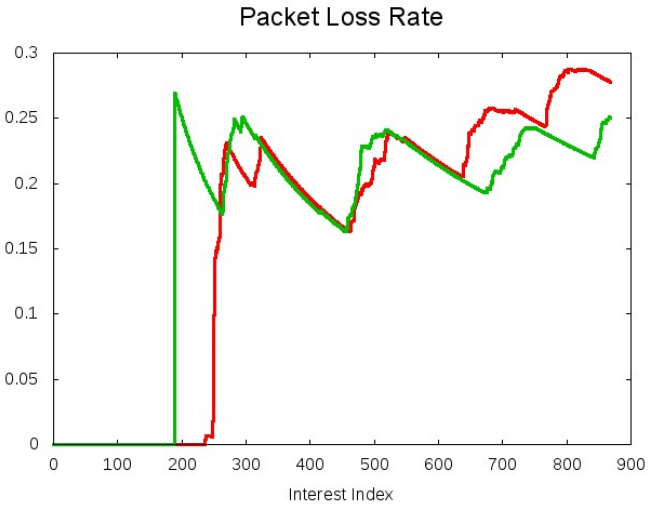


NDN without “content fetch layer”









WIN=256

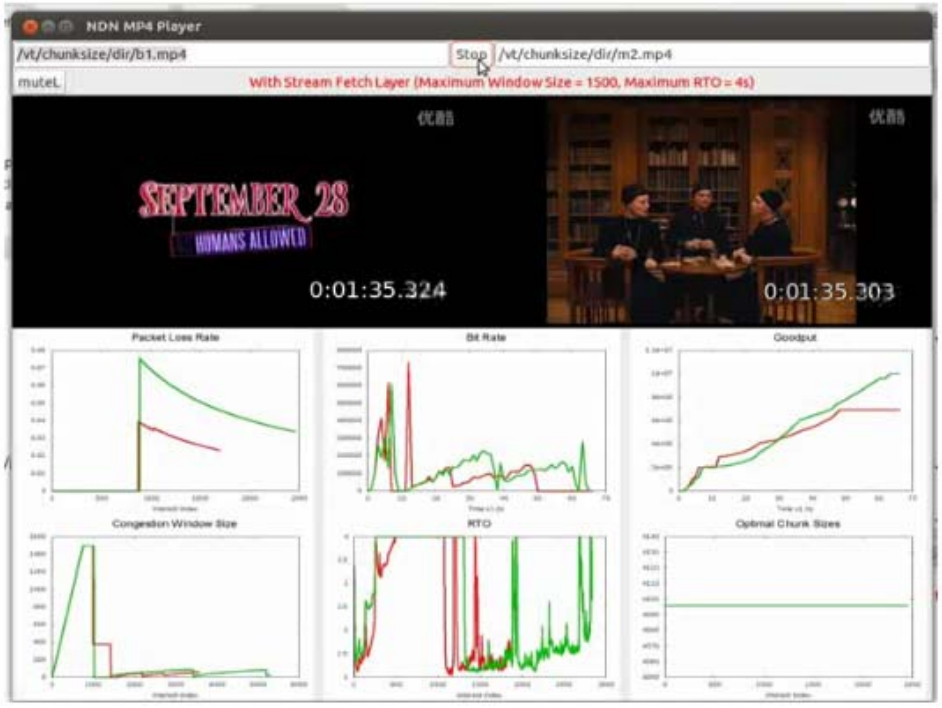
WIN=512

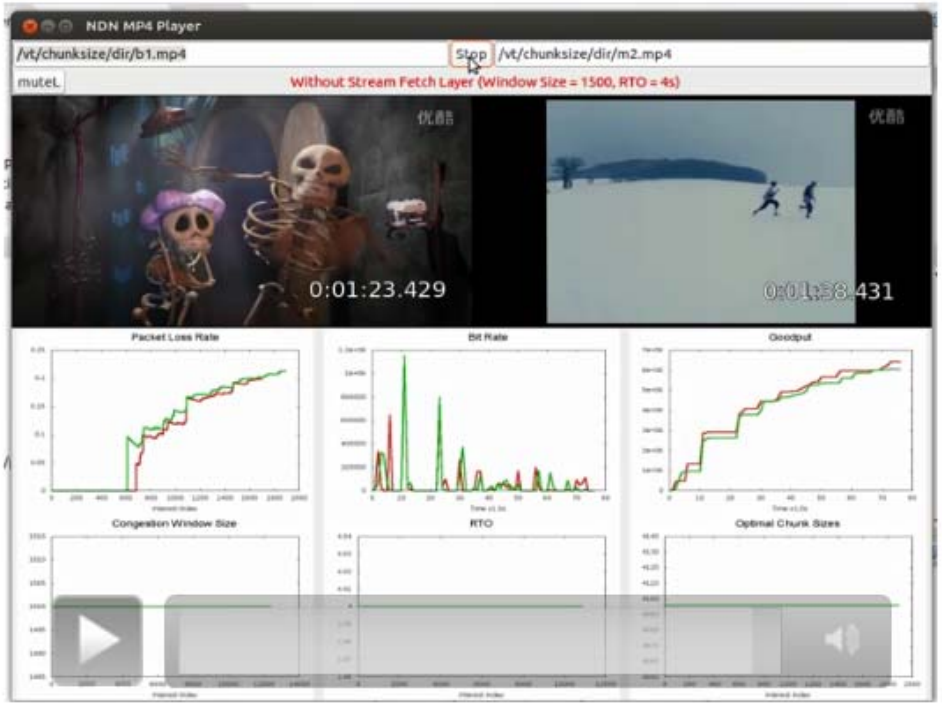
High Loss Rate

Demo: Adaptive Window-size/RTO

Max WindowSize=1500
Max RTO=4s
ChunkSize=4KB

Fixed WindowSize=1500
Fixed RTO=4s
ChunkSize=4KB

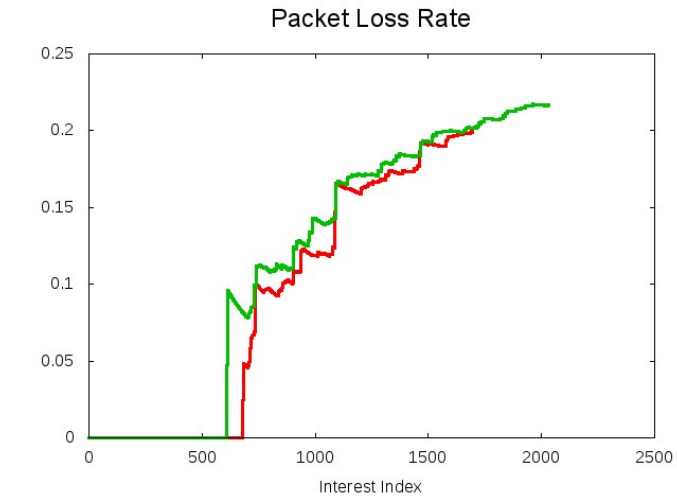


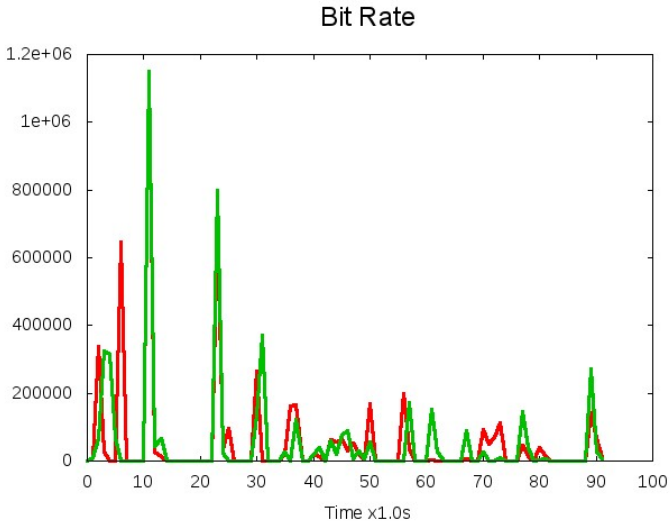


Loss Rate Bit Rate Goodput

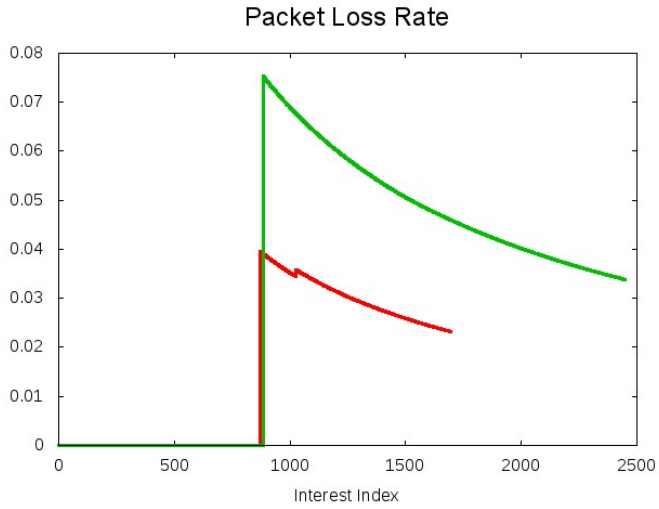
Window Size RTO Chunk Size

