

# Routing State Abstraction

## Using Declarative Equivalence

draft-gao-alto-routing-state-abstraction-01

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- ▶ A general objective of ALTO is to provide generic network state to applications for better traffic optimization
- ▶ It is important that ALTO provide **abstract network state**
  - ▶ Protect information privacy
  - ▶ Improve scalability



- ▶ The current ALTO standard can provide
  - ▶ any network information for a **single flow**
  - ▶ **flow-irrelevant** network information for **multiple flows**, such as *hopcount*
  - ▶ statistical network information based on the **Law of Large Numbers** for **multiple flows**, such as the average RTT between PIDs

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- ▶ Generally speaking, where the decisions for each flow are **independent**

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  - ▶ Map-Reduce scheduling in data centers
  - ▶ Traffic engineering in an ISP network
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- ▶ However, many applications require multi-flow coordination
  - ▶ Map-Reduce scheduling in data centers
  - ▶ Traffic engineering in an ISP network
  - ▶ ...
- ▶ Path vector can solve this by providing **network state** with common **network elements** for all the flows
  - ▶ **network element**: link/AS/...
  - ▶ **network state**: properties/statistics/...

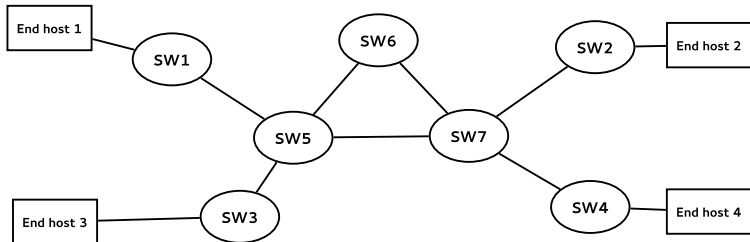


Figure: Example Topology

```
"PID1": {  
  "PID2": ["ne15", "ne56", "ne67", "ne27"],  
  "PID4": ["ne15", "ne57", "ne47"]  
},  
"PID2": {  
  "PID1": ["ne27", "ne57", "ne15"],  
  "PID3": ["ne27", "ne57", "ne35"]  
},
```

```
"PID3": {  
  "PID2": ["ne35", "ne57", "ne27"],  
  "PID4": ["ne35", "ne57", "ne47"]  
},  
"PID4": {  
  "PID1": ["ne47", "ne75", "ne15"],  
  "PID3": ["ne47", "ne57", "ne35"]  
}
```

How to compute **abstract** network state



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- ▶ Return **dynamic** network state
- ▶ Return **minimal** network state
- ▶ Return **equivalent** network state

## A Generic Definition:

*The abstract network state **A** for a user request is **equivalent** to the raw network state **R**, if and only if the user can make the same optimized decision with **A** as with **R**.*

The RSADE, **Routing State Abstraction using Declarative Equivalence**, is proposed to provide such a network state abstraction service for a certain family of optimization: utilizing the objective function.

### Objective Function

*An expression containing **variables** and mathematical constants.*

### Variable

*Just like a math variable but usually has a specific physical significance.*

## Optimal Equivalence

*If the object function has the same solution using  $\mathbf{A}$  and  $\mathbf{R}$ ,  $\mathbf{A}$  and  $\mathbf{R}$  are considered **equivalent**.*

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## Range Equivalence

*If the values of all linear combinations of variables, computed with  $\mathbf{A}$  and  $\mathbf{R}$ , have the same range,  $\mathbf{A}$  and  $\mathbf{R}$  are considered **equivalent**.*

## Flow descriptor

*Specify the relevant flows.*

- ▶ Legacy: use EndpointFilter
- ▶ New: use FlowFilter\*

## Equivalence Condition

*Describe how the network can effect the decision making.*

- ▶ In RSADE, we limit this to *linear inequalities per link*.



## Abstract Network State

- ▶ **Path vector**

Return path vectors and let application construct the constraints.

- ▶ **Constraints\***

Construct the constraints for the application using the format defined in *equivalence conditions* and return them.

## How to specify FlowFilter

- ▶ A list of flows
- ▶ Consider possible OpenFlow use case: use tuples instead of destinations alone
- ▶ Each flow can be described as a (src, dst) combination

FlowFilter := flow-list

flow-list := flow-spec, [flow-list]

flow-spec := generic-match-condition

- ▶ Extension 1: more advanced endpoint address descriptors
  - ▶ `draft-wang-alto-ecs-flows-00`
- ▶ Extension 2: more fields in the flow specification
  - ▶ Examples: **web-proxy**, **qos-group**, etc.

## How to specify Equivalence Conditions

- ▶ Two kinds of inequalities
  - ▶ network-irrelevant: the constraint is application-specific
  - ▶ network-relevant: the constraint uses properties in the network
- ▶ Consider the structure of a network-relevant inequality
  - ▶ Math constants (provided by application)
  - ▶ Variables (provided by application)
  - ▶ Link properties (provided by network)
  - ▶ The routing information (provided by network)

$$R[1] * \text{flow1} + R[2] * \text{flow2} \leq 0.8 * \text{bandwidth}$$

- ▶ Optional: provide the objective function

```

equiv-cond           := variable-list X0 link-constraint-list

variable-list        := variable-name[, variable-list]

X0                   := simple-constraint[, simple-constraint]
simple-constraint     := simple-expr CMP-OP simple-expr
simple-expr           := constant * variable-name[ + simple-expr]

link-constraints-list := link-constraint[, link-constraint-list]
link-constraint       := link-expr CMP-OP link-expr
link-expr              := constant | attribute-name | variable-name
                       | constant * link-expr
                       | attribute-name * link-expr
                       | link-expr + link-expr
  
```

See [draft-gao-alto-routing-state-abstraction-01](#) for details.

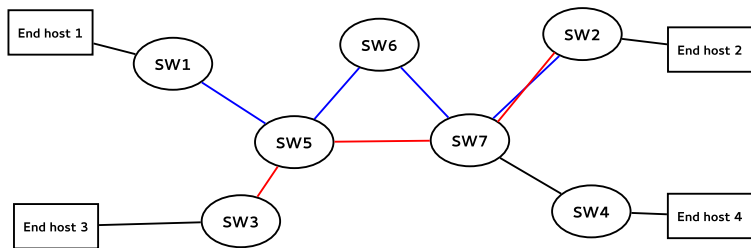


Figure: Example Topology

Assume each link is 100Mbps and apply

- ▶ Flow descriptor:  
flows eh1→eh2(blue) and eh3→eh2(red)
- ▶ Equivalence condition:  
 $R[1] * flow1 + R[2] * flow2 \leq bandwidth$

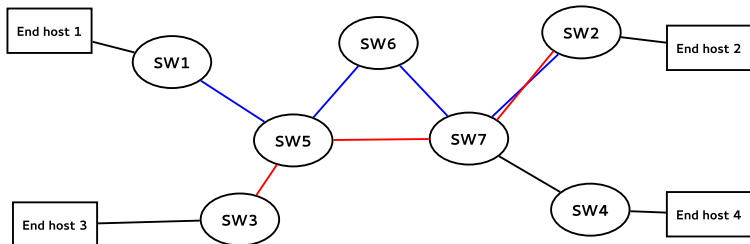


Figure: Example Topology

We get

```

ne15: 1 * flow1 + 0 * flow2 <= 100M
ne56: 1 * flow1 + 0 * flow2 <= 100M
ne67: 1 * flow1 + 0 * flow2 <= 100M
ne27: 1 * flow1 + 1 * flow2 <= 100M
ne57: 0 * flow1 + 1 * flow2 <= 100M
ne35: 0 * flow1 + 1 * flow2 <= 100M

```

In order to satisfy the **minimal** and **equivalent** criteria, we have defined the following terms:

[Equivalence] Two constraint sets  $\mathbf{S}_1 : \{\vec{x} | \mathbf{A}_1 \vec{x} \leq \vec{b}_1\}$  and  $\mathbf{S}_2 : \{\vec{x} | \mathbf{A}_2 \vec{x} \leq \vec{b}_2\}$  of a network function are equivalent if and only if they limit the decision variables in the same way:  $\mathbf{X}_0 \cap \mathbf{S}_1 = \mathbf{X}_0 \cap \mathbf{S}_2$ .

[Redundant] A constraint  $s$  is redundant to a constraint set  $\mathbf{S}$  if and only if  $s \in \mathbf{S}$  and the two sets  $\mathbf{S}$  and  $\mathbf{S} \setminus \{s\}$  are equivalent.