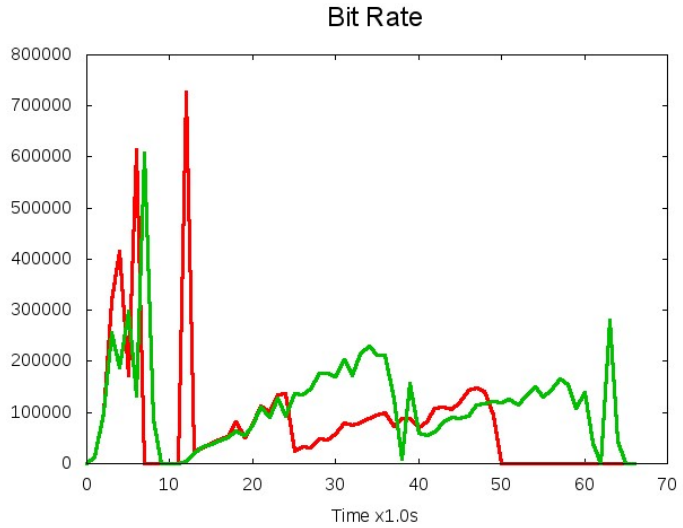
Comparison 1





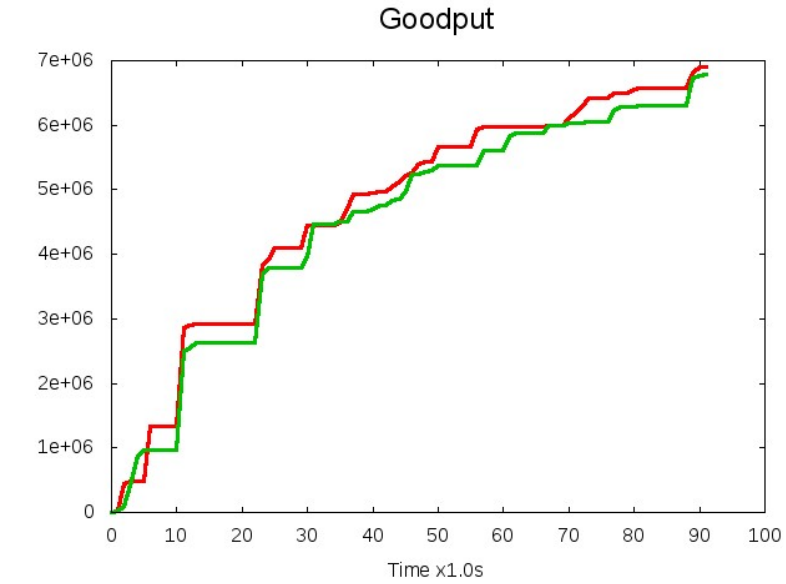
Fixed

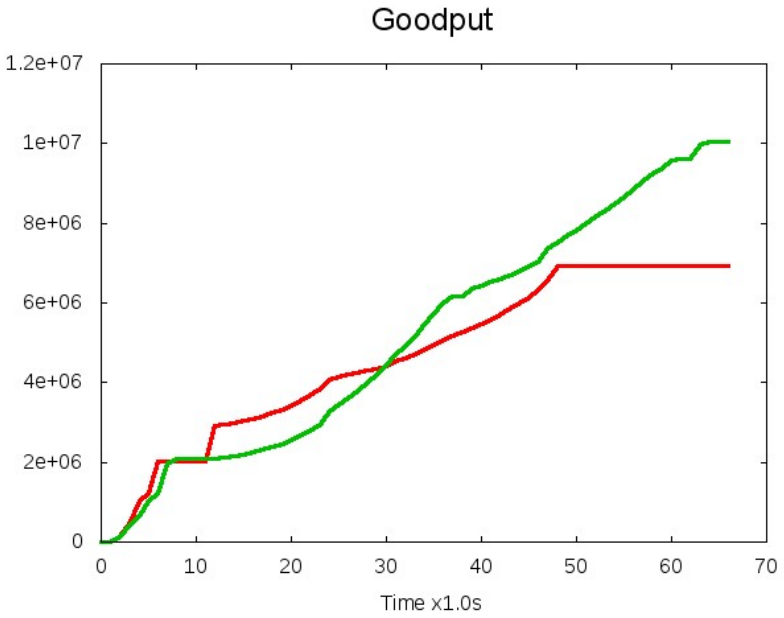




Adaptive

Comparison 2





Fixed WindowSize=1500
Fixed RTO=4s
ChunkSize=4KB

Adaptive window size and RTO will improve the Goodput.

Max WindowSize =1500
Max RTO=4s
ChunkSize=4KB

Adaptive Chunk Size Mechanism

- Chunk Size
  - too small: low payload ratio
    - large header overhead, ~500B: signature, key locator, ...
  - too big: high chunk loss rate and more re-transmission
    - underlying layers need to do fragmentation
    - one piece of fragmentation loss will cause the whole big chunk loss

Why Adaptive Chunk Size

- Say, the chunk header size is 1, underlying packet size is 3; and size of requested "application content" is 8. underlying packet loss rate p
  - Solution1: 4 chunks:  chunk size=2+1
  - Solution2: 1 chunks:  chunk size=8+1

	Solution 1	Solution 2
Payload Ratio	2/3=66.7%	8/9=88.9% ✔
Chunk Loss Rate (p=1%)	1% per chunk ✔	1-(1-p)^3 = 2.9%
Chunk Loss Rate (p=3%)	3% ✔	8.7% ✗
Size if Retransmission	3	9 ✗

Demo: Adaptive Chunk Size Mechanism

Summary

- We may need a “content fetch layer” at end systems to better serve applications
- Some possible mechanisms in this content fetch layer were implemented in our prototype and testbed
  - end-to-network mode (for most applications)
    - Adaptive Congestion Window
    - Adaptive Interests’ lifetime (RTO)
  - end-to-end mode (Only for point-point applications) and (might) provider-to-network mode (e..g online video provider)
    - Adaptive Chunk Size
  - Maybe more functions

Thanks !