[Minimal Constraint Set] A constraint set  $\mathbf{S}$  is minimal if and only if  $\forall s \in \mathbf{S}$ ,  $s$  is not redundant.



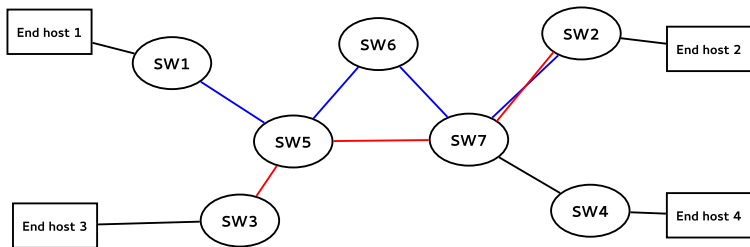


Figure: Example Topology

The minimal constraint set is

$$\text{ne27: } 1 * \text{flow1} + 1 * \text{flow2} \leq 100\text{M}$$

And the corresponding **path vector** response is

eh1 -> eh2: [ ne27 ],

eh3 -> eh2: [ ne27 ]

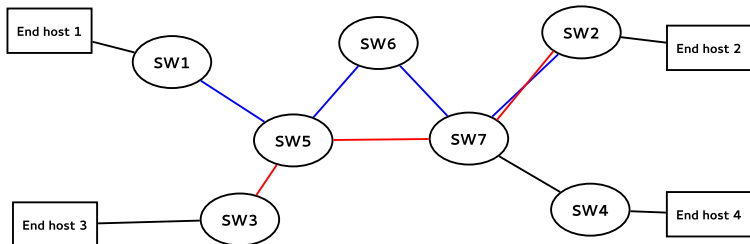


Figure: Example Topology

Change the bandwidth of ne57 to 70Mbps, we have

$$\text{ne15: } 1 * \text{flow1} + 0 * \text{flow2} \leq 100\text{M}$$

$$\text{ne56: } 1 * \text{flow1} + 0 * \text{flow2} \leq 100\text{M}$$

$$\text{ne67: } 1 * \text{flow1} + 0 * \text{flow2} \leq 100\text{M}$$

$$\text{ne27: } 1 * \text{flow1} + 1 * \text{flow2} \leq 100\text{M}$$

$$\text{ne57: } 0 * \text{flow1} + 1 * \text{flow2} \leq 70\text{M}$$

$$\text{ne35: } 0 * \text{flow1} + 1 * \text{flow2} \leq 100\text{M}$$

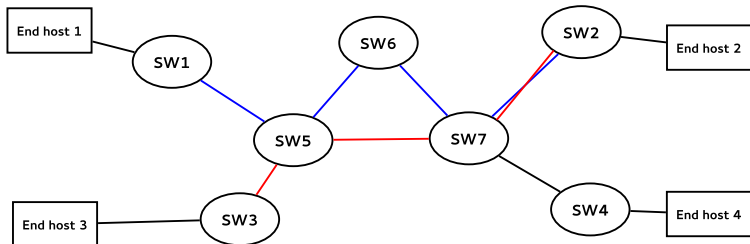


Figure: Example Topology

In this case, the minimal constraint set is

$$\text{ne27: } 1 * \text{flow1} + 1 * \text{flow2} \leq 100\text{M}$$

$$\text{ne57: } 0 * \text{flow1} + 1 * \text{flow2} \leq 70\text{M}$$

And the corresponding **path vector** response is

$$\text{eh1} \rightarrow \text{eh2: } [ \text{ne27} ],$$

$$\text{eh3} \rightarrow \text{eh2: } [ \text{ne27}, \text{ne57} ]$$

The **path vector** form of the second response in the example is demonstrated below:

```
"endpoint-cost-map": {
  "eh1": [ "eh2" : [ "ane1" ] ],
  "eh3": [ "eh2" : [ "ane1", "ane2" ] ]
},
"network-elements": {
  "ane1": { "bandwidth": "100 Mbps" },
  "ane2": { "bandwidth": "70 Mbps" }
}
```

The **constraint** form of the second response in the example is demonstrated below:

```
"flow-constraints": [  
  "flow1 + flow2 <= 100000000",  
  "flow2 <= 700000000"  
]
```

Thank